

# Yosys Manual

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# Abstract

Most of today's digital design is done in HDL code (mostly Verilog or VHDL) and with the help of HDL synthesis tools.

In special cases such as synthesis for coarse-grain cell libraries or when testing new synthesis algorithms it might be necessary to write a custom HDL synthesis tool or add new features to an existing one. In these cases the availability of a Free and Open Source (FOSS) synthesis tool that can be used as basis for custom tools would be helpful.

In the absence of such a tool, the Yosys Open SYnthesis Suite (Yosys) was developed. This document covers the design and implementation of this tool. At the moment the main focus of Yosys lies on the high-level aspects of digital synthesis. The pre-existing FOSS logic-synthesis tool ABC is used by Yosys to perform advanced gate-level optimizations.

An evaluation of Yosys based on real-world designs is included. It is shown that Yosys can be used as-is to synthesize such designs. The results produced by Yosys in this tests were successfully verified using formal verification and are comparable in quality to the results produced by a commercial synthesis tool.

This document was originally published as bachelor thesis at the Vienna University of Technology [Wol13].

# Abbreviations

AIG	And-Inverter-Graph
ASIC	Application-Specific Integrated Circuit
AST	Abstract Syntax Tree
BDD	Binary Decision Diagram
BLIF	Berkeley Logic Interchange Format
EDA	Electronic Design Automation
EDIF	Electronic Design Interchange Format
ER Diagram	Entity-Relationship Diagram
FOSS	Free and Open-Source Software
FPGA	Field-Programmable Gate Array
FSM	Finite-state machine
HDL	Hardware Description Language
LPM	Library of Parameterized Modules
RTLIL	RTL Intermediate Language
RTL	Register Transfer Level
SAT	Satisfiability Problem
VHDL	VHSIC Hardware Description Language
VHSIC	Very-High-Speed Integrated Circuit
YOSYS	Yosys Open SYnthesis Suite

# Contents

<b>1</b>	<b>Introduction</b>	<b>15</b>
1.1	History of Yosys . . . . .	15
1.2	Structure of this Document . . . . .	16
<b>2</b>	<b>Basic Principles</b>	<b>17</b>
2.1	Levels of Abstraction . . . . .	17
2.1.1	System Level . . . . .	18
2.1.2	High Level . . . . .	18
2.1.3	Behavioural Level . . . . .	18
2.1.4	Register-Transfer Level (RTL) . . . . .	19
2.1.5	Logical Gate Level . . . . .	19
2.1.6	Physical Gate Level . . . . .	20
2.1.7	Switch Level . . . . .	20
2.1.8	Yosys . . . . .	20
2.2	Features of Synthesizable Verilog . . . . .	20
2.2.1	Structural Verilog . . . . .	21
2.2.2	Expressions in Verilog . . . . .	21
2.2.3	Behavioural Modelling . . . . .	21
2.2.4	Functions and Tasks . . . . .	22
2.2.5	Conditionals, Loops and Generate-Statements . . . . .	22
2.2.6	Arrays and Memories . . . . .	23
2.3	Challenges in Digital Circuit Synthesis . . . . .	23
2.3.1	Standards Compliance . . . . .	23
2.3.2	Optimizations . . . . .	24
2.3.3	Technology Mapping . . . . .	24
2.4	Script-Based Synthesis Flows . . . . .	24
2.5	Methods from Compiler Design . . . . .	25
2.5.1	Lexing and Parsing . . . . .	25
2.5.2	Multi-Pass Compilation . . . . .	27

## CONTENTS

<b>3</b>	<b>Approach</b>	<b>28</b>
3.1	Data- and Control-Flow	28
3.2	Internal Formats in Yosys	29
3.3	Typical Use Case	29
<b>4</b>	<b>Implementation Overview</b>	<b>31</b>
4.1	Simplified Data Flow	31
4.2	The RTL Intermediate Language	32
4.2.1	RTLIL Identifiers	33
4.2.2	RTLIL::Design and RTLIL::Module	34
4.2.3	RTLIL::Cell and RTLIL::Wire	35
4.2.4	RTLIL::SigSpec	35
4.2.5	RTLIL::Process	36
4.2.6	RTLIL::Memory	38
4.3	Command Interface and Synthesis Scripts	39
4.4	Source Tree and Build System	39
<b>5</b>	<b>Internal Cell Library</b>	<b>41</b>
5.1	RTL Cells	41
5.1.1	Unary Operators	41
5.1.2	Binary Operators	42
5.1.3	Multiplexers	43
5.1.4	Registers	44
5.1.5	Memories	46
5.1.6	Finite State Machines	50
5.1.7	Specify rules	50
5.1.8	Formal verification cells	50
5.2	Gates	50
<b>6</b>	<b>Programming Yosys Extensions</b>	<b>56</b>
6.1	Guidelines	56
6.2	The “stubsnets” Example Module	61

## CONTENTS

<b>7</b>	<b>The Verilog and AST Frontends</b>	<b>65</b>
7.1	Transforming Verilog to AST	65
7.1.1	The Verilog Preprocessor	66
7.1.2	The Verilog Lexer	66
7.1.3	The Verilog Parser	66
7.2	Transforming AST to RTLIL	67
7.2.1	AST Simplification	67
7.2.2	Generating RTLIL	69
7.3	Synthesizing Verilog always Blocks	69
7.3.1	The ProcessGenerator Algorithm	71
7.3.2	The proc pass	74
7.4	Synthesizing Verilog Arrays	74
7.5	Synthesizing Parametric Designs	74
<b>8</b>	<b>Optimizations</b>	<b>75</b>
8.1	Simple Optimizations	75
8.1.1	The opt_expr pass	75
8.1.2	The opt_muxtree pass	76
8.1.3	The opt_reduce pass	76
8.1.4	The opt_rmdff pass	77
8.1.5	The opt_clean pass	77
8.1.6	The opt_merge pass	77
8.2	FSM Extraction and Encoding	77
8.2.1	FSM Detection	78
8.2.2	FSM Extraction	78
8.2.3	FSM Optimization	79
8.2.4	FSM Recoding	80
8.3	Logic Optimization	80
<b>9</b>	<b>Technology Mapping</b>	<b>81</b>
9.1	Cell Substitution	81
9.2	Subcircuit Substitution	81
9.3	Gate-Level Technology Mapping	82
<b>A</b>	<b>Auxiliary Libraries</b>	<b>83</b>
A.1	SHA1	83
A.2	BigInt	83
A.3	SubCircuit	83
A.4	ezSAT	83

## CONTENTS

<b>B</b>	<b>Auxiliary Programs</b>	<b>84</b>
B.1	yosys-config . . . . .	84
B.2	yosys-filterlib . . . . .	84
B.3	yosys-abc . . . . .	84
<b>C</b>	<b>Command Reference Manual</b>	<b>85</b>
C.1	abc – use ABC for technology mapping . . . . .	85
C.2	abc9 – use ABC9 for technology mapping . . . . .	88
C.3	abc9_exe – use ABC9 for technology mapping . . . . .	91
C.4	abc9_ops – helper functions for ABC9 . . . . .	92
C.5	add – add objects to the design . . . . .	94
C.6	aigmap – map logic to and-inverter-graph circuit . . . . .	94
C.7	alumacc – extract ALU and MACC cells . . . . .	95
C.8	anlogic_eqn – Anlogic: Calculate equations for luts . . . . .	95
C.9	anlogic_fixcarry – Anlogic: fix carry chain . . . . .	95
C.10	assertpmux – adds asserts for parallel muxes . . . . .	95
C.11	async2sync – convert async FF inputs to sync circuits . . . . .	96
C.12	attrmap – renaming attributes . . . . .	96
C.13	attrmvcp – move or copy attributes from wires to driving cells . . . . .	96
C.14	autoname – automatically assign names to objects . . . . .	97
C.15	blackbox – convert modules into blackbox modules . . . . .	97
C.16	bmuxmap – transform \$bmux cells to trees of \$mux cells . . . . .	97
C.17	bugpoint – minimize testcases . . . . .	97
C.18	cd – a shortcut for 'select -module <name>' . . . . .	99
C.19	check – check for obvious problems in the design . . . . .	99
C.20	chformal – change formal constraints of the design . . . . .	100
C.21	chparam – re-evaluate modules with new parameters . . . . .	100
C.22	chtype – change type of cells in the design . . . . .	101
C.23	clean – remove unused cells and wires . . . . .	101
C.24	clean_zerowidth – clean zero-width connections from the design . . . . .	101
C.25	clk2fflogic – convert clocked FFs to generic \$ff cells . . . . .	101
C.26	clkbufmap – insert clock buffers on clock networks . . . . .	102
C.27	connect – create or remove connections . . . . .	102
C.28	connect_rpc – connect to RPC frontend . . . . .	103
C.29	connwrappers – match width of input-output port pairs . . . . .	104
C.30	coolrunner2_fixup – insert necessary buffer cells for CoolRunner-II architecture . . . . .	104
C.31	coolrunner2_sop – break \$sop cells into ANDTERM/ORTERM cells . . . . .	104

## CONTENTS

C.32 copy – copy modules in the design . . . . .	104
C.33 cover – print code coverage counters . . . . .	104
C.34 cutpoint – adds formal cut points to the design . . . . .	105
C.35 debug – run command with debug log messages enabled . . . . .	106
C.36 delete – delete objects in the design . . . . .	106
C.37 deminout – demote inout ports to input or output . . . . .	106
C.38 demuxmap – transform \$demux cells to \$eq + \$mux cells . . . . .	106
C.39 design – save, restore and reset current design . . . . .	106
C.40 dffinit – set INIT param on FF cells . . . . .	108
C.41 dfflegalize – convert FFs to types supported by the target . . . . .	108
C.42 dfflibmap – technology mapping of flip-flops . . . . .	109
C.43 dffunmap – unmap clock enable and synchronous reset from FFs . . . . .	110
C.44 dump – print parts of the design in RTLIL format . . . . .	110
C.45 echo – turning echoing back of commands on and off . . . . .	111
C.46 ecp5_gsr – ECP5: handle GSR . . . . .	111
C.47 edgetypes – list all types of edges in selection . . . . .	111
C.48 efinix_fixcarry – Efinix: fix carry chain . . . . .	111
C.49 equiv_add – add a \$equiv cell . . . . .	111
C.50 equiv_induct – proving \$equiv cells using temporal induction . . . . .	112
C.51 equiv_make – prepare a circuit for equivalence checking . . . . .	112
C.52 equiv_mark – mark equivalence checking regions . . . . .	113
C.53 equiv_miter – extract miter from equiv circuit . . . . .	113
C.54 equiv_opt – prove equivalence for optimized circuit . . . . .	113
C.55 equiv_purge – purge equivalence checking module . . . . .	114
C.56 equiv_remove – remove \$equiv cells . . . . .	114
C.57 equiv_simple – try proving simple \$equiv instances . . . . .	115
C.58 equiv_status – print status of equivalent checking module . . . . .	115
C.59 equiv_struct – structural equivalence checking . . . . .	115
C.60 eval – evaluate the circuit given an input . . . . .	116
C.61 exec – execute commands in the operating system shell . . . . .	116
C.62 expose – convert internal signals to module ports . . . . .	117
C.63 extract – find subcircuits and replace them with cells . . . . .	118
C.64 extract_counter – Extract GreenPak4 counter cells . . . . .	119
C.65 extract_fa – find and extract full/half adders . . . . .	120
C.66 extract_reduce – converts gate chains into \$reduce_* cells . . . . .	120
C.67 extractinv – extract explicit inverter cells for invertible cell pins . . . . .	120



## CONTENTS

C.68 flatten – flatten design . . . . .	121
C.69 flowmap – pack LUTs with FlowMap . . . . .	121
C.70 fmcombine – combine two instances of a cell into one . . . . .	122
C.71 fminit – set init values/sequences for formal . . . . .	122
C.72 freduce – perform functional reduction . . . . .	123
C.73 fsm – extract and optimize finite state machines . . . . .	123
C.74 fsm_detect – finding FSMs in design . . . . .	124
C.75 fsm_expand – expand FSM cells by merging logic into it . . . . .	124
C.76 fsm_export – exporting FSMs to KISS2 files . . . . .	125
C.77 fsm_extract – extracting FSMs in design . . . . .	125
C.78 fsm_info – print information on finite state machines . . . . .	125
C.79 fsm_map – mapping FSMs to basic logic . . . . .	126
C.80 fsm_opt – optimize finite state machines . . . . .	126
C.81 fsm_recode – recoding finite state machines . . . . .	126
C.82 fst2tb – generate testbench out of fst file . . . . .	126
C.83 gatamate_foldinv – fold inverters into Gatamate LUT trees . . . . .	127
C.84 glift – create GLIFT models and optimization problems . . . . .	127
C.85 greenpak4_dffinv – merge greenpak4 inverters and DFF/latches . . . . .	129
C.86 help – display help messages . . . . .	129
C.87 hierarchy – check, expand and clean up design hierarchy . . . . .	129
C.88 hilomap – technology mapping of constant hi- and/or lo-drivers . . . . .	131
C.89 history – show last interactive commands . . . . .	131
C.90 ice40_braminit – iCE40: perform SB_RAM40_4K initialization from file . . . . .	131
C.91 ice40_dsp – iCE40: map multipliers . . . . .	132
C.92 ice40_opt – iCE40: perform simple optimizations . . . . .	132
C.93 ice40_wrapcarry – iCE40: wrap carries . . . . .	132
C.94 insbuf – insert buffer cells for connected wires . . . . .	133
C.95 iopadmap – technology mapping of i/o pads (or buffers) . . . . .	133
C.96 jny – write design and metadata . . . . .	134
C.97 json – write design in JSON format . . . . .	134
C.98 log – print text and log files . . . . .	134
C.99 logger – set logger properties . . . . .	135
C.100ls – list modules or objects in modules . . . . .	136
C.101ltp – print longest topological path . . . . .	136
C.102lut2mux – convert \$lut to \$_MUX_ . . . . .	136
C.103naccmap – mapping macc cells . . . . .	136

## CONTENTS

C.104	memory – translate memories to basic cells . . . . .	137
C.105	memory_bmux2rom – convert muxes to ROMs . . . . .	137
C.106	memory_bram – map memories to block rams . . . . .	137
C.107	memory_collect – creating multi-port memory cells . . . . .	139
C.108	memory_dff – merge input/output DFFs into memory read ports . . . . .	139
C.109	memory_libmap – map memories to cells . . . . .	140
C.110	memory_map – translate multiport memories to basic cells . . . . .	140
C.111	memory_memx – emulate vlog sim behavior for mem ports . . . . .	141
C.112	memory_narrow – split up wide memory ports . . . . .	141
C.113	memory_nordff – extract read port FFs from memories . . . . .	141
C.114	memory_share – consolidate memory ports . . . . .	141
C.115	memory_unpack – unpack multi-port memory cells . . . . .	142
C.116	miter – automatically create a miter circuit . . . . .	142
C.117	mutate – generate or apply design mutations . . . . .	142
C.118	muxcover – cover trees of MUX cells with wider MUXes . . . . .	144
C.119	muxpack – \$mux/\$pmux cascades to \$pmux . . . . .	144
C.120	lutmap – map to LUTs of different sizes . . . . .	145
C.121	onehot – optimize \$eq cells for onehot signals . . . . .	145
C.122	bpt – perform simple optimizations . . . . .	145
C.123	bpt_clean – remove unused cells and wires . . . . .	146
C.124	bpt_demorgan – Optimize reductions with DeMorgan equivalents . . . . .	146
C.125	opt_dff – perform DFF optimizations . . . . .	146
C.126	opt_expr – perform const folding and simple expression rewriting . . . . .	147
C.127	opt_ffinv – push inverters through FFs . . . . .	147
C.128	opt_lut – optimize LUT cells . . . . .	148
C.129	opt_lut_ins – discard unused LUT inputs . . . . .	148
C.130	opt_mem – optimize memories . . . . .	148
C.131	opt_mem_feedback – convert memory read-to-write port feedback paths to write enables . . . . .	148
C.132	opt_mem_priority – remove priority relations between write ports that can never collide . . . . .	149
C.133	opt_mem_widen – optimize memories where all ports are wide . . . . .	149
C.134	opt_merge – consolidate identical cells . . . . .	149
C.135	opt_muxtree – eliminate dead trees in multiplexer trees . . . . .	149
C.136	opt_reduce – simplify large MUXes and AND/OR gates . . . . .	150
C.137	opt_share – merge mutually exclusive cells of the same type that share an input signal . . . . .	150
C.138	paramap – renaming cell parameters . . . . .	150
C.139	peepopt – collection of peephole optimizers . . . . .	151

## CONTENTS

C.140	plugin – load and list loaded plugins . . . . .	151
C.141	pmux2shiftx – transform \$pmux cells to \$shiftx cells . . . . .	151
C.142	pmuxtree – transform \$pmux cells to trees of \$mux cells . . . . .	152
C.143	portlist – list (top-level) ports . . . . .	152
C.144	prep – generic synthesis script . . . . .	152
C.145	printattrs – print attributes of selected objects . . . . .	153
C.146	proc – translate processes to netlists . . . . .	154
C.147	proc_arst – detect asynchronous resets . . . . .	154
C.148	proc_clean – remove empty parts of processes . . . . .	155
C.149	proc_dff – extract flip-flops from processes . . . . .	155
C.150	proc_dlatch – extract latches from processes . . . . .	155
C.151	proc_init – convert initial block to init attributes . . . . .	155
C.152	proc_memwr – extract memory writes from processes . . . . .	155
C.153	proc_mux – convert decision trees to multiplexers . . . . .	156
C.154	proc_prune – remove redundant assignments . . . . .	156
C.155	proc_rmdead – eliminate dead trees in decision trees . . . . .	156
C.156	proc_rom – convert switches to ROMs . . . . .	156
C.157	qbsat – solve a 2QBF-SAT problem in the circuit . . . . .	156
C.158	qwp – quadratic wirelength placer . . . . .	158
C.159	read – load HDL designs . . . . .	158
C.160	read_aiger – read AIGER file . . . . .	159
C.161	read_blif – read BLIF file . . . . .	159
C.162	read_ilang – (deprecated) alias of read_rtlil . . . . .	160
C.163	read_json – read JSON file . . . . .	160
C.164	read_liberty – read cells from liberty file . . . . .	160
C.165	read_rtlil – read modules from RTLIL file . . . . .	161
C.166	read_verilog – read modules from Verilog file . . . . .	161
C.167	rename – rename object in the design . . . . .	164
C.168	rmports – remove module ports with no connections . . . . .	165
C.169	sat – solve a SAT problem in the circuit . . . . .	165
C.170	scatter – add additional intermediate nets . . . . .	168
C.171	scc – detect strongly connected components (logic loops) . . . . .	168
C.172	scratchpad – get/set values in the scratchpad . . . . .	169
C.173	script – execute commands from file or wire . . . . .	170
C.174	select – modify and view the list of selected objects . . . . .	170
C.175	setattr – set/unset attributes on objects . . . . .	174

## CONTENTS

C.176	setparam – set/unset parameters on objects . . . . .	175
C.177	setundef – replace undef values with defined constants . . . . .	175
C.178	share – perform sat-based resource sharing . . . . .	176
C.179	shell – enter interactive command mode . . . . .	176
C.180	show – generate schematics using graphviz . . . . .	177
C.181	shregmap – map shift registers . . . . .	178
C.182	sim – simulate the circuit . . . . .	179
C.183	simplemap – mapping simple coarse-grain cells . . . . .	181
C.184	splice – create explicit splicing cells . . . . .	181
C.185	splitnets – split up multi-bit nets . . . . .	182
C.186	sta – perform static timing analysis . . . . .	183
C.187	stat – print some statistics . . . . .	183
C.188	submod – moving part of a module to a new submodule . . . . .	183
C.189	supercover – add hi/lo cover cells for each wire bit . . . . .	184
C.190	synth – generic synthesis script . . . . .	184
C.191	synth_achronix – synthesis for Achronix Speedster2i FPGAs. . . . .	186
C.192	synth_anlogic – synthesis for Anlogic FPGAs . . . . .	187
C.193	synth_coolrunner2 – synthesis for Xilinx Coolrunner-II CPLDs . . . . .	189
C.194	synth_easic – synthesis for eASIC platform . . . . .	191
C.195	synth_ecp5 – synthesis for ECP5 FPGAs . . . . .	192
C.196	synth_efinix – synthesis for Efinix FPGAs . . . . .	195
C.197	synth_gatamate – synthesis for Cologne Chip GateMate FPGAs . . . . .	197
C.198	synth_gowin – synthesis for Gowin FPGAs . . . . .	199
C.199	synth_greenpak4 – synthesis for GreenPAK4 FPGAs . . . . .	202
C.200	synth_ice40 – synthesis for iCE40 FPGAs . . . . .	203
C.201	synth_intel – synthesis for Intel (Altera) FPGAs. . . . .	207
C.202	synth_intel_alm – synthesis for ALM-based Intel (Altera) FPGAs. . . . .	209
C.203	synth_machxo2 – synthesis for MachXO2 FPGAs. This work is experimental. . . . .	211
C.204	synth_nexus – synthesis for Lattice Nexus FPGAs . . . . .	213
C.205	synth_quicklogic – Synthesis for QuickLogic FPGAs . . . . .	216
C.206	synth_sf2 – synthesis for SmartFusion2 and IGLOO2 FPGAs . . . . .	218
C.207	synth_xilinx – synthesis for Xilinx FPGAs . . . . .	220
C.208	cl – execute a TCL script file . . . . .	223
C.209	techmap – generic technology mapper . . . . .	224
C.210	tee – redirect command output to file . . . . .	227
C.211	test_abcloop – automatically test handling of loops in abc command . . . . .	227

## CONTENTS

C.212	test_autotb – generate simple test benches	227
C.213	test_cell – automatically test the implementation of a cell type	228
C.214	test_pmggen – test pass for pmgen	229
C.215	order – print cells in topological order	229
C.216	trace – redirect command output to file	230
C.217	tribuf – infer tri-state buffers	230
C.218	uniquify – create unique copies of modules	230
C.219	verific – load Verilog and VHDL designs using Verific	231
C.220	verilog_defaults – set default options for read_verilog	234
C.221	verilog_defines – define and undefine verilog defines	235
C.222	wbflip – flip the whitebox attribute	235
C.223	wreduce – reduce the word size of operations if possible	235
C.224	write_aiger – write design to AIGER file	236
C.225	write_blif – write design to BLIF file	236
C.226	write_btor – write design to BTOR file	238
C.227	write_cxxrtl – convert design to C++ RTL simulation	238
C.228	write_edif – write design to EDIF netlist file	243
C.229	write_file – write a text to a file	243
C.230	write_firrtl – write design to a FIRRTL file	244
C.231	write_ilang – (deprecated) alias of write_rtlil	244
C.232	write_intersynth – write design to InterSynth netlist file	244
C.233	write_jny – generate design metadata	244
C.234	write_json – write design to a JSON file	245
C.235	write_rtlil – write design to RTLIL file	249
C.236	write_simplec – convert design to simple C code	249
C.237	write_smt2 – write design to SMT-LIBv2 file	250
C.238	write_smv – write design to SMV file	253
C.239	write_spice – write design to SPICE netlist file	253
C.240	write_table – write design as connectivity table	254
C.241	write_verilog – write design to Verilog file	254
C.242	write_xaiger – write design to XAIGER file	256
C.243	xilinx_dffopt – Xilinx: optimize FF control signal usage	256
C.244	xilinx_dsp – Xilinx: pack resources into DSPs	256
C.245	xilinx_srl – Xilinx shift register extraction	257
C.246	zinit – add inverters so all FF are zero-initialized	257

## CONTENTS

<b>D RTLIL Text Representation</b>	<b>258</b>
D.1 Lexical elements . . . . .	258
D.1.1 Characters . . . . .	258
D.1.2 Identifiers . . . . .	258
D.1.3 Values . . . . .	259
D.1.4 Strings . . . . .	259
D.1.5 Comments . . . . .	260
D.2 File . . . . .	260
D.2.1 Autoindex statements . . . . .	260
D.2.2 Modules . . . . .	260
D.2.3 Attribute statements . . . . .	260
D.2.4 Signal specifications . . . . .	261
D.2.5 Connections . . . . .	261
D.2.6 Wires . . . . .	261
D.2.7 Memories . . . . .	262
D.2.8 Cells . . . . .	262
D.2.9 Processes . . . . .	262
D.2.10 Switches . . . . .	263
D.2.11 Syncs . . . . .	263
<b>E Application Notes</b>	<b>264</b>

# Chapter 1

## Introduction

This document presents the Free and Open Source (FOSS) Verilog HDL synthesis tool “Yosys”. Its design and implementation as well as its performance on real-world designs is discussed in this document.

### 1.1 History of Yosys

A Hardware Description Language (HDL) is a computer language used to describe circuits. A HDL synthesis tool is a computer program that takes a formal description of a circuit written in an HDL as input and generates a netlist that implements the given circuit as output.

Currently the most widely used and supported HDLs for digital circuits are Verilog [Ver06][Ver02] and VHDL<sup>1</sup> [VHD09][VHD04]. Both HDLs are used for test and verification purposes as well as logic synthesis, resulting in a set of synthesizable and a set of non-synthesizable language features. In this document we only look at the synthesizable subset of the language features.

In recent work on heterogeneous coarse-grain reconfigurable logic [WGS<sup>+</sup>12] the need for a custom application-specific HDL synthesis tool emerged. It was soon realised that a synthesis tool that understood Verilog or VHDL would be preferred over a synthesis tool for a custom HDL. Given an existing Verilog or VHDL front end, the work for writing the necessary additional features and integrating them in an existing tool can be estimated to be about the same as writing a new tool with support for a minimalistic custom HDL.

The proposed custom HDL synthesis tool should be licensed under a Free and Open Source Software (FOSS) licence. So an existing FOSS Verilog or VHDL synthesis tool would have been needed as basis to build upon. The main advantages of choosing Verilog or VHDL is the ability to synthesize existing HDL code and to mitigate the requirement for circuit-designers to learn a new language. In order to take full advantage of any existing FOSS Verilog or VHDL tool, such a tool would have to provide a feature-complete implementation of the synthesizable HDL subset.

Basic RTL synthesis is a well understood field [HS96]. Lexing, parsing and processing of computer languages [ASU86] is a thoroughly researched field. All the information required to write such tools has been openly available for a long time, and it is therefore likely that a FOSS HDL synthesis tool with a feature-complete Verilog or VHDL front end must exist which can be used as a basis for a custom RTL synthesis tool.

Due to the author’s preference for Verilog over VHDL it was decided early on to go for Verilog instead of VHDL<sup>2</sup>. So the existing FOSS Verilog synthesis tools were evaluated (see App. ??). The results of this evaluation are utterly devastating. Therefore a completely new Verilog synthesis tool was implemented and is recommended as basis for custom synthesis tools. This is the tool that is discussed in this document.

---

<sup>1</sup>VHDL is an acronym for “VHSIC hardware description language” and VHSIC is an acronym for “Very-High-Speed Integrated Circuits”.

<sup>2</sup>A quick investigation into FOSS VHDL tools yielded similar grim results for FOSS VHDL synthesis tools.

## 1.2 Structure of this Document

The structure of this document is as follows:

Chapter 1 is this introduction.

Chapter 2 covers a short introduction to the world of HDL synthesis. Basic principles and the terminology are outlined in this chapter.

Chapter 3 gives the quickest possible outline to how the problem of implementing a HDL synthesis tool is approached in the case of Yosys.

Chapter 4 contains a more detailed overview of the implementation of Yosys. This chapter covers the data structures used in Yosys to represent a design in detail and is therefore recommended reading for everyone who is interested in understanding the Yosys internals.

Chapter 5 covers the internal cell library used by Yosys. This is especially important knowledge for anyone who wants to understand the intermediate netlists used internally by Yosys.

Chapter 6 gives a tour to the internal APIs of Yosys. This is recommended reading for everyone who actually wants to read or write Yosys source code. The chapter concludes with an example loadable module for Yosys.

Chapters 7, 8, and 9 cover three important pieces of the synthesis pipeline: The Verilog frontend, the optimization passes and the technology mapping to the target architecture, respectively.

Chapter ?? covers the evaluation of the performance (correctness and quality) of Yosys on real-world input data. The chapter concludes the main part of this document with conclusions and outlook to future work.

Various appendices, including a command reference manual (App. C) and an evaluation of pre-existing FOSS Verilog synthesis tools (App. ??) complete this document.



## Chapter 2

# Basic Principles

This chapter contains a short introduction to the basic principles of digital circuit synthesis.

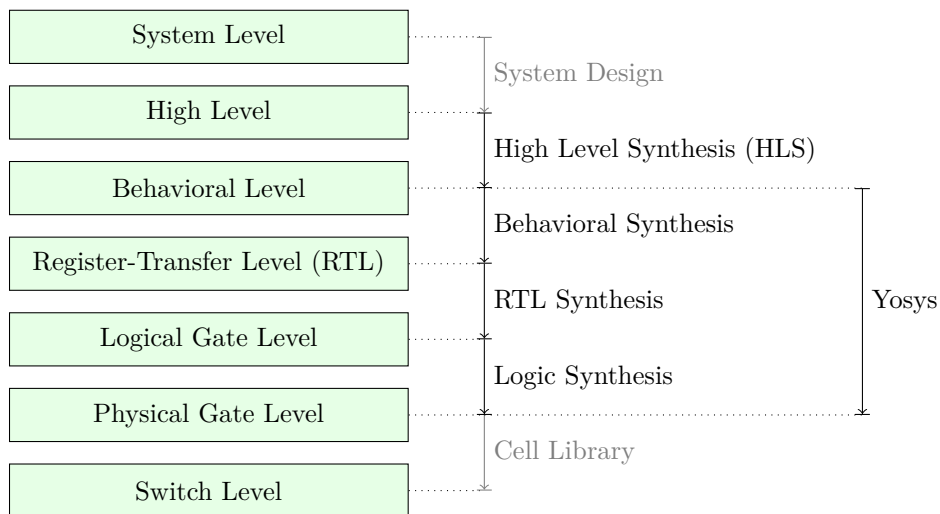
### 2.1 Levels of Abstraction

Digital circuits can be represented at different levels of abstraction. During the design process a circuit is usually first specified using a higher level abstraction. Implementation can then be understood as finding a functionally equivalent representation at a lower abstraction level. When this is done automatically using software, the term *synthesis* is used.

So synthesis is the automatic conversion of a high-level representation of a circuit to a functionally equivalent low-level representation of a circuit. Figure 2.1 lists the different levels of abstraction and how they relate to different kinds of synthesis.

Regardless of the way a lower level representation of a circuit is obtained (synthesis or manual design), the lower level representation is usually verified by comparing simulation results of the lower level and the higher level representation <sup>1</sup>. Therefore even if no synthesis is used, there must still be a simulatable representation of the circuit in all levels to allow for verification of the design.

<sup>1</sup>In recent years formal equivalence checking also became an important verification method for validating RTL and lower abstraction representation of the design.



**Figure 2.1:** Different levels of abstraction and synthesis.

Note: The exact meaning of terminology such as “High-Level” is of course not fixed over time. For example the HDL “ABEL” was first introduced in 1985 as “A High-Level Design Language for Programmable Logic Devices” [LHBB85], but would not be considered a “High-Level Language” today.

### 2.1.1 System Level

The System Level abstraction of a system only looks at its biggest building blocks like CPUs and computing cores. At this level the circuit is usually described using traditional programming languages like C/C++ or Matlab. Sometimes special software libraries are used that are aimed at simulation circuits on the system level, such as SystemC.

Usually no synthesis tools are used to automatically transform a system level representation of a circuit to a lower-level representation. But system level design tools exist that can be used to connect system level building blocks.

The IEEE 1685-2009 standard defines the IP-XACT file format that can be used to represent designs on the system level and building blocks that can be used in such system level designs. [IP-10]

### 2.1.2 High Level

The high-level abstraction of a system (sometimes referred to as *algorithmic* level) is also often represented using traditional programming languages, but with a reduced feature set. For example when representing a design at the high level abstraction in C, pointers can only be used to mimic concepts that can be found in hardware, such as memory interfaces. Full featured dynamic memory management is not allowed as it has no corresponding concept in digital circuits.

Tools exist to synthesize high level code (usually in the form of C/C++/SystemC code with additional metadata) to behavioural HDL code (usually in the form of Verilog or VHDL code). Aside from the many commercial tools for high level synthesis there are also a number of FOSS tools for high level synthesis [16] [19].

### 2.1.3 Behavioural Level

At the behavioural abstraction level a language aimed at hardware description such as Verilog or VHDL is used to describe the circuit, but so-called *behavioural modelling* is used in at least part of the circuit description. In behavioural modelling there must be a language feature that allows for imperative programming to be used to describe data paths and registers. This is the `always`-block in Verilog and the `process`-block in VHDL.

In behavioural modelling, code fragments are provided together with a *sensitivity list*; a list of signals and conditions. In simulation, the code fragment is executed whenever a signal in the sensitivity list changes its value or a condition in the sensitivity list is triggered. A synthesis tool must be able to transfer this representation into an appropriate datapath followed by the appropriate types of register.

For example consider the following Verilog code fragment:

```
1 always @(posedge clk)
2     y <= a + b;
```

In simulation the statement `y <= a + b` is executed whenever a positive edge on the signal `clk` is detected. The synthesis result however will contain an adder that calculates the sum `a + b` all the time, followed by a d-type flip-flop with the adder output on its D-input and the signal `y` on its Q-output.

Usually the imperative code fragments used in behavioural modelling can contain statements for conditional execution (**if**- and **case**-statements in Verilog) as well as loops, as long as those loops can be completely unrolled.

Interestingly there seems to be no other FOSS Tool that is capable of performing Verilog or VHDL behavioural syntheses besides Yosys (see App. ??).

#### 2.1.4 Register-Transfer Level (RTL)

On the Register-Transfer Level the design is represented by combinatorial data paths and registers (usually d-type flip flops). The following Verilog code fragment is equivalent to the previous Verilog example, but is in RTL representation:

```

1 assign tmp = a + b;           // combinatorial data path
2
3 always @(posedge clk)        // register
4     y <= tmp;
```

A design in RTL representation is usually stored using HDLs like Verilog and VHDL. But only a very limited subset of features is used, namely minimalistic `always`-blocks (Verilog) or `process`-blocks (VHDL) that model the register type used and unconditional assignments for the datapath logic. The use of HDLs on this level simplifies simulation as no additional tools are required to simulate a design in RTL representation.

Many optimizations and analyses can be performed best at the RTL level. Examples include FSM detection and optimization, identification of memories or other larger building blocks and identification of shareable resources.

Note that RTL is the first abstraction level in which the circuit is represented as a graph of circuit elements (registers and combinatorial cells) and signals. Such a graph, when encoded as list of cells and connections, is called a netlist.

RTL synthesis is easy as each circuit node element in the netlist can simply be replaced with an equivalent gate-level circuit. However, usually the term *RTL synthesis* does not only refer to synthesizing an RTL netlist to a gate level netlist but also to performing a number of highly sophisticated optimizations within the RTL representation, such as the examples listed above.

A number of FOSS tools exist that can perform isolated tasks within the domain of RTL synthesis steps. But there seems to be no FOSS tool that covers a wide range of RTL synthesis operations.

#### 2.1.5 Logical Gate Level

At the logical gate level the design is represented by a netlist that uses only cells from a small number of single-bit cells, such as basic logic gates (AND, OR, NOT, XOR, etc.) and registers (usually D-Type Flip-flops).

A number of netlist formats exists that can be used on this level, e.g. the Electronic Design Interchange Format (EDIF), but for ease of simulation often a HDL netlist is used. The latter is a HDL file (Verilog or VHDL) that only uses the most basic language constructs for instantiation and connecting of cells.

There are two challenges in logic synthesis: First finding opportunities for optimizations within the gate level netlist and second the optimal (or at least good) mapping of the logic gate netlist to an equivalent netlist of physically available gate types.

The simplest approach to logic synthesis is *two-level logic synthesis*, where a logic function is converted into a sum-of-products representation, e.g. using a Karnaugh map. This is a simple approach, but has

exponential worst-case effort and cannot make efficient use of physical gates other than AND/NAND-, OR/NOR- and NOT-Gates.

Therefore modern logic synthesis tools utilize much more complicated *multi-level logic synthesis* algorithms [BHSV90]. Most of these algorithms convert the logic function to a Binary-Decision-Diagram (BDD) or And-Inverter-Graph (AIG) and work from that representation. The former has the advantage that it has a unique normalized form. The latter has much better worst case performance and is therefore better suited for the synthesis of large logic functions.

Good FOSS tools exists for multi-level logic synthesis [27] [26] [28].

Yosys contains basic logic synthesis functionality but can also use ABC [27] for the logic synthesis step. Using ABC is recommended.

### 2.1.6 Physical Gate Level

On the physical gate level only gates are used that are physically available on the target architecture. In some cases this may only be NAND, NOR and NOT gates as well as D-Type registers. In other cases this might include cells that are more complex than the cells used at the logical gate level (e.g. complete half-adders). In the case of an FPGA-based design the physical gate level representation is a netlist of LUTs with optional output registers, as these are the basic building blocks of FPGA logic cells.

For the synthesis tool chain this abstraction is usually the lowest level. In case of an ASIC-based design the cell library might contain further information on how the physical cells map to individual switches (transistors).

### 2.1.7 Switch Level

A switch level representation of a circuit is a netlist utilizing single transistors as cells. Switch level modelling is possible in Verilog and VHDL, but is seldom used in modern designs, as in modern digital ASIC or FPGA flows the physical gates are considered the atomic build blocks of the logic circuit.

### 2.1.8 Yosys

Yosys is a Verilog HDL synthesis tool. This means that it takes a behavioural design description as input and generates an RTL, logical gate or physical gate level description of the design as output. Yosys' main strengths are behavioural and RTL synthesis. A wide range of commands (synthesis passes) exist within Yosys that can be used to perform a wide range of synthesis tasks within the domain of behavioural, rtl and logic synthesis. Yosys is designed to be extensible and therefore is a good basis for implementing custom synthesis tools for specialised tasks.

## 2.2 Features of Synthesizable Verilog

The subset of Verilog [Ver06] that is synthesizable is specified in a separate IEEE standards document, the IEEE standard 1364.1-2002 [Ver02]. This standard also describes how certain language constructs are to be interpreted in the scope of synthesis.

This section provides a quick overview of the most important features of synthesizable Verilog, structured in order of increasing complexity.

### 2.2.1 Structural Verilog

*Structural Verilog* (also known as *Verilog Netlists*) is a Netlist in Verilog syntax. Only the following language constructs are used in this case:

- Constant values
- Wire and port declarations
- Static assignments of signals to other signals
- Cell instantiations

Many tools (especially at the back end of the synthesis chain) only support structural Verilog as input. ABC is an example of such a tool. Unfortunately there is no standard specifying what *Structural Verilog* actually is, leading to some confusion about what syntax constructs are supported in structural Verilog when it comes to features such as attributes or multi-bit signals.

### 2.2.2 Expressions in Verilog

In all situations where Verilog accepts a constant value or signal name, expressions using arithmetic operations such as +, - and \*, boolean operations such as & (AND), | (OR) and ^ (XOR) and many others (comparison operations, unary operator, etc.) can also be used.

During synthesis these operators are replaced by cells that implement the respective function.

Many FOSS tools that claim to be able to process Verilog in fact only support basic structural Verilog and simple expressions. Yosys can be used to convert full featured synthesizable Verilog to this simpler subset, thus enabling such applications to be used with a richer set of Verilog features.

### 2.2.3 Behavioural Modelling

Code that utilizes the Verilog always statement is using *Behavioural Modelling*. In behavioural modelling, a circuit is described by means of imperative program code that is executed on certain events, namely any change, a rising edge, or a falling edge of a signal. This is a very flexible construct during simulation but is only synthesizable when one of the following is modelled:

- **Asynchronous or latched logic**

In this case the sensitivity list must contain all expressions that are used within the always block. The syntax @\* can be used for these cases. Examples of this kind include:

```

1 // asynchronous
2 always @* begin
3     if (add_mode)
4         y <= a + b;
5     else
6         y <= a - b;
7 end
8
9 // latched
10 always @* begin
11     if (!hold)
12         y <= a + b;
13 end

```

Note that latched logic is often considered bad style and in many cases just the result of sloppy HDL design. Therefore many synthesis tools generate warnings whenever latched logic is generated.

- **Synchronous logic (with optional synchronous reset)**

This is logic with d-type flip-flops on the output. In this case the sensitivity list must only contain the respective clock edge. Example:

```

1 // counter with synchronous reset
2 always @(posedge clk) begin
3     if (reset)
4         y <= 0;
5     else
6         y <= y + 1;
7 end

```

- **Synchronous logic with asynchronous reset**

This is logic with d-type flip-flops with asynchronous resets on the output. In this case the sensitivity list must only contain the respective clock and reset edges. The values assigned in the reset branch must be constant. Example:

```

1 // counter with asynchronous reset
2 always @(posedge clk, posedge reset) begin
3     if (reset)
4         y <= 0;
5     else
6         y <= y + 1;
7 end

```

Many synthesis tools support a wider subset of flip-flops that can be modelled using `always`-statements (including Yosys). But only the ones listed above are covered by the Verilog synthesis standard and when writing new designs one should limit herself or himself to these cases.

In behavioural modelling, blocking assignments (`=`) and non-blocking assignments (`<=`) can be used. The concept of blocking vs. non-blocking assignment is one of the most misunderstood constructs in Verilog [CI00].

The blocking assignment behaves exactly like an assignment in any imperative programming language, while with the non-blocking assignment the right hand side of the assignment is evaluated immediately but the actual update of the left hand side register is delayed until the end of the time-step. For example the Verilog code `a <= b; b <= a;` exchanges the values of the two registers. See Sec. ?? for a more detailed description of this behaviour.

## 2.2.4 Functions and Tasks

Verilog supports *Functions* and *Tasks* to bundle statements that are used in multiple places (similar to *Procedures* in imperative programming). Both constructs can be implemented easily by substituting the function/task-call with the body of the function or task.

## 2.2.5 Conditionals, Loops and Generate-Statements

Verilog supports **if-else**-statements and **for**-loops inside **always**-statements.

It also supports both features in **generate**-statements on the module level. This can be used to selectively enable or disable parts of the module based on the module parameters (**if-else**) or to generate a set of similar subcircuits (**for**).

While the **if-else**-statement inside an **always**-block is part of behavioural modelling, the three other cases are (at least for a synthesis tool) part of a built-in macro processor. Therefore it must be possible for the synthesis tool to completely unroll all loops and evaluate the condition in all **if-else**-statement in **generate**-statements using const-folding.

Examples for this can be found in Fig. ?? and Fig. ?? in App. ??.

### 2.2.6 Arrays and Memories

Verilog supports arrays. This is in general a synthesizable language feature. In most cases arrays can be synthesized by generating addressable memories. However, when complex or asynchronous access patterns are used, it is not possible to model an array as memory. In these cases the array must be modelled using individual signals for each word and all accesses to the array must be implemented using large multiplexers.

In some cases it would be possible to model an array using memories, but it is not desired. Consider the following delay circuit:

```

1  module (clk, in_data, out_data);
2
3  parameter BITS = 8;
4  parameter STAGES = 4;
5
6  input clk;
7  input [BITS-1:0] in_data;
8  output [BITS-1:0] out_data;
9  reg [BITS-1:0] ffs [STAGES-1:0];
10
11 integer i;
12 always @(posedge clk) begin
13     ffs[0] <= in_data;
14     for (i = 1; i < STAGES; i = i+1)
15         ffs[i] <= ffs[i-1];
16 end
17
18 assign out_data = ffs[STAGES-1];
19
20 endmodule

```

This could be implemented using an addressable memory with **STAGES** input and output ports. A better implementation would be to use a simple chain of flip-flops (a so-called shift register). This better implementation can either be obtained by first creating a memory-based implementation and then optimizing it based on the static address signals for all ports or directly identifying such situations in the language front end and converting all memory accesses to direct accesses to the correct signals.

## 2.3 Challenges in Digital Circuit Synthesis

This section summarizes the most important challenges in digital circuit synthesis. Tools can be characterized by how well they address these topics.

### 2.3.1 Standards Compliance

The most important challenge is compliance with the HDL standards in question (in case of Verilog the IEEE Standards 1364.1-2002 and 1364-2005). This can be broken down in two items:

- Completeness of implementation of the standard
- Correctness of implementation of the standard

Completeness is mostly important to guarantee compatibility with existing HDL code. Once a design has been verified and tested, HDL designers are very reluctant regarding changes to the design, even if it is only about a few minor changes to work around a missing feature in a new synthesis tool.

Correctness is crucial. In some areas this is obvious (such as correct synthesis of basic behavioural models). But it is also crucial for the areas that concern minor details of the standard, such as the exact rules for handling signed expressions, even when the HDL code does not target different synthesis tools. This is because (unlike software source code that is only processed by compilers), in most design flows HDL code is not only processed by the synthesis tool but also by one or more simulators and sometimes even a formal verification tool. It is key for this verification process that all these tools use the same interpretation for the HDL code.

### 2.3.2 Optimizations

Generally it is hard to give a one-dimensional description of how well a synthesis tool optimizes the design. First of all because not all optimizations are applicable to all designs and all synthesis tasks. Some optimizations work (best) on a coarse-grained level (with complex cells such as adders or multipliers) and others work (best) on a fine-grained level (single bit gates). Some optimizations target area and others target speed. Some work well on large designs while others don't scale well and can only be applied to small designs.

A good tool is capable of applying a wide range of optimizations at different levels of abstraction and gives the designer control over which optimizations are performed (or skipped) and what the optimization goals are.

### 2.3.3 Technology Mapping

Technology mapping is the process of converting the design into a netlist of cells that are available in the target architecture. In an ASIC flow this might be the process-specific cell library provided by the fab. In an FPGA flow this might be LUT cells as well as special function units such as dedicated multipliers. In a coarse-grain flow this might even be more complex special function units.

An open and vendor independent tool is especially of interest if it supports a wide range of different types of target architectures.

## 2.4 Script-Based Synthesis Flows

A digital design is usually started by implementing a high-level or system-level simulation of the desired function. This description is then manually transformed (or re-implemented) into a synthesizable lower-level description (usually at the behavioural level) and the equivalence of the two representations is verified by simulating both and comparing the simulation results.

Then the synthesizable description is transformed to lower-level representations using a series of tools and the results are again verified using simulation. This process is illustrated in Fig. 2.2.

In this example the System Level Model and the Behavioural Model are both manually written design files. After the equivalence of system level model and behavioural model has been verified, the lower level representations of the design can be generated using synthesis tools. Finally the RTL Model and the Gate-Level Model are verified and the design process is finished.





**Figure 2.2:** Typical design flow. Green boxes represent manually created models. Orange boxes represent models generated by synthesis tools.

However, in any real-world design effort there will be multiple iterations for this design process. The reason for this can be the late change of a design requirement or the fact that the analysis of a low-abstraction model (e.g. gate-level timing analysis) revealed that a design change is required in order to meet the design requirements (e.g. maximum possible clock speed).

Whenever the behavioural model or the system level model is changed their equivalence must be re-verified by re-running the simulations and comparing the results. Whenever the behavioural model is changed the synthesis must be re-run and the synthesis results must be re-verified.

In order to guarantee reproducibility it is important to be able to re-run all automatic steps in a design project with a fixed set of settings easily. Because of this, usually all programs used in a synthesis flow can be controlled using scripts. This means that all functions are available via text commands. When such a tool provides a GUI, this is complementary to, and not instead of, a command line interface.

Usually a synthesis flow in an UNIX/Linux environment would be controlled by a shell script that calls all required tools (synthesis and simulation/verification in this example) in the correct order. Each of these tools would be called with a script file containing commands for the respective tool. All settings required for the tool would be provided by these script files so that no manual interaction would be necessary. These script files are considered design sources and should be kept under version control just like the source code of the system level and the behavioural model.

## 2.5 Methods from Compiler Design

Some parts of synthesis tools involve problem domains that are traditionally known from compiler design. This section addresses some of these domains.

### 2.5.1 Lexing and Parsing

The best known concepts from compiler design are probably *lexing* and *parsing*. These are two methods that together can be used to process complex computer languages easily. [ASU86]

A *lexer* consumes single characters from the input and generates a stream of *lexical tokens* that consist of a *type* and a *value*. For example the Verilog input “**assign** foo = bar + 42;” might be translated by the lexer to the list of lexical tokens given in Tab. 2.1.

The lexer is usually generated by a lexer generator (e.g. flex [17]) from a description file that is using regular expressions to specify the text pattern that should match the individual tokens.

The lexer is also responsible for skipping ignored characters (such as whitespace outside string constants and comments in the case of Verilog) and converting the original text snippet to a token value.

Note that individual keywords use different token types (instead of a keyword type with different token values). This is because the parser usually can only use the Token-Type to make a decision on the grammatical role of a token.

## CHAPTER 2. BASIC PRINCIPLES

Token-Type	Token-Value
TOK_ASSIGN	-
TOK_IDENTIFIER	"foo"
TOK_EQ	-
TOK_IDENTIFIER	"bar"
TOK_PLUS	-
TOK_NUMBER	42
TOK_SEMICOLON	-

**Table 2.1:** Exemplary token list for the statement “**assign** foo = bar + 42;”.

The parser then transforms the list of tokens into a parse tree that closely resembles the productions from the computer languages grammar. As the lexer, the parser is also typically generated by a code generator (e.g. `bison` [18]) from a grammar description in Backus-Naur Form (BNF).

Let's consider the following BNF (in Bison syntax):

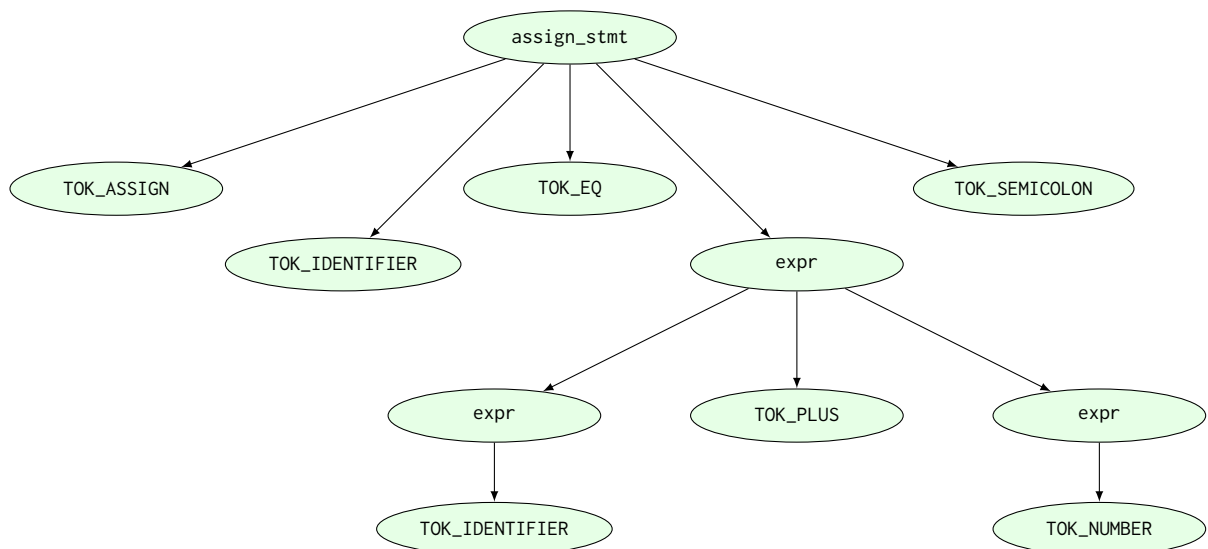
```
1 assign_stmt: TOK_ASSIGN TOK_IDENTIFIER TOK_EQ expr TOK_SEMICOLON;
2 expr: TOK_IDENTIFIER | TOK_NUMBER | expr TOK_PLUS expr;
```

The parser converts the token list to the parse tree in Fig. 2.3. Note that the parse tree never actually exists as a whole as data structure in memory. Instead the parser calls user-specified code snippets (so-called *reduce-functions*) for all inner nodes of the parse tree in depth-first order.

In some very simple applications (e.g. code generation for stack machines) it is possible to perform the task at hand directly in the reduce functions. But usually the reduce functions are only used to build an in-memory data structure with the relevant information from the parse tree. This data structure is called an *abstract syntax tree* (AST).

The exact format for the abstract syntax tree is application specific (while the format of the parse tree and token list are mostly dictated by the grammar of the language at hand). Figure 2.4 illustrates what an AST for the parse tree in Fig. 2.3 could look like.

Usually the AST is then converted into yet another representation that is more suitable for further processing. In compilers this is often an assembler-like three-address-code intermediate representation. [ASU86]



**Figure 2.3:** Example parse tree for the Verilog expression “`assign foo = bar + 42;`”.



**Figure 2.4:** Example abstract syntax tree for the Verilog expression “`assign foo = bar + 42;`”.

### 2.5.2 Multi-Pass Compilation

Complex problems are often best solved when split up into smaller problems. This is certainly true for compilers as well as for synthesis tools. The components responsible for solving the smaller problems can be connected in two different ways: through *Single-Pass Pipelining* and by using *Multiple Passes*.

Traditionally a parser and lexer are connected using the pipelined approach: The lexer provides a function that is called by the parser. This function reads data from the input until a complete lexical token has been read. Then this token is returned to the parser. So the lexer does not first generate a complete list of lexical tokens and then pass it to the parser. Instead they run concurrently and the parser can consume tokens as the lexer produces them.

The single-pass pipelining approach has the advantage of lower memory footprint (at no time must the complete design be kept in memory) but has the disadvantage of tighter coupling between the interacting components.

Therefore single-pass pipelining should only be used when the lower memory footprint is required or the components are also conceptually tightly coupled. The latter certainly is the case for a parser and its lexer. But when data is passed between two conceptually loosely coupled components it is often beneficial to use a multi-pass approach.

In the multi-pass approach the first component processes all the data and the result is stored in a in-memory data structure. Then the second component is called with this data. This reduces complexity, as only one component is running at a time. It also improves flexibility as components can be exchanged easier.

Most modern compilers are multi-pass compilers.

## Chapter 3

# Approach

Yosys is a tool for synthesising (behavioural) Verilog HDL code to target architecture netlists. Yosys aims at a wide range of application domains and thus must be flexible and easy to adapt to new tasks. This chapter covers the general approach followed in the effort to implement this tool.

### 3.1 Data- and Control-Flow

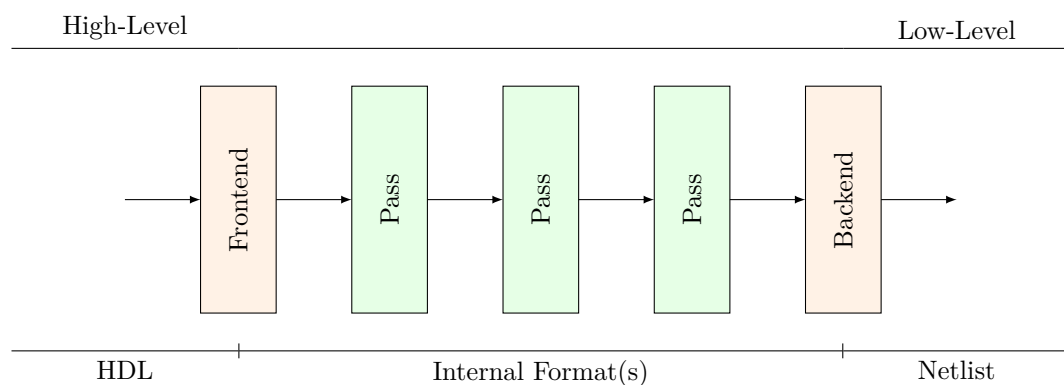
The data- and control-flow of a typical synthesis tool is very similar to the data- and control-flow of a typical compiler: different subsystems are called in a predetermined order, each consuming the data generated by the last subsystem and generating the data for the next subsystem (see Fig. 3.1).

The first subsystem to be called is usually called a *frontend*. It does not process the data generated by another subsystem but instead reads the user input—in the case of a HDL synthesis tool, the behavioural HDL code.

The subsystems that consume data from previous subsystems and produce data for the next subsystems (usually in the same or a similar format) are called *passes*.

The last subsystem that is executed transforms the data generated by the last pass into a suitable output format and writes it to a disk file. This subsystem is usually called the *backend*.

In Yosys all frontends, passes and backends are directly available as commands in the synthesis script. Thus the user can easily create a custom synthesis flow just by calling passes in the right order in a synthesis script.



**Figure 3.1:** General data- and control-flow of a synthesis tool

## 3.2 Internal Formats in Yosys

Yosys uses two different internal formats. The first is used to store an abstract syntax tree (AST) of a Verilog input file. This format is simply called *AST* and is generated by the Verilog Frontend. This data structure is consumed by a subsystem called *AST Frontend*<sup>1</sup>. This AST Frontend then generates a design in Yosys' main internal format, the Register-Transfer-Level-Intermediate-Language (RTLIL) representation. It does that by first performing a number of simplifications within the AST representation and then generating RTLIL from the simplified AST data structure.

The RTLIL representation is used by all passes as input and outputs. This has the following advantages over using different representational formats between different passes:

- The passes can be rearranged in a different order and passes can be removed or inserted.
- Passes can simply pass-thru the parts of the design they don't change without the need to convert between formats. In fact Yosys passes output the same data structure they received as input and performs all changes in place.
- All passes use the same interface, thus reducing the effort required to understand a pass when reading the Yosys source code, e.g. when adding additional features.

The RTLIL representation is basically a netlist representation with the following additional features:

- An internal cell library with fixed-function cells to represent RTL datapath and register cells as well as logical gate-level cells (single-bit gates and registers).
- Support for multi-bit values that can use individual bits from wires as well as constant bits to represent coarse-grain netlists.
- Support for basic behavioural constructs (if-then-else structures and multi-case switches with a sensitivity list for updating the outputs).
- Support for multi-port memories.

The use of RTLIL also has the disadvantage of having a very powerful format between all passes, even when doing gate-level synthesis where the more advanced features are not needed. In order to reduce complexity for passes that operate on a low-level representation, these passes check the features used in the input RTLIL and fail to run when unsupported high-level constructs are used. In such cases a pass that transforms the higher-level constructs to lower-level constructs must be called from the synthesis script first.

## 3.3 Typical Use Case

The following example script may be used in a synthesis flow to convert the behavioural Verilog code from the input file `design.v` to a gate-level netlist `synth.v` using the cell library described by the Liberty file [\[25\]](#) `cells.lib`:

```

1 # read input file to internal representation
2 read_verilog design.v
3
4 # convert high-level behavioral parts ("processes") to d-type flip-flops and muxes
5 proc

```

<sup>1</sup>In Yosys the term *pass* is only used to refer to commands that operate on the RTLIL data structure.

## CHAPTER 3. APPROACH

```
6
7 # perform some simple optimizations
8 opt
9
10 # convert high-level memory constructs to d-type flip-flops and multiplexers
11 memory
12
13 # perform some simple optimizations
14 opt
15
16 # convert design to (logical) gate-level netlists
17 techmap
18
19 # perform some simple optimizations
20 opt
21
22 # map internal register types to the ones from the cell library
23 dfflibmap -liberty cells.lib
24
25 # use ABC to map remaining logic to cells from the cell library
26 abc -liberty cells.lib
27
28 # cleanup
29 opt
30
31 # write results to output file
32 write_verilog synth.v
```

A detailed description of the commands available in Yosys can be found in App. [C](#).

## Chapter 4

# Implementation Overview

Yosys is an extensible open source hardware synthesis tool. It is aimed at designers who are looking for an easily accessible, universal, and vendor-independent synthesis tool, as well as scientists who do research in electronic design automation (EDA) and are looking for an open synthesis framework that can be used to test algorithms on complex real-world designs.

Yosys can synthesize a large subset of Verilog 2005 and has been tested with a wide range of real-world designs, including the OpenRISC 1200 CPU [23], the openMSP430 CPU [22], the OpenCores I<sup>2</sup>C master [20] and the k68 CPU [21].

As of this writing a Yosys VHDL frontend is in development.

Yosys is written in C++ (using some features from the new C++11 standard). This chapter describes some of the fundamental Yosys data structures. For the sake of simplicity the C++ type names used in the Yosys implementation are used in this chapter, even though the chapter only explains the conceptual idea behind it and can be used as reference to implement a similar system in any language.

### 4.1 Simplified Data Flow

Figure 4.1 shows the simplified data flow within Yosys. Rectangles in the figure represent program modules and ellipses internal data structures that are used to exchange design data between the program modules.

Design data is read in using one of the frontend modules. The high-level HDL frontends for Verilog and VHDL code generate an abstract syntax tree (AST) that is then passed to the AST frontend. Note that both HDL frontends use the same AST representation that is powerful enough to cover the Verilog HDL and VHDL language.

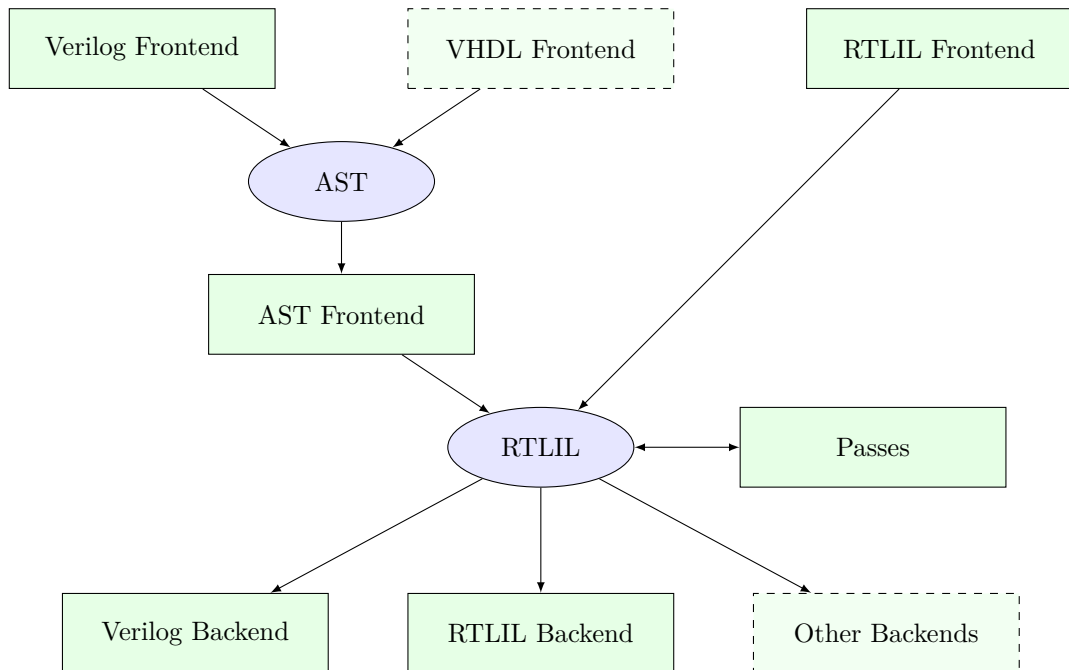
The AST Frontend then compiles the AST to Yosys's main internal data format, the RTL Intermediate Language (RTLIL). A more detailed description of this format is given in the next section.

There is also a text representation of the RTLIL data structure that can be parsed using the RTLIL Frontend.

The design data may then be transformed using a series of passes that all operate on the RTLIL representation of the design.

Finally the design in RTLIL representation is converted back to text by one of the backends, namely the Verilog Backend for generating Verilog netlists and the RTLIL Backend for writing the RTLIL data in the same format that is understood by the RTLIL Frontend.

With the exception of the AST Frontend, which is called by the high-level HDL frontends and can't be called directly by the user, all program modules are called by the user (usually using a synthesis script that contains text commands for Yosys).



**Figure 4.1:** Yosys simplified data flow (ellipses: data structures, rectangles: program modules)

By combining passes in different ways and/or adding additional passes to Yosys it is possible to adapt Yosys to a wide range of applications. For this to be possible it is key that (1) all passes operate on the same data structure (RTLIL) and (2) that this data structure is powerful enough to represent the design in different stages of the synthesis.

## 4.2 The RTL Intermediate Language

All frontends, passes and backends in Yosys operate on a design in RTLIL representation. The only exception are the high-level frontends that use the AST representation as an intermediate step before generating RTLIL data.

In order to avoid reinventing names for the RTLIL classes, they are simply referred to by their full C++ name, i.e. including the `RTLIL::` namespace prefix, in this document.

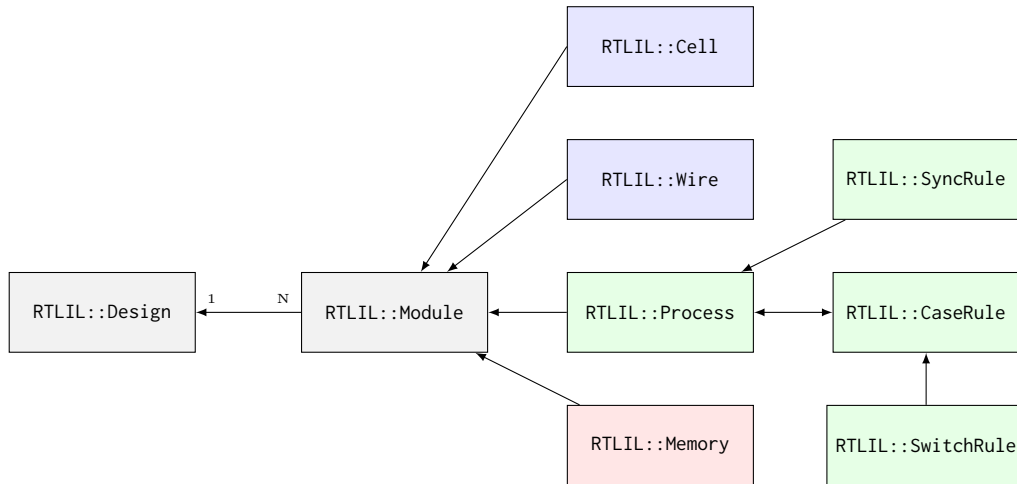
Figure 4.2 shows a simplified Entity-Relationship Diagram (ER Diagram) of RTLIL. In  $1 : N$  relationships the arrow points from the  $N$  side to the 1. For example one `RTLIL::Design` contains  $N$  (zero to many) instances of `RTLIL::Module`. A two-pointed arrow indicates a  $1 : 1$  relationship.

The `RTLIL::Design` is the root object of the RTLIL data structure. There is always one “current design” in memory which passes operate on, frontends add data to and backends convert to exportable formats. But in some cases passes internally generate additional `RTLIL::Design` objects. For example when a pass is reading an auxiliary Verilog file such as a cell library, it might create an additional `RTLIL::Design` object and call the Verilog frontend with this other object to parse the cell library.

There is only one active `RTLIL::Design` object that is used by all frontends, passes and backends called by the user, e.g. using a synthesis script. The `RTLIL::Design` then contains zero to many `RTLIL::Module` objects. This corresponds to modules in Verilog or entities in VHDL. Each module in turn contains objects from three different categories:

- `RTLIL::Cell` and `RTLIL::Wire` objects represent classical netlist data.





**Figure 4.2:** Simplified RTLIL Entity-Relationship Diagram

- RTLIL::Process objects represent the decision trees (if-then-else statements, etc.) and synchronization declarations (clock signals and sensitivity) from Verilog always and VHDL process blocks.
- RTLIL::Memory objects represent addressable memories (arrays).

Usually the output of the synthesis procedure is a netlist, i.e. all RTLIL::Process and RTLIL::Memory objects must be replaced by RTLIL::Cell and RTLIL::Wire objects by synthesis passes.

All features of the HDL that cannot be mapped directly to these RTLIL classes must be transformed to an RTLIL-compatible representation by the HDL frontend. This includes Verilog-features such as generate-blocks, loops and parameters.

The following sections contain a more detailed description of the different parts of RTLIL and rationale behind some of the design decisions.

### 4.2.1 RTLIL Identifiers

All identifiers in RTLIL (such as module names, port names, signal names, cell types, etc.) follow the following naming convention: they must either start with a backslash (\) or a dollar sign (\$).

Identifiers starting with a backslash are public visible identifiers. Usually they originate from one of the HDL input files. For example the signal name “\sig42” is most likely a signal that was declared using the name “sig42” in an HDL input file. On the other hand the signal name “\$sig42” is an auto-generated signal name. The backends convert all identifiers that start with a dollar sign to identifiers that do not collide with identifiers that start with a backslash.

This has three advantages:

- First, it is impossible that an auto-generated identifier collides with an identifier that was provided by the user.
- Second, the information about which identifiers were originally provided by the user is always available which can help guide some optimizations. For example the “opt\_rmunused” tries to preserve signals with a user-provided name but doesn’t hesitate to delete signals that have auto-generated names when they just duplicate other signals.

- Third, the delicate job of finding suitable auto-generated public visible names is deferred to one central location. Internally auto-generated names that may hold important information for Yosys developers can be used without disturbing external tools. For example the Verilog backend assigns names in the form `_integer_`.

Whitespace and control characters (any character with an ASCII code 32 or less) are not allowed in RTLIL identifiers; most frontends and backends cannot support these characters in identifiers.

In order to avoid programming errors, the RTLIL data structures check if all identifiers start with either a backslash or a dollar sign, and contain no whitespace or control characters. Violating these rules results in a runtime error.

All RTLIL identifiers are case sensitive.

Some transformations, such as flattening, may have to change identifiers provided by the user to avoid name collisions. When that happens, attribute `hdlname` is attached to the object with the changed identifier. This attribute contains one name (if emitted directly by the frontend, or is a result of disambiguation) or multiple names separated by spaces (if a result of flattening). All names specified in the `hdlname` attribute are public and do not include the leading ```.

#### 4.2.2 RTLIL::Design and RTLIL::Module

The RTLIL::Design object is basically just a container for RTLIL::Module objects. In addition to a list of RTLIL::Module objects the RTLIL::Design also keeps a list of *selected objects*, i.e. the objects that passes should operate on. In most cases the whole design is selected and therefore passes operate on the whole design. But this mechanism can be useful for more complex synthesis jobs in which only parts of the design should be affected by certain passes.

Besides the objects shown in the ER diagram in Fig. 4.2 an RTLIL::Module object contains the following additional properties:

- The module name
- A list of attributes
- A list of connections between wires
- An optional frontend callback used to derive parametrized variations of the module

The attributes can be Verilog attributes imported by the Verilog frontend or attributes assigned by passes. They can be used to store additional metadata about modules or just mark them to be used by certain part of the synthesis script but not by others.

Verilog and VHDL both support parametric modules (known as “generic entities” in VHDL). The RTLIL format does not support parametric modules itself. Instead each module contains a callback function into the AST frontend to generate a parametrized variation of the RTLIL::Module as needed. This callback then returns the auto-generated name of the parametrized variation of the module. (A hash over the parameters and the module name is used to prohibit the same parametrized variation from being generated twice. For modules with only a few parameters, a name directly containing all parameters is generated instead of a hash string.)

### 4.2.3 RTLIL::Cell and RTLIL::Wire

A module contains zero to many RTLIL::Cell and RTLIL::Wire objects. Objects of these types are used to model netlists. Usually the goal of all synthesis efforts is to convert all modules to a state where the functionality of the module is implemented only by cells from a given cell library and wires to connect these cells with each other. Note that module ports are just wires with a special property.

An RTLIL::Wire object has the following properties:

- The wire name
- A list of attributes
- A width (buses are just wires with a width > 1)
- Bus direction (MSB to LSB or vice versa)
- Lowest valid bit index (LSB or MSB depending on bus direction)
- If the wire is a port: port number and direction (input/output/inout)

As with modules, the attributes can be Verilog attributes imported by the Verilog frontend or attributes assigned by passes.

In Yosys, buses (signal vectors) are represented using a single wire object with a width > 1. So Yosys does not convert signal vectors to individual signals. This makes some aspects of RTLIL more complex but enables Yosys to be used for coarse grain synthesis where the cells of the target architecture operate on entire signal vectors instead of single bit wires.

In Verilog and VHDL, busses may have arbitrary bounds, and LSB can have either the lowest or the highest bit index. In RTLIL, bit 0 always corresponds to LSB; however, information from the HDL frontend is preserved so that the bus will be correctly indexed in error messages, backend output, constraint files, etc.

An RTLIL::Cell object has the following properties:

- The cell name and type
- A list of attributes
- A list of parameters (for parametric cells)
- Cell ports and the connections of ports to wires and constants

The connections of ports to wires are coded by assigning an RTLIL::SigSpec to each cell port. The RTLIL::SigSpec data type is described in the next section.

### 4.2.4 RTLIL::SigSpec

A “signal” is everything that can be applied to a cell port. I.e.

- Any constant value of arbitrary bit-width  
For example: 1337, 16'b0000010100111001, 1'b1, 1'bx
- All bits of a wire or a selection of bits from a wire  
For example: mywire, mywire[24], mywire[15:8]
- Concatenations of the above  
For example: {16'd1337, mywire[15:8]}

The `RTLIL::SigSpec` data type is used to represent signals. The `RTLIL::Cell` object contains one `RTLIL::SigSpec` for each cell port.

In addition, connections between wires are represented using a pair of `RTLIL::SigSpec` objects. Such pairs are needed in different locations. Therefore the type name `RTLIL::SigSig` was defined for such a pair.

#### 4.2.5 RTLIL::Process

When a high-level HDL frontend processes behavioural code it splits it up into data path logic (e.g. the expression `a + b` is replaced by the output of an adder that takes `a` and `b` as inputs) and an `RTLIL::Process` that models the control logic of the behavioural code. Let's consider a simple example:

```

1 module ff_with_en_and_async_reset(clock, reset, enable, d, q);
2   input clock, reset, enable, d;
3   output reg q;
4   always @(posedge clock, posedge reset)
5       if (reset)
6           q <= 0;
7       else if (enable)
8           q <= d;
9 endmodule

```

In this example there is no data path and therefore the `RTLIL::Module` generated by the frontend only contains a few `RTLIL::Wire` objects and an `RTLIL::Process`. The `RTLIL::Process` in `RTLIL` syntax:

```

1 process $proc$ff_with_en_and_async_reset.v:4$1
2     assign $0\q[0:0] \q
3     switch \reset
4         case 1'1
5             assign $0\q[0:0] 1'0
6         case
7             switch \enable
8                 case 1'1
9                     assign $0\q[0:0] \d
10                case
11                end
12        end
13    sync posedge \clock
14        update \q $0\q[0:0]
15    sync posedge \reset
16        update \q $0\q[0:0]
17 end

```

This `RTLIL::Process` contains two `RTLIL::SyncRule` objects, two `RTLIL::SwitchRule` objects and five `RTLIL::CaseRule` objects. The wire `$0\q[0:0]` is an automatically created wire that holds the next value of `\q`. The lines 2...12 describe how `$0\q[0:0]` should be calculated. The lines 13...16 describe how the value of `$0\q[0:0]` is used to update `\q`.

An `RTLIL::Process` is a container for zero or more `RTLIL::SyncRule` objects and exactly one `RTLIL::CaseRule` object, which is called the *root case*.

An `RTLIL::SyncRule` object contains an (optional) synchronization condition (signal and edge-type), zero or more assignments (`RTLIL::SigSig`), and zero or more memory writes (`RTLIL::MemWriteAction`). The always synchronization condition is used to break combinatorial loops when a latch should be inferred instead.

## CHAPTER 4. IMPLEMENTATION OVERVIEW

An `RTLIL::CaseRule` is a container for zero or more assignments (`RTLIL::SigSig`) and zero or more `RTLIL::SwitchRule` objects. An `RTLIL::SwitchRule` objects is a container for zero or more `RTLIL::CaseRule` objects.

In the above example the lines 2...12 are the root case. Here `$0\q[0:0]` is first assigned the old value `\q` as default value (line 2). The root case also contains an `RTLIL::SwitchRule` object (lines 3...12). Such an object is very similar to the C `switch` statement as it uses a control signal (`\reset` in this case) to determine which of its cases should be active. The `RTLIL::SwitchRule` object then contains one `RTLIL::CaseRule` object per case. In this example there is a case<sup>1</sup> for `\reset == 1` that causes `$0\q[0:0]` to be set (lines 4 and 5) and a default case that in turn contains a switch that sets `$0\q[0:0]` to the value of `\d` if `\enable` is active (lines 6...11).

A case can specify zero or more compare values that will determine whether it matches. Each of the compare values must be the exact same width as the control signal. When more than one compare value is specified, the case matches if any of them matches the control signal; when zero compare values are specified, the case always matches (i.e. it is the default case).

A switch prioritizes cases from first to last: multiple cases can match, but only the first matched case becomes active. This normally synthesizes to a priority encoder. The `parallel_case` attribute allows passes to assume that no more than one case will match, and `full_case` attribute allows passes to assume that exactly one case will match; if these invariants are ever dynamically violated, the behavior is undefined. These attributes are useful when an invariant invisible to the synthesizer causes the control signal to never take certain bit patterns.

The lines 13...16 then cause `\q` to be updated whenever there is a positive clock edge on `\clock` or `\reset`.

In order to generate such a representation, the language frontend must be able to handle blocking and nonblocking assignments correctly. However, the language frontend does not need to identify the correct type of storage element for the output signal or generate multiplexers for the decision tree. This is done by passes that work on the RTLIL representation. Therefore it is relatively easy to substitute these steps with other algorithms that target different target architectures or perform optimizations or other transformations on the decision trees before further processing them.

One of the first actions performed on a design in RTLIL representation in most synthesis scripts is identifying asynchronous resets. This is usually done using the `proc_arst` pass. This pass transforms the above example to the following `RTLIL::Process`:

```

1  process $proc$ff_with_en_and_async_reset.v:4$1
2      assign $0\q[0:0] \q
3      switch \enable
4          case 1'1
5              assign $0\q[0:0] \d
6          case
7      end
8      sync posedge \clock
9          update \q $0\q[0:0]
10     sync high \reset
11         update \q 1'0
12 end

```

This pass has transformed the outer `RTLIL::SwitchRule` into a modified `RTLIL::SyncRule` object for the `\reset` signal. Further processing converts the `RTLIL::Process` into e.g. a d-type flip-flop with asynchronous reset and a multiplexer for the enable signal:

<sup>1</sup>The syntax `1'1` in the RTLIL code specifies a constant with a length of one bit (the first “1”), and this bit is a one (the second “1”).

```

1  cell $adff $procdff$6
2      parameter \ARST_POLARITY 1'1
3      parameter \ARST_VALUE 1'0
4      parameter \CLK_POLARITY 1'1
5      parameter \WIDTH 1
6      connect \ARST \reset
7      connect \CLK \clock
8      connect \D $0\q[0:0]
9      connect \Q \q
10 end
11 cell $mux $procmux$3
12     parameter \WIDTH 1
13     connect \A \q
14     connect \B \d
15     connect \S \enable
16     connect \Y $0\q[0:0]
17 end

```

Different combinations of passes may yield different results. Note that `$adff` and `$mux` are internal cell types that still need to be mapped to cell types from the target cell library.

Some passes refuse to operate on modules that still contain `RTLIL::Process` objects as the presence of these objects in a module increases the complexity. Therefore the passes to translate processes to a netlist of cells are usually called early in a synthesis script. The `proc` pass calls a series of other passes that together perform this conversion in a way that is suitable for most synthesis tasks.

#### 4.2.6 RTLIL::Memory

For every array (memory) in the HDL code an `RTLIL::Memory` object is created. A memory object has the following properties:

- The memory name
- A list of attributes
- The width of an addressable word
- The size of the memory in number of words

All read accesses to the memory are transformed to `$memrd` cells and all write accesses to `$memwr` cells by the language frontend. These cells consist of independent read- and write-ports to the memory. Memory initialization is transformed to `$meminit` cells by the language frontend. The `\MEMID` parameter on these cells is used to link them together and to the `RTLIL::Memory` object they belong to.

The rationale behind using separate cells for the individual ports versus creating a large multiport memory cell right in the language frontend is that the separate `$memrd` and `$memwr` cells can be consolidated using resource sharing. As resource sharing is a non-trivial optimization problem where different synthesis tasks can have different requirements it lends itself to do the optimisation in separate passes and merge the `RTLIL::Memory` objects and `$memrd` and `$memwr` cells to multiport memory blocks after resource sharing is completed.

The `memory` pass performs this conversion and can (depending on the options passed to it) transform the memories directly to d-type flip-flops and address logic or yield multiport memory blocks (represented using `$mem` cells).

See Sec. 5.1.5 for details about the memory cell types.

### 4.3 Command Interface and Synthesis Scripts

Yosys reads and processes commands from synthesis scripts, command line arguments and an interactive command prompt. Yosys commands consist of a command name and an optional whitespace separated list of arguments. Commands are terminated using the newline character or a semicolon (;). Empty lines and lines starting with the hash sign (#) are ignored. See Sec. 3.3 for an example synthesis script.

The command help can be used to access the command reference manual.

Most commands can operate not only on the entire design but also specifically on *selected* parts of the design. For example the command dump will print all selected objects in the current design while dump foobar will only print the module foobar and dump \* will print the entire design regardless of the current selection.

The selection mechanism is very powerful. For example the command dump \*/t:\$add %x:+[A] \*/w:\* %i will print all wires that are connected to the \A port of a \$add cell. Detailed documentation of the select framework can be found in the command reference for the select command.

### 4.4 Source Tree and Build System

The Yosys source tree is organized into the following top-level directories:

- backends/  
This directory contains a subdirectory for each of the backend modules.
- frontends/  
This directory contains a subdirectory for each of the frontend modules.
- kernel/  
This directory contains all the core functionality of Yosys. This includes the functions and definitions for working with the RTLIL data structures (rtlil.h and rtlil.cc), the main() function (driver.cc), the internal framework for generating log messages (log.h and log.cc), the internal framework for registering and calling passes (register.h and register.cc), some core commands that are not really passes (select.cc, show.cc, ...) and a couple of other small utility libraries.
- passes/  
This directory contains a subdirectory for each pass or group of passes. For example as of this writing the directory passes/opt/ contains the code for seven passes: opt, opt\_expr, opt\_muxtree, opt\_reduce, opt\_rmdff, opt\_rmunused and opt\_merge.
- techlibs/  
This directory contains simulation models and standard implementations for the cells from the internal cell library.
- tests/  
This directory contains a couple of test cases. Most of the smaller tests are executed automatically when make test is called. The larger tests must be executed manually. Most of the larger tests require downloading external HDL source code and/or external tools. The tests range from comparing simulation results of the synthesized design to the original sources to logic equivalence checking of entire CPU cores.

The top-level Makefile includes frontends/\*/Makefile.inc, passes/\*/Makefile.inc and backends/\*/Makefile.inc. So when extending Yosys it is enough to create a new directory in frontends/, passes/ or backends/ with your sources and a Makefile.inc. The Yosys kernel automatically detects all commands linked with Yosys. So it is not needed to add additional commands to a central list of commands.

## CHAPTER 4. IMPLEMENTATION OVERVIEW

Good starting points for reading example source code to learn how to write passes are `passes/opt/opt_rmdff.cc` and `passes/opt/opt_merge.cc`.

See the top-level README file for a quick *Getting Started* guide and build instructions. The Yosys build is based solely on Makefiles.

Users of the Qt Creator IDE can generate a QT Creator project file using `make qtcreator`. Users of the Eclipse IDE can use the “Makefile Project with Existing Code” project type in the Eclipse “New Project” dialog (only available after the CDT plugin has been installed) to create an Eclipse project in order to programming extensions to Yosys or just browse the Yosys code base.



# Chapter 5

## Internal Cell Library

Most of the passes in Yosys operate on netlists, i.e. they only care about the `RTLIL::Wire` and `RTLIL::Cell` objects in an `RTLIL::Module`. This chapter discusses the cell types used by Yosys to represent a behavioural design internally.

This chapter is split in two parts. In the first part the internal RTL cells are covered. These cells are used to represent the design on a coarse grain level. Like in the original HDL code on this level the cells operate on vectors of signals and complex cells like adders exist. In the second part the internal gate cells are covered. These cells are used to represent the design on a fine-grain gate-level. All cells from this category operate on single bit signals.

### 5.1 RTL Cells

Most of the RTL cells closely resemble the operators available in HDLs such as Verilog or VHDL. Therefore Verilog operators are used in the following sections to define the behaviour of the RTL cells.

Note that all RTL cells have parameters indicating the size of inputs and outputs. When passes modify RTL cells they must always keep the values of these parameters in sync with the size of the signals connected to the inputs and outputs.

Simulation models for the RTL cells can be found in the file `techlibs/common/simlib.v` in the Yosys source tree.

#### 5.1.1 Unary Operators

All unary RTL cells have one input port `\A` and one output port `\Y`. They also have the following parameters:

- `\A_SIGNED`  
Set to a non-zero value if the input `\A` is signed and therefore should be sign-extended when needed.
- `\A_WIDTH`  
The width of the input port `\A`.
- `\Y_WIDTH`  
The width of the output port `\Y`.

Verilog	Cell Type
$Y = \sim A$	\$not
$Y = +A$	\$pos
$Y = -A$	\$neg
$Y = \&A$	\$reduce_and
$Y =  A$	\$reduce_or
$Y = ^A$	\$reduce_xor
$Y = \sim^A$	\$reduce_xnor
$Y =  A$	\$reduce_bool
$Y = !A$	\$logic_not

**Table 5.1:** Cell types for unary operators with their corresponding Verilog expressions.

Table 5.1 lists all cells for unary RTL operators.

For the unary cells that output a logical value (\$reduce\_and, \$reduce\_or, \$reduce\_xor, \$reduce\_xnor, \$reduce\_bool, \$logic\_not), when the `\Y_WIDTH` parameter is greater than 1, the output is zero-extended, and only the least significant bit varies.

Note that \$reduce\_or and \$reduce\_bool actually represent the same logic function. But the HDL frontends generate them in different situations. A \$reduce\_or cell is generated when the prefix `|` operator is being used. A \$reduce\_bool cell is generated when a bit vector is used as a condition in an if-statement or `?:-`expression.

### 5.1.2 Binary Operators

All binary RTL cells have two input ports `\A` and `\B` and one output port `\Y`. They also have the following parameters:

- `\A_SIGNED`  
Set to a non-zero value if the input `\A` is signed and therefore should be sign-extended when needed.
- `\A_WIDTH`  
The width of the input port `\A`.
- `\B_SIGNED`  
Set to a non-zero value if the input `\B` is signed and therefore should be sign-extended when needed.
- `\B_WIDTH`  
The width of the input port `\B`.
- `\Y_WIDTH`  
The width of the output port `\Y`.

Table 5.2 lists all cells for binary RTL operators.

The \$shl and \$shr cells implement logical shifts, whereas the \$ssh1 and \$sshr cells implement arithmetic shifts. The \$shl and \$ssh1 cells implement the same operation. All four of these cells interpret the second operand as unsigned, and require `\B_SIGNED` to be zero.

Two additional shift operator cells are available that do not directly correspond to any operator in Verilog, \$shift and \$shiftx. The \$shift cell performs a right logical shift if the second operand is positive (or unsigned), and a left logical shift if it is negative. The \$shiftx cell performs the same operation as the \$shift cell, but the vacated bit positions are filled with undef (x) bits, and corresponds to the Verilog indexed part-select expression.

Verilog	Cell Type	Verilog	Cell Type
$Y = A \& B$	<code>\$and</code>	$Y = A < B$	<code>\$lt</code>
$Y = A   B$	<code>\$or</code>	$Y = A <= B$	<code>\$le</code>
$Y = A \wedge B$	<code>\$xor</code>	$Y = A == B$	<code>\$eq</code>
$Y = A \sim \wedge B$	<code>\$xnor</code>	$Y = A != B$	<code>\$ne</code>
$Y = A << B$	<code>\$shl</code>	$Y = A >= B$	<code>\$ge</code>
$Y = A >> B$	<code>\$shr</code>	$Y = A > B$	<code>\$gt</code>
$Y = A <<< B$	<code>\$sshl</code>	$Y = A + B$	<code>\$add</code>
$Y = A >>> B$	<code>\$sshr</code>	$Y = A - B$	<code>\$sub</code>
$Y = A \&\& B$	<code>\$logic_and</code>	$Y = A * B$	<code>\$mul</code>
$Y = A    B$	<code>\$logic_or</code>	$Y = A / B$	<code>\$div</code>
$Y = A === B$	<code>\$eqx</code>	$Y = A \% B$	<code>\$mod</code>
$Y = A !== B$	<code>\$nex</code>	[N/A]	<code>\$divfloor</code>
		[N/A]	<code>\$modfloor</code>
		$Y = A ** B$	<code>\$pow</code>

**Table 5.2:** Cell types for binary operators with their corresponding Verilog expressions.

For the binary cells that output a logical value (`$logic_and`, `$logic_or`, `$eqx`, `$nex`, `$lt`, `$le`, `$eq`, `$ne`, `$ge`, `$gt`), when the `\Y_WIDTH` parameter is greater than 1, the output is zero-extended, and only the least significant bit varies.

Division and modulo cells are available in two rounding modes. The original `$div` and `$mod` cells are based on truncating division, and correspond to the semantics of the verilog `/` and `%` operators. The `$divfloor` and `$modfloor` cells represent flooring division and flooring modulo, the latter of which is also known as “remainder” in several languages. See table 5.3 for a side-by-side comparison between the different semantics.

Division	Result	Truncating		Flooring	
		<code>\$div</code>	<code>\$mod</code>	<code>\$divfloor</code>	<code>\$modfloor</code>
$-10 / 3$	-3.3	-3	-1	-4	2
$10 / -3$	-3.3	-3	1	-4	-2
$-10 / -3$	3.3	3	-1	3	-1
$10 / 3$	3.3	3	1	3	1

**Table 5.3:** Comparison between different rounding modes for division and modulo cells.

### 5.1.3 Multiplexers

Multiplexers are generated by the Verilog HDL frontend for `?:-`expressions. Multiplexers are also generated by the `proc` pass to map the decision trees from `RTLIL::Process` objects to logic.

The simplest multiplexer cell type is `$mux`. Cells of this type have a `\WIDTH` parameter and data inputs `\A` and `\B` and a data output `\Y`, all of the specified width. This cell also has a single bit control input `\S`. If `\S` is 0 the value from the `\A` input is sent to the output, if it is 1 the value from the `\B` input is sent to the output. So the `$mux` cell implements the function  $Y = S ? B : A$ .

The `$pmux` cell is used to multiplex between many inputs using a one-hot select signal. Cells of this type have a `\WIDTH` and a `\S_WIDTH` parameter and inputs `\A`, `\B`, and `\S` and an output `\Y`. The `\S` input is `\S_WIDTH` bits wide. The `\A` input and the output are both `\WIDTH` bits wide and the `\B` input is `\WIDTH*\S_WIDTH` bits wide. When all bits of `\S` are zero, the value from `\A` input is sent to the output. If the  $n$ ’th bit from `\S` is set, the value  $n$ ’th `\WIDTH` bits wide slice of the `\B` input is sent to the output. When more than one bit from `\S` is set the output is undefined. Cells of this type are used to model “parallel cases” (defined by using the `parallel_case` attribute or detected by an optimization).

The `$tribuf` cell is used to implement tristate logic. Cells of this type have a `\WIDTH` parameter and inputs `\A` and `\EN` and an output `\Y`. The `\A` input and `\Y` output are `\WIDTH` bits wide, and the `\EN` input is one bit wide. When `\EN` is 0, the output `\Y` is not driven. When `\EN` is 1, the value from `\A` input is sent to the `\Y` output. Therefore, the `$tribuf` cell implements the function  $Y = EN ? A : 'bz$ .

Behavioural code with cascaded if-then-else- and case-statements usually results in trees of multiplexer cells. Many passes (from various optimizations to FSM extraction) heavily depend on these multiplexer trees to understand dependencies between signals. Therefore optimizations should not break these multiplexer trees (e.g. by replacing a multiplexer between a calculated signal and a constant zero with an `$and` gate).

### 5.1.4 Registers

SR-type latches are represented by `$sr` cells. These cells have input ports `\SET` and `\CLR` and an output port `\Q`. They have the following parameters:

- `\WIDTH`  
The width of inputs `\SET` and `\CLR` and output `\Q`.
- `\SET_POLARITY`  
The set input bits are active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.
- `\CLR_POLARITY`  
The reset input bits are active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.

Both set and reset inputs have separate bits for every output bit. When both the set and reset inputs of an `$sr` cell are active for a given bit index, the reset input takes precedence.

D-type flip-flops are represented by `$dff` cells. These cells have a clock port `\CLK`, an input port `\D` and an output port `\Q`. The following parameters are available for `$dff` cells:

- `\WIDTH`  
The width of input `\D` and output `\Q`.
- `\CLK_POLARITY`  
Clock is active on the positive edge if this parameter has the value `1'b1` and on the negative edge if this parameter is `1'b0`.

D-type flip-flops with asynchronous reset are represented by `$adff` cells. As the `$dff` cells they have `\CLK`, `\D` and `\Q` ports. In addition they also have a single-bit `\ARST` input port for the reset pin and the following additional two parameters:

- `\ARST_POLARITY`  
The asynchronous reset is active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.
- `\ARST_VALUE`  
The state of `\Q` will be set to this value when the reset is active.

Usually these cells are generated by the proc pass using the information in the designs `RTLIL::Process` objects.

D-type flip-flops with synchronous reset are represented by `$sdff` cells. As the `$dff` cells they have `\CLK`, `\D` and `\Q` ports. In addition they also have a single-bit `\SRST` input port for the reset pin and the following additional two parameters:

## CHAPTER 5. INTERNAL CELL LIBRARY

- `\SRST_POLARITY`

The synchronous reset is active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.

- `\SRST_VALUE`

The state of `\Q` will be set to this value when the reset is active.

Note that the `$dff` and `$sdff` cells can only be used when the reset value is constant.

D-type flip-flops with asynchronous load are represented by `$aldff` cells. As the `$dff` cells they have `\CLK`, `\D` and `\Q` ports. In addition they also have a single-bit `\ALOAD` input port for the async load enable pin, a `\AD` input port with the same width as data for the async load data, and the following additional parameter:

- `\ALOAD_POLARITY`

The asynchronous load is active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.

D-type flip-flops with asynchronous set and reset are represented by `$dffsr` cells. As the `$dff` cells they have `\CLK`, `\D` and `\Q` ports. In addition they also have multi-bit `\SET` and `\CLR` input ports and the corresponding polarity parameters, like `$sr` cells.

D-type flip-flops with enable are represented by `$dfffe`, `$adfffe`, `$aldfffe`, `$dffsre`, `$sdfffe`, and `$sdfffce` cells, which are enhanced variants of `$dff`, `$adff`, `$aldff`, `$dffsr`, `$sdff` (with reset over enable) and `$sdff` (with enable over reset) cells, respectively. They have the same ports and parameters as their base cell. In addition they also have a single-bit `\EN` input port for the enable pin and the following parameter:

- `\EN_POLARITY`

The enable input is active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.

D-type latches are represented by `$dlatch` cells. These cells have an enable port `\EN`, an input port `\D`, and an output port `\Q`. The following parameters are available for `$dlatch` cells:

- `\WIDTH`

The width of input `\D` and output `\Q`.

- `\EN_POLARITY`

The enable input is active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.

The latch is transparent when the `\EN` input is active.

D-type latches with reset are represented by `$adlatch` cells. In addition to `$dlatch` ports and parameters, they also have a single-bit `\ARST` input port for the reset pin and the following additional parameters:

- `\ARST_POLARITY`

The asynchronous reset is active-high if this parameter has the value `1'b1` and active-low if this parameter is `1'b0`.

- `\ARST_VALUE`

The state of `\Q` will be set to this value when the reset is active.

D-type latches with set and reset are represented by `$dlatchsr` cells. In addition to `$dlatch` ports and parameters, they also have multi-bit `\SET` and `\CLR` input ports and the corresponding polarity parameters, like `$sr` cells.

### 5.1.5 Memories

Memories are either represented using RTLIL::Memory objects, \$memrd\_v2, \$memwr\_v2, and \$meminit\_v2 cells, or by \$mem\_v2 cells alone.

In the first alternative the RTLIL::Memory objects hold the general metadata for the memory (bit width, size in number of words, etc.) and for each port a \$memrd\_v2 (read port) or \$memwr\_v2 (write port) cell is created. Having individual cells for read and write ports has the advantage that they can be consolidated using resource sharing passes. In some cases this drastically reduces the number of required ports on the memory cell. In this alternative, memory initialization data is represented by \$meminit\_v2 cells, which allow delaying constant folding for initialization addresses and data until after the frontend finishes.

The \$memrd\_v2 cells have a clock input \CLK, an enable input \EN, an address input \ADDR, a data output \DATA, an asynchronous reset input \ARST, and a synchronous reset input \SRST. They also have the following parameters:

- \MEMID  
The name of the RTLIL::Memory object that is associated with this read port.
- \ABITS  
The number of address bits (width of the \ADDR input port).
- \WIDTH  
The number of data bits (width of the \DATA output port). Note that this may be a power-of-two multiple of the underlying memory's width – such ports are called wide ports and access an aligned group of cells at once. In this case, the corresponding low bits of \ADDR must be tied to 0.
- \CLK\_ENABLE  
When this parameter is non-zero, the clock is used. Otherwise this read port is asynchronous and the \CLK input is not used.
- \CLK\_POLARITY  
Clock is active on the positive edge if this parameter has the value 1'b1 and on the negative edge if this parameter is 1'b0.
- \TRANSPARENCY\_MASK  
This parameter is a bitmask of write ports that this read port is transparent with. The bits of this parameter are indexed by the write port's \PORTID parameter. Transparency can only be enabled between synchronous ports sharing a clock domain. When transparency is enabled for a given port pair, a read and write to the same address in the same cycle will return the new value. Otherwise the old value is returned.
- \COLLISION\_X\_MASK  
This parameter is a bitmask of write ports that have undefined collision behavior with this port. The bits of this parameter are indexed by the write port's \PORTID parameter. This behavior can only be enabled between synchronous ports sharing a clock domain. When undefined collision is enabled for a given port pair, a read and write to the same address in the same cycle will return the undefined (all-X) value. This option is exclusive (for a given port pair) with the transparency option.
- \ARST\_VALUE  
Whenever the \ARST input is asserted, the data output will be reset to this value. Only used for synchronous ports.
- \SRST\_VALUE  
Whenever the \SRST input is synchronously asserted, the data output will be reset to this value. Only used for synchronous ports.
- \INIT\_VALUE  
The initial value of the data output, for synchronous ports.

## CHAPTER 5. INTERNAL CELL LIBRARY

- `\CE_OVER_SRST`

If this parameter is non-zero, the `\SRST` input is only recognized when `\EN` is true. Otherwise, `\SRST` is recognized regardless of `\EN`.

The `$memwr_v2` cells have a clock input `\CLK`, an enable input `\EN` (one enable bit for each data bit), an address input `\ADDR` and a data input `\DATA`. They also have the following parameters:

- `\MEMID`

The name of the RTLIL::Memory object that is associated with this write port.

- `\ABITS`

The number of address bits (width of the `\ADDR` input port).

- `\WIDTH`

The number of data bits (width of the `\DATA` output port). Like with `$memrd_v2` cells, the width is allowed to be any power-of-two multiple of memory width, with the corresponding restriction on address.

- `\CLK_ENABLE`

When this parameter is non-zero, the clock is used. Otherwise this write port is asynchronous and the `\CLK` input is not used.

- `\CLK_POLARITY`

Clock is active on positive edge if this parameter has the value `1'b1` and on the negative edge if this parameter is `1'b0`.

- `\PORTID`

An identifier for this write port, used to index write port bit mask parameters.

- `\PRIORITY_MASK`

This parameter is a bitmask of write ports that this write port has priority over in case of writing to the same address. The bits of this parameter are indexed by the other write port's `\PORTID` parameter. Write ports can only have priority over write ports with lower port ID. When two ports write to the same address and neither has priority over the other, the result is undefined. Priority can only be set between two synchronous ports sharing the same clock domain.

The `$meminit_v2` cells have an address input `\ADDR`, a data input `\DATA`, with the width of the `\DATA` port equal to `\WIDTH` parameter times `\WORDS` parameter, and a bit enable mask input `\EN` with width equal to `\WIDTH` parameter. All three of the inputs must resolve to a constant for synthesis to succeed.

- `\MEMID`

The name of the RTLIL::Memory object that is associated with this initialization cell.

- `\ABITS`

The number of address bits (width of the `\ADDR` input port).

- `\WIDTH`

The number of data bits per memory location.

- `\WORDS`

The number of consecutive memory locations initialized by this cell.

- `\PRIORITY`

The cell with the higher integer value in this parameter wins an initialization conflict.

## CHAPTER 5. INTERNAL CELL LIBRARY

The HDL frontend models a memory using RTLIL::Memory objects and asynchronous `$memrd_v2` and `$memwr_v2` cells. The memory pass (i.e. its various sub-passes) migrates `$dff` cells into the `$memrd_v2` and `$memwr_v2` cells making them synchronous, then converts them to a single `$mem_v2` cell and (optionally) maps this cell type to `$dff` cells for the individual words and multiplexer-based address decoders for the read and write interfaces. When the last step is disabled or not possible, a `$mem_v2` cell is left in the design.

The `$mem_v2` cell provides the following parameters:

- `\MEMID`  
The name of the original RTLIL::Memory object that became this `$mem_v2` cell.
- `\SIZE`  
The number of words in the memory.
- `\ABITS`  
The number of address bits.
- `\WIDTH`  
The number of data bits per word.
- `\INIT`  
The initial memory contents.
- `\RD_PORTS`  
The number of read ports on this memory cell.
- `\RD_WIDE_CONTINUATION`  
This parameter is `\RD_PORTS` bits wide, containing a bitmask of “wide continuation” read ports. Such ports are used to represent the extra data bits of wide ports in the combined cell, and must have all control signals identical with the preceding port, except for address, which must have the proper sub-cell address encoded in the low bits.
- `\RD_CLK_ENABLE`  
This parameter is `\RD_PORTS` bits wide, containing a clock enable bit for each read port.
- `\RD_CLK_POLARITY`  
This parameter is `\RD_PORTS` bits wide, containing a clock polarity bit for each read port.
- `\RD_TRANSPARENCY_MASK`  
This parameter is `\RD_PORTS*WR_PORTS` bits wide, containing a concatenation of all `\TRANSPARENCY_MASK` values of the original `$memrd_v2` cells.
- `\RD_COLLISION_X_MASK`  
This parameter is `\RD_PORTS*WR_PORTS` bits wide, containing a concatenation of all `\COLLISION_X_MASK` values of the original `$memrd_v2` cells.
- `\RD_CE_OVER_SRST`  
This parameter is `\RD_PORTS` bits wide, determining relative synchronous reset and enable priority for each read port.
- `\RD_INIT_VALUE`  
This parameter is `\RD_PORTS*WIDTH` bits wide, containing the initial value for each synchronous read port.
- `\RD_ARST_VALUE`  
This parameter is `\RD_PORTS*WIDTH` bits wide, containing the asynchronous reset value for each synchronous read port.



## CHAPTER 5. INTERNAL CELL LIBRARY

- `\RD_SRST_VALUE`  
This parameter is `\RD_PORTS*WIDTH` bits wide, containing the synchronous reset value for each synchronous read port.
- `\WR_PORTS`  
The number of write ports on this memory cell.
- `\WR_WIDE_CONTINUATION`  
This parameter is `\WR_PORTS` bits wide, containing a bitmask of “wide continuation” write ports.
- `\WR_CLK_ENABLE`  
This parameter is `\WR_PORTS` bits wide, containing a clock enable bit for each write port.
- `\WR_CLK_POLARITY`  
This parameter is `\WR_PORTS` bits wide, containing a clock polarity bit for each write port.
- `\WR_PRIORITY_MASK`  
This parameter is `\WR_PORTS*WR_PORTS` bits wide, containing a concatenation of all `\VPRIORITY_MASK` values of the original `$memwr_v2` cells.

The `$mem_v2` cell has the following ports:

- `\RD_CLK`  
This input is `\RD_PORTS` bits wide, containing all clock signals for the read ports.
- `\RD_EN`  
This input is `\RD_PORTS` bits wide, containing all enable signals for the read ports.
- `\RD_ADDR`  
This input is `\RD_PORTS*\ABITS` bits wide, containing all address signals for the read ports.
- `\RD_DATA`  
This input is `\RD_PORTS*WIDTH` bits wide, containing all data signals for the read ports.
- `\RD_ARST`  
This input is `\RD_PORTS` bits wide, containing all asynchronous reset signals for the read ports.
- `\RD_SRST`  
This input is `\RD_PORTS` bits wide, containing all synchronous reset signals for the read ports.
- `\WR_CLK`  
This input is `\WR_PORTS` bits wide, containing all clock signals for the write ports.
- `\WR_EN`  
This input is `\WR_PORTS*WIDTH` bits wide, containing all enable signals for the write ports.
- `\WR_ADDR`  
This input is `\WR_PORTS*\ABITS` bits wide, containing all address signals for the write ports.
- `\WR_DATA`  
This input is `\WR_PORTS*WIDTH` bits wide, containing all data signals for the write ports.

The `memory_collect` pass can be used to convert discrete `$memrd_v2`, `$memwr_v2`, and `$meminit_v2` cells belonging to the same memory to a single `$mem_v2` cell, whereas the `memory_unpack` pass performs the inverse operation. The `memory_dff` pass can combine asynchronous memory ports that are fed by or feeding registers into synchronous memory ports. The `memory_bram` pass can be used to recognize `$mem_v2` cells that can be implemented with a block RAM resource on an FPGA. The `memory_map` pass can be used to implement `$mem_v2` cells as basic logic: word-wide DFFs and address decoders.

### 5.1.6 Finite State Machines

#### FIXME:

Add a brief description of the `$fsm` cell type.

### 5.1.7 Specify rules

#### FIXME:

Add information about `$specify2`, `$specify3`, and `$specrule` cells.

### 5.1.8 Formal verification cells

#### FIXME:

Add information about `$assert`, `$assume`, `$live`, `$fair`, `$cover`, `$equiv`, `$initstate`, `$anyconst`, `$anyseq`, `$allconst`, `$allseq` cells.

#### FIXME:

Add information about `$ff` and `$_FF_` cells.

## 5.2 Gates

For gate level logic networks, fixed function single bit cells are used that do not provide any parameters.

Simulation models for these cells can be found in the file `techlibs/common/simcells.v` in the Yosys source tree.

Tables 5.4, 5.6, 5.5, 5.7, 5.8, 5.9, 5.10, 5.11 and 5.12 list all cell types used for gate level logic. The cell types `$_BUF_`, `$_NOT_`, `$_AND_`, `$_NAND_`, `$_ANDNOT_`, `$_OR_`, `$_NOR_`, `$_ORNOT_`, `$_XOR_`, `$_XNOR_`, `$_AOI3_`, `$_OAI3_`, `$_AOI4_`, `$_OAI4_`, `$_MUX_`, `$_MUX4_`, `$_MUX8_`, `$_MUX16_` and `$_NMUX_` are used to model combinatorial logic. The cell type `$_TBUF_` is used to model tristate logic.

The `$_MUX4_`, `$_MUX8_` and `$_MUX16_` cells are used to model wide muxes, and correspond to the following Verilog code:

```
// $_MUX4_
assign Y = T ? (S ? D : C) :
           (S ? B : A);

// $_MUX8_
assign Y = U ? T ? (S ? H : G) :
           (S ? F : E) :
           T ? (S ? D : C) :
           (S ? B : A);

// $_MUX16_
assign Y = V ? U ? T ? (S ? P : O) :
           (S ? N : M) :
           T ? (S ? L : K) :
           (S ? J : I) :
           U ? T ? (S ? H : G) :
           (S ? F : E) :
           T ? (S ? D : C) :
           (S ? B : A);
```

The cell types `$_DFF_N_` and `$_DFF_P_` represent d-type flip-flops.

The cell types `$_DFFE_[NP][NP]_` implement d-type flip-flops with enable. The values in the table for these cell types relate to the following Verilog code template.

## CHAPTER 5. INTERNAL CELL LIBRARY

Verilog	Cell Type
$Y = A$	<code>\$_BUF_</code>
$Y = \sim A$	<code>\$_NOT_</code>
$Y = A \& B$	<code>\$_AND_</code>
$Y = \sim(A \& B)$	<code>\$_NAND_</code>
$Y = A \& \sim B$	<code>\$_ANDNOT_</code>
$Y = A   B$	<code>\$_OR_</code>
$Y = \sim(A   B)$	<code>\$_NOR_</code>
$Y = A   \sim B$	<code>\$_ORNOT_</code>
$Y = A \wedge B$	<code>\$_XOR_</code>
$Y = \sim(A \wedge B)$	<code>\$_XNOR_</code>
$Y = \sim((A \& B)   C)$	<code>\$_AOI3_</code>
$Y = \sim((A   B) \& C)$	<code>\$_OAI3_</code>
$Y = \sim((A \& B)   (C \& D))$	<code>\$_AOI4_</code>
$Y = \sim((A   B) \& (C   D))$	<code>\$_OAI4_</code>
$Y = S ? B : A$	<code>\$_MUX_</code>
$Y = \sim(S ? B : A)$	<code>\$_NMUX_</code>
(see below)	<code>\$_MUX4_</code>
(see below)	<code>\$_MUX8_</code>
(see below)	<code>\$_MUX16_</code>
$Y = EN ? A : 1'bz$	<code>\$_TBUF_</code>
<b>always</b> @(negedge C) Q <= D	<code>\$_DFF_N_</code>
<b>always</b> @(posedge C) Q <= D	<code>\$_DFF_P_</code>
<b>always</b> @* if (!E) Q <= D	<code>\$_DLATCH_N_</code>
<b>always</b> @* if (E) Q <= D	<code>\$_DLATCH_P_</code>

**Table 5.4:** Cell types for gate level logic networks (main list)

```

always @(ClkEdge c)
    if (EN == EnLvl)
        Q <= D;

```

The cell types `$_DFF_[NP][NP][01]_` implement d-type flip-flops with asynchronous reset. The values in the table for these cell types relate to the following Verilog code template, where *RstEdge* is **posedge** if *RstLvl* is 1, and **negedge** otherwise.

```

always @(ClkEdge c, RstEdge R)
    if (R == RstLvl)
        Q <= RstVal;
    else
        Q <= D;

```

The cell types `$_SDFF_[NP][NP][01]_` implement d-type flip-flops with synchronous reset. The values in the table for these cell types relate to the following Verilog code template:

```

always @(ClkEdge c)
    if (R == RstLvl)
        Q <= RstVal;
    else
        Q <= D;

```

The cell types `$_DFFE_[NP][NP][01][NP]_` implement d-type flip-flops with asynchronous reset and enable. The values in the table for these cell types relate to the following Verilog code template, where *RstEdge* is **posedge** if *RstLvl* is 1, and **negedge** otherwise.

<i>ClkEdge</i>	<i>RstLvl</i>	<i>RstVal</i>	Cell Type
<b>negedge</b>	0	0	<code>\$_DFF_NN0_</code> , <code>\$_SDFF_NN0_</code>
<b>negedge</b>	0	1	<code>\$_DFF_NN1_</code> , <code>\$_SDFF_NN1_</code>
<b>negedge</b>	1	0	<code>\$_DFF_NP0_</code> , <code>\$_SDFF_NP0_</code>
<b>negedge</b>	1	1	<code>\$_DFF_NP1_</code> , <code>\$_SDFF_NP1_</code>
<b>posedge</b>	0	0	<code>\$_DFF_PN0_</code> , <code>\$_SDFF_PN0_</code>
<b>posedge</b>	0	1	<code>\$_DFF_PN1_</code> , <code>\$_SDFF_PN1_</code>
<b>posedge</b>	1	0	<code>\$_DFF_PP0_</code> , <code>\$_SDFF_PP0_</code>
<b>posedge</b>	1	1	<code>\$_DFF_PP1_</code> , <code>\$_SDFF_PP1_</code>

**Table 5.5:** Cell types for gate level logic networks (FFs with reset)

<i>ClkEdge</i>	<i>EnLvl</i>	Cell Type
<b>negedge</b>	0	<code>\$_DFFE_NN_</code>
<b>negedge</b>	1	<code>\$_DFFE_NP_</code>
<b>posedge</b>	0	<code>\$_DFFE_PN_</code>
<b>posedge</b>	1	<code>\$_DFFE_PP_</code>

**Table 5.6:** Cell types for gate level logic networks (FFs with enable)

```

always @(ClkEdge c, RstEdge R)
    if (R == RstLvl)
        Q <= RstVal;
    else if (EN == EnLvl)
        Q <= D;

```

The cell types `$_SDFFE_[NP][NP][01][NP]_` implement d-type flip-flops with synchronous reset and enable, with reset having priority over enable. The values in the table for these cell types relate to the following Verilog code template:

```

always @(ClkEdge c)
    if (R == RstLvl)
        Q <= RstVal;
    else if (EN == EnLvl)
        Q <= D;

```

The cell types `$_SDFFCE_[NP][NP][01][NP]_` implement d-type flip-flops with synchronous reset and enable, with enable having priority over reset. The values in the table for these cell types relate to the following Verilog code template:

```

always @(ClkEdge c)
    if (EN == EnLvl)
        if (R == RstLvl)
            Q <= RstVal;
        else
            Q <= D;

```

The cell types `$_DFFSR_[NP][NP][NP]_` implement d-type flip-flops with asynchronous set and reset. The values in the table for these cell types relate to the following Verilog code template, where *RstEdge* is **posedge** if *RstLvl* if 1, **negedge** otherwise, and *SetEdge* is **posedge** if *SetLvl* if 1, **negedge** otherwise.

```

always @(ClkEdge c, RstEdge R, SetEdge S)
    if (R == RstLvl)
        Q <= 0;
    else if (S == SetLvl)
        Q <= 1;

```

## CHAPTER 5. INTERNAL CELL LIBRARY

<i>ClkEdge</i>	<i>RstLvl</i>	<i>RstVal</i>	<i>EnLvl</i>	Cell Type
<b>negedge</b>	0	0	0	<code>\$_DFFE_NN0N_</code> , <code>\$_SDFFE_NN0N_</code> , <code>\$_SDFFCE_NN0N_</code>
<b>negedge</b>	0	0	1	<code>\$_DFFE_NN0P_</code> , <code>\$_SDFFE_NN0P_</code> , <code>\$_SDFFCE_NN0P_</code>
<b>negedge</b>	0	1	0	<code>\$_DFFE_NN1N_</code> , <code>\$_SDFFE_NN1N_</code> , <code>\$_SDFFCE_NN1N_</code>
<b>negedge</b>	0	1	1	<code>\$_DFFE_NN1P_</code> , <code>\$_SDFFE_NN1P_</code> , <code>\$_SDFFCE_NN1P_</code>
<b>negedge</b>	1	0	0	<code>\$_DFFE_NP0N_</code> , <code>\$_SDFFE_NP0N_</code> , <code>\$_SDFFCE_NP0N_</code>
<b>negedge</b>	1	0	1	<code>\$_DFFE_NP0P_</code> , <code>\$_SDFFE_NP0P_</code> , <code>\$_SDFFCE_NP0P_</code>
<b>negedge</b>	1	1	0	<code>\$_DFFE_NP1N_</code> , <code>\$_SDFFE_NP1N_</code> , <code>\$_SDFFCE_NP1N_</code>
<b>negedge</b>	1	1	1	<code>\$_DFFE_NP1P_</code> , <code>\$_SDFFE_NP1P_</code> , <code>\$_SDFFCE_NP1P_</code>
<b>posedge</b>	0	0	0	<code>\$_DFFE_PN0N_</code> , <code>\$_SDFFE_PN0N_</code> , <code>\$_SDFFCE_PN0N_</code>
<b>posedge</b>	0	0	1	<code>\$_DFFE_PN0P_</code> , <code>\$_SDFFE_PN0P_</code> , <code>\$_SDFFCE_PN0P_</code>
<b>posedge</b>	0	1	0	<code>\$_DFFE_PN1N_</code> , <code>\$_SDFFE_PN1N_</code> , <code>\$_SDFFCE_PN1N_</code>
<b>posedge</b>	0	1	1	<code>\$_DFFE_PN1P_</code> , <code>\$_SDFFE_PN1P_</code> , <code>\$_SDFFCE_PN1P_</code>
<b>posedge</b>	1	0	0	<code>\$_DFFE_PP0N_</code> , <code>\$_SDFFE_PP0N_</code> , <code>\$_SDFFCE_PP0N_</code>
<b>posedge</b>	1	0	1	<code>\$_DFFE_PP0P_</code> , <code>\$_SDFFE_PP0P_</code> , <code>\$_SDFFCE_PP0P_</code>
<b>posedge</b>	1	1	0	<code>\$_DFFE_PP1N_</code> , <code>\$_SDFFE_PP1N_</code> , <code>\$_SDFFCE_PP1N_</code>
<b>posedge</b>	1	1	1	<code>\$_DFFE_PP1P_</code> , <code>\$_SDFFE_PP1P_</code> , <code>\$_SDFFCE_PP1P_</code>

**Table 5.7:** Cell types for gate level logic networks (FFs with reset and enable)

<i>ClkEdge</i>	<i>SetLvl</i>	<i>RstLvl</i>	Cell Type
<b>negedge</b>	0	0	<code>\$_DFFSR_NNN_</code>
<b>negedge</b>	0	1	<code>\$_DFFSR_NNP_</code>
<b>negedge</b>	1	0	<code>\$_DFFSR_NPN_</code>
<b>negedge</b>	1	1	<code>\$_DFFSR_NPP_</code>
<b>posedge</b>	0	0	<code>\$_DFFSR_PNN_</code>
<b>posedge</b>	0	1	<code>\$_DFFSR_PNP_</code>
<b>posedge</b>	1	0	<code>\$_DFFSR_PPN_</code>
<b>posedge</b>	1	1	<code>\$_DFFSR_PPP_</code>

**Table 5.8:** Cell types for gate level logic networks (FFs with set and reset)

```

else
    Q <= D;

```

The cell types `$_DFFSRE_[NP][NP][NP][NP]_` implement d-type flip-flops with asynchronous set and reset and enable. The values in the table for these cell types relate to the following Verilog code template, where *RstEdge* is **posedge** if *RstLvl* is 1, **negedge** otherwise, and *SetEdge* is **posedge** if *SetLvl* is 1, **negedge** otherwise.

```

always @(ClkEdge c, RstEdge r, SetEdge s)
    if (R == RstLvl)
        Q <= 0;
    else if (S == SetLvl)
        Q <= 1;
    else if (E == EnLvl)
        Q <= D;

```

The cell types `$_DLATCH_N_` and `$_DLATCH_P_` represent d-type latches.

The cell types `$_DLATCH_[NP][NP][01]_` implement d-type latches with reset. The values in the table for these cell types relate to the following Verilog code template:

```

always @*
    if (R == RstLvl)
        Q <= RstVal;
    else if (E == EnLvl)

```

## CHAPTER 5. INTERNAL CELL LIBRARY

<i>ClkEdge</i>	<i>SetLvl</i>	<i>RstLvl</i>	<i>EnLvl</i>	Cell Type
<b>negedge</b>	0	0	0	<code>\$_DFFSRE_NNNN_</code>
<b>negedge</b>	0	0	1	<code>\$_DFFSRE_NNPN_</code>
<b>negedge</b>	0	1	0	<code>\$_DFFSRE_NPNP_</code>
<b>negedge</b>	0	1	1	<code>\$_DFFSRE_NPPN_</code>
<b>negedge</b>	1	0	0	<code>\$_DFFSRE_NPNN_</code>
<b>negedge</b>	1	0	1	<code>\$_DFFSRE_NPNP_</code>
<b>negedge</b>	1	1	0	<code>\$_DFFSRE_NPPN_</code>
<b>negedge</b>	1	1	1	<code>\$_DFFSRE_NPPP_</code>
<b>posedge</b>	0	0	0	<code>\$_DFFSRE_PNNN_</code>
<b>posedge</b>	0	0	1	<code>\$_DFFSRE_PNPN_</code>
<b>posedge</b>	0	1	0	<code>\$_DFFSRE_PNP_</code>
<b>posedge</b>	0	1	1	<code>\$_DFFSRE_PNPP_</code>
<b>posedge</b>	1	0	0	<code>\$_DFFSRE_PPNN_</code>
<b>posedge</b>	1	0	1	<code>\$_DFFSRE_PPNP_</code>
<b>posedge</b>	1	1	0	<code>\$_DFFSRE_PPPN_</code>
<b>posedge</b>	1	1	1	<code>\$_DFFSRE_PPPP_</code>

**Table 5.9:** Cell types for gate level logic networks (FFs with set and reset and enable)

<i>EnLvl</i>	<i>RstLvl</i>	<i>RstVal</i>	Cell Type
0	0	0	<code>\$_DLATCH_NN0_</code>
0	0	1	<code>\$_DLATCH_NN1_</code>
0	1	0	<code>\$_DLATCH_NP0_</code>
0	1	1	<code>\$_DLATCH_NP1_</code>
1	0	0	<code>\$_DLATCH_PN0_</code>
1	0	1	<code>\$_DLATCH_PN1_</code>
1	1	0	<code>\$_DLATCH_PP0_</code>
1	1	1	<code>\$_DLATCH_PP1_</code>

**Table 5.10:** Cell types for gate level logic networks (latches with reset)

```
Q <= D;
```

The cell types `$_DLATCHSR_[NP][NP]_` implement d-type latches with set and reset. The values in the table for these cell types relate to the following Verilog code template:

```
always @*
    if (R == RstLvl)
        Q <= 0;
    else if (S == SetLvl)
        Q <= 1;
    else if (E == EnLvl)
        Q <= D;
```

The cell types `$_SR_[NP][NP]_` implement sr-type latches. The values in the table for these cell types relate to the following Verilog code template:

```
always @*
    if (R == RstLvl)
        Q <= 0;
    else if (S == SetLvl)
        Q <= 1;
```

In most cases gate level logic networks are created from RTL networks using the `techmap` pass. The flip-flop cells from the gate level logic network can be mapped to physical flip-flop cells from a Liberty file using

## CHAPTER 5. INTERNAL CELL LIBRARY

<i>EnLvl</i>	<i>SetLvl</i>	<i>RstLvl</i>	Cell Type
0	0	0	<code>\$_DLATCHSR_NNN_</code>
0	0	1	<code>\$_DLATCHSR_NNP_</code>
0	1	0	<code>\$_DLATCHSR_NPN_</code>
0	1	1	<code>\$_DLATCHSR_NPP_</code>
1	0	0	<code>\$_DLATCHSR_PNN_</code>
1	0	1	<code>\$_DLATCHSR_PNP_</code>
1	1	0	<code>\$_DLATCHSR_PPN_</code>
1	1	1	<code>\$_DLATCHSR_PPP_</code>

**Table 5.11:** Cell types for gate level logic networks (latches with set and reset)

<i>SetLvl</i>	<i>RstLvl</i>	Cell Type
0	0	<code>\$_SR_NN_</code>
0	1	<code>\$_SR_NP_</code>
1	0	<code>\$_SR_PN_</code>
1	1	<code>\$_SR_PP_</code>

**Table 5.12:** Cell types for gate level logic networks (SR latches)

the `dfflibmap` pass. The combinatorial logic cells can be mapped to physical cells from a Liberty file via ABC [27] using the `abc` pass.

**FIXME:**

Add information about `$slice` and `$concat` cells.

**FIXME:**

Add information about `$lut` and `$sop` cells.

**FIXME:**

Add information about `$alu`, `$macc`, `$fa`, and `$lcu` cells.

## Chapter 6

# Programming Yosys Extensions

This chapter contains some bits and pieces of information about programming yosys extensions. Also consult the section on programming in the “Yosys Presentation” (can be downloaded from the Yosys website as PDF) and don’t be afraid to ask questions on the YosysHQ Slack.

### 6.1 Guidelines

The `guidelines` directory contains notes on various aspects of Yosys development. The files `GettingStarted` and `CodingStyle` may be of particular interest, and are reproduced here.

#### GettingStarted

```
1  Getting Started
2  =====
3
4
5  Outline of a Yosys command
6  -----
7
8  Here is a the C++ code for a "hello_world" Yosys command (hello.cc):
9
10     #include "kernel/yosys.h"
11
12     USING_YOSYS_NAMESPACE
13     PRIVATE_NAMESPACE_BEGIN
14
15     struct HelloWorldPass : public Pass {
16         HelloWorldPass() : Pass("hello_world") { }
17         void execute(vector<string>, Design*) override {
18             log("Hello World!\n");
19         }
20     } HelloWorldPass;
21
22     PRIVATE_NAMESPACE_END
23
24  This can be built into a Yosys module using the following command:
25
26     yosys-config --exec --cxx --cxxflags --ldflags -o hello.so -shared hello.cc --ldlibs
```



## CHAPTER 6. PROGRAMMING YOSYS EXTENSIONS

```
27
28 Or short:
29
30     yosys-config --build hello.so hello.cc
31
32 And then executed using the following command:
33
34     yosys -m hello.so -p hello_world
35
36
37 Yosys Data Structures
38 -----
39
40 Here is a short list of data structures that you should make yourself familiar
41 with before you write C++ code for Yosys. The following data structures are all
42 defined when "kernel/yosys.h" is included and USING_YOSYS_NAMESPACE is used.
43
44     1. Yosys Container Classes
45
46 Yosys uses dict<K, T> and pool<T> as main container classes. dict<K, T> is
47 essentially a replacement for std::unordered_map<K, T> and pool<T> is a
48 replacement for std::unordered_set<T>. The main characteristics are:
49
50     - dict<K, T> and pool<T> are about 2x faster than the std containers
51
52     - references to elements in a dict<K, T> or pool<T> are invalidated by
53       insert and remove operations (similar to std::vector<T> on push_back()).
54
55     - some iterators are invalidated by erase(). specifically, iterators
56       that have not passed the erased element yet are invalidated. (erase()
57       itself returns valid iterator to the next element.)
58
59     - no iterators are invalidated by insert(). elements are inserted at
60       begin(). i.e. only a new iterator that starts at begin() will see the
61       inserted elements.
62
63     - the method .count(key, iterator) is like .count(key) but only
64       considers elements that can be reached via the iterator.
65
66     - iterators can be compared. it1 < it2 means that the position of t2
67       can be reached via t1 but not vice versa.
68
69     - the method .sort() can be used to sort the elements in the container
70       the container stays sorted until elements are added or removed.
71
72     - dict<K, T> and pool<T> will have the same order of iteration across
73       all compilers, standard libraries and architectures.
74
75 In addition to dict<K, T> and pool<T> there is also an idict<K> that
76 creates a bijective map from K to the integers. For example:
77
78     idict<string, 42> si;
79     log("%d\n", si("hello"));      // will print 42
80     log("%d\n", si("world"));      // will print 43
```

## CHAPTER 6. PROGRAMMING YOSYS EXTENSIONS

```
81     log("%d\n", si.at("world"));    // will print 43
82     log("%d\n", si.at("dummy"));    // will throw exception
83     log("%s\n", si[42].c_str());    // will print hello
84     log("%s\n", si[43].c_str());    // will print world
85     log("%s\n", si[44].c_str());    // will throw exception
86
87 It is not possible to remove elements from an idict.
88
89 Finally mfp<K> implements a merge-find set data structure (aka. disjoint-set or
90 union-find) over the type K ("mfp" = merge-find-promote).
91
92     2. Standard STL data types
93
94 In Yosys we use std::vector<T> and std::string whenever applicable. When
95 dict<K, T> and pool<T> are not suitable then std::map<K, T> and std::set<T>
96 are used instead.
97
98 The types std::vector<T> and std::string are also available as vector<T>
99 and string in the Yosys namespace.
100
101     3. RTLIL objects
102
103 The current design (essentially a collection of modules, each defined by a
104 netlist) is stored in memory using RTLIL object (declared in kernel/rtlil.h,
105 automatically included by kernel/yosys.h). You should glance over at least
106 the declarations for the following types in kernel/rtlil.h:
107
108     RTLIL::IdString
109         This is a handle for an identifier (e.g. cell or wire name).
110         It feels a lot like a std::string, but is only a single int
111         in size. (The actual string is stored in a global lookup
112         table.)
113
114     RTLIL::SigBit
115         A single signal bit. I.e. either a constant state (0, 1,
116         x, z) or a single bit from a wire.
117
118     RTLIL::SigSpec
119         Essentially a vector of SigBits.
120
121     RTLIL::Wire
122     RTLIL::Cell
123         The building blocks of the netlist in a module.
124
125     RTLIL::Module
126     RTLIL::Design
127         The module is a container with connected cells and wires
128         in it. The design is a container with modules in it.
129
130 All this types are also available without the RTLIL:: prefix in the Yosys
131 namespace.
132
133     4. SigMap and other Helper Classes
134
```

## CHAPTER 6. PROGRAMMING YOSYS EXTENSIONS

There are a couple of additional helper classes that are in wide use in Yosys. Most importantly there is SigMap (declared in kernel/sigtools.h).

When a design has many wires in it that are connected to each other, then a single signal bit can have multiple valid names. The SigMap object can be used to map SigSpecs or SigBits to unique SigSpecs and SigBits that consistently only use one wire from such a group of connected wires. For example:

```
SigBit a = module->addWire(NEW_ID);
SigBit b = module->addWire(NEW_ID);
module->connect(a, b);

log("%d\n", a == b); // will print 0

SigMap sigmap(module);
log("%d\n", sigmap(a) == sigmap(b)); // will print 1
```

Using the RTLIL Netlist Format

In the RTLIL netlist format the cell ports contain SigSpecs that point to the Wires. There are no references in the other direction. This has two direct consequences:

- (1) It is very easy to go from cells to wires but hard to go in the other way.
- (2) There is no danger in removing cells from the netlists, but removing wires can break the netlist format when there are still references to the wire somewhere in the netlist.

The solution to (1) is easy: Create custom indexes that allow you to make fast lookups for the wire-to-cell direction. You can either use existing generic index structures to do that (such as the ModIndex class) or write your own index. For many application it is simplest to construct a custom index. For example:

```
SigMap sigmap(module);
dict<SigBit, Cell*> sigbit_to_driver_index;

for (auto cell : module->cells())
    for (auto &conn : cell->connections())
        if (cell->output(conn.first))
            for (auto bit : sigmap(conn.second))
                sigbit_to_driver_index[bit] = cell;
```

Regarding (2): There is a general theme in Yosys that you don't remove wires from the design. You can rename them, unconnect them, but you do not actually remove the Wire object from the module. Instead you let the "clean" command take care of the dangling wires. On the other hand it is safe to remove cells (as long as you make sure this does not invalidate a custom index you are using in your code).

Example Code

## CHAPTER 6. PROGRAMMING YOSYS EXTENSIONS

```
189 -----
190
191 The following yosys commands are a good starting point if you are looking for examples
192 of how to use the Yosys API:
193
194     manual/CHAPTER_Prog/stubnets.cc
195     manual/PRESENTATION_Prog/my_cmd.cc
196
197
198 Script Passes
199 -----
200
201 The ScriptPass base class can be used to implement passes that just call other passes,
202 like a script. Examples for such passes are:
203
204     techlibs/common/prep.cc
205     techlibs/common/synth.cc
206
207 In some cases it is easier to implement such a pass as regular pass, for example when
208 ScriptPass doesn't provide the type of flow control desired. (But many of the
209 script passes in Yosys that don't use ScriptPass simply predate the ScriptPass base
210 class.) Examples for such passes are:
211
212     passes/opt/opt.cc
213     passes/proc/proc.cc
214
215 Whether they use the ScriptPass base-class or not, a pass should always either
216 call other passes without doing any non-trivial work itself, or should implement
217 a non-trivial algorithm but not call any other passes. The reason for this is that
218 this helps containing complexity in individual passes and simplifies debugging the
219 entire system.
220
221 Exceptions to this rule should be rare and limited to cases where calling other
222 passes is optional and only happens when requested by the user (such as for
223 example `techmap -autoproc`), or where it is about commands that are "top-level
224 commands" in their own right, not components to be used in regular synthesis
225 flows (such as the `bugpoint` command).
226
227 A pass that would "naturally" call other passes and also do some work itself
228 should be re-written in one of two ways:
229
230 1) It could be re-written as script pass with the parts that are not calls
231 to other passes factored out into individual new passes. Usually in those
232 cases the new sub passes share the same prefix as the top-level script pass.
233
234 2) It could be re-written so that it already expects the design in a certain
235 state, expecting the calling script to set up this state before calling the
236 pass in questions.
237
238 Many back-ends are examples for the 2nd approach. For example, `write_aiger`
239 does not convert the design into AIG representation, but expects the design
240 to be already in this form, and prints an `Unsupported cell type` error
241 message otherwise.
242
```

```

243
244 Notes on the existing codebase
245 -----
246
247 For historical reasons not all parts of Yosys adhere to the current coding
248 style. When adding code to existing parts of the system, adhere to this guide
249 for the new code instead of trying to mimic the style of the surrounding code.

```

#### CodingStyle

```

1 Coding Style
2 =====
3
4
5 Formatting of code
6 -----
7
8 - Yosys code is using tabs for indentation. Tabs are 8 characters.
9
10 - A continuation of a statement in the following line is indented by
11   two additional tabs.
12
13 - Lines are as long as you want them to be. A good rule of thumb is
14   to break lines at about column 150.
15
16 - Opening braces can be put on the same or next line as the statement
17   opening the block (if, switch, for, while, do). Put the opening brace
18   on its own line for larger blocks, especially blocks that contains
19   blank lines.
20
21 - Otherwise stick to the Linux Kernel Coding Style:
22   https://www.kernel.org/doc/Documentation/CodingStyle
23
24
25 C++ Language
26 -----
27
28 Yosys is written in C++11. At the moment only constructs supported by
29 gcc 4.8 are allowed in Yosys code. This will change in future releases.
30
31 In general Yosys uses "int" instead of "size_t". To avoid compiler
32 warnings for implicit type casts, always use "GetSize(foobar)" instead
33 of "foobar.size()". (GetSize() is defined in kernel/yosys.h)
34
35 Use range-based for loops whenever applicable.

```

## 6.2 The “stubsnets” Example Module

The following is the complete code of the “stubsnets” example module. It is included in the Yosys source distribution as `manual/CHAPTER_Prog/stubnets.cc`.

stubsnets.cc

## CHAPTER 6. PROGRAMMING YOSYS EXTENSIONS

```
1 // This is free and unencumbered software released into the public domain.
2 //
3 // Anyone is free to copy, modify, publish, use, compile, sell, or
4 // distribute this software, either in source code form or as a compiled
5 // binary, for any purpose, commercial or non-commercial, and by any
6 // means.
7
8 #include "kernel/yosys.h"
9 #include "kernel/sigtools.h"
10
11 #include <string>
12 #include <map>
13 #include <set>
14
15 USING_YOSYS_NAMESPACE
16 PRIVATE_NAMESPACE_BEGIN
17
18 // this function is called for each module in the design
19 static void find_stub_nets(RTLIL::Design *design, RTLIL::Module *module, bool report_bits)
20 {
21     // use a SigMap to convert nets to a unique representation
22     SigMap sigmap(module);
23
24     // count how many times a single-bit signal is used
25     std::map<RTLIL::SigBit, int> bit_usage_count;
26
27     // count output lines for this module (needed only for summary output at the end)
28     int line_count = 0;
29
30     log("Looking for stub wires in module %s:\n", RTLIL::id2cstr(module->name));
31
32     // For all ports on all cells
33     for (auto &cell_iter : module->cells_)
34     for (auto &conn : cell_iter.second->connections())
35     {
36         // Get the signals on the port
37         // (use sigmap to get a unique signal name)
38         RTLIL::SigSpec sig = sigmap(conn.second);
39
40         // add each bit to bit_usage_count, unless it is a constant
41         for (auto &bit : sig)
42             if (bit.wire != NULL)
43                 bit_usage_count[bit]++;
44     }
45
46     // for each wire in the module
47     for (auto &wire_iter : module->wires_)
48     {
49         RTLIL::Wire *wire = wire_iter.second;
50
51         // .. but only selected wires
52         if (!design->selected(module, wire))
53             continue;
54     }
```

## CHAPTER 6. PROGRAMMING YOSYS EXTENSIONS

```

55 // add +1 usage if this wire actually is a port
56 int usage_offset = wire->port_id > 0 ? 1 : 0;
57
58 // we will record which bits of the (possibly multi-bit) wire are stub signals
59 std::set<int> stub_bits;
60
61 // get a signal description for this wire and split it into separate bits
62 RTLIL::SigSpec sig = sigmap(wire);
63
64 // for each bit (unless it is a constant):
65 // check if it is used at least two times and add to stub_bits otherwise
66 for (int i = 0; i < GetSize(sig); i++)
67     if (sig[i].wire != NULL && (bit_usage_count[sig[i]] + usage_offset) < 2)
68         stub_bits.insert(i);
69
70 // continue if no stub bits found
71 if (stub_bits.size() == 0)
72     continue;
73
74 // report stub bits and/or stub wires, don't report single bits
75 // if called with report_bits set to false.
76 if (GetSize(stub_bits) == GetSize(sig)) {
77     log("__found_stub_wire:_%s\n", RTLIL::id2cstr(wire->name));
78 } else {
79     if (!report_bits)
80         continue;
81     log("__found_wire_with_stub_bits:_%s_", RTLIL::id2cstr(wire->name));
82     for (int bit : stub_bits)
83         log("%sd", bit == *stub_bits.begin() ? "" : ", ", bit);
84     log("]\n");
85 }
86
87 // we have outputted a line, increment summary counter
88 line_count++;
89 }
90
91 // report summary
92 if (report_bits)
93     log("__found_%d_stub_wires_or_wires_with_stub_bits.\n", line_count);
94 else
95     log("__found_%d_stub_wires.\n", line_count);
96 }
97
98 // each pass contains a singleton object that is derived from Pass
99 struct StubnetsPass : public Pass {
100     StubnetsPass() : Pass("stubnets") { }
101     void execute(std::vector<std::string> args, RTLIL::Design *design) override
102     {
103         // variables to mirror information from passed options
104         bool report_bits = 0;
105
106         log_header(design, "Executing_STUBNETS_pass_(find_stub_nets).\n");
107
108         // parse options

```

## CHAPTER 6. PROGRAMMING YOSYS EXTENSIONS

```
109         size_t argidx;
110         for (argidx = 1; argidx < args.size(); argidx++) {
111             std::string arg = args[argidx];
112             if (arg == "--report_bits") {
113                 report_bits = true;
114                 continue;
115             }
116             break;
117         }
118
119         // handle extra options (e.g. selection)
120         extra_args(args, argidx, design);
121
122         // call find_stub_nets() for each module that is either
123         // selected as a whole or contains selected objects.
124         for (auto &it : design->modules_)
125             if (design->selected_module(it.first))
126                 find_stub_nets(design, it.second, report_bits);
127     }
128 } StubnetsPass;
129
130 PRIVATE_NAMESPACE_END
```

### Makefile

```
1 test: stubnets.so
2     yosys -ql test1.log -m ./stubnets.so test.v -p "stubnets"
3     yosys -ql test2.log -m ./stubnets.so test.v -p "opt;_stubnets"
4     yosys -ql test3.log -m ./stubnets.so test.v -p "techmap;_opt;_stubnets_--report_bits"
5     tail test1.log test2.log test3.log
6
7 stubnets.so: stubnets.cc
8     yosys-config --exec --cxx --cxxflags --ldflags -o $@ -shared $^ --ldlibs
9
10 clean:
11     rm -f test1.log test2.log test3.log
12     rm -f stubnets.so stubnets.d
```

### test.v

```
1 module uut(in1, in2, in3, out1, out2);
2
3 input [8:0] in1, in2, in3;
4 output [8:0] out1, out2;
5
6 assign out1 = in1 + in2 + (in3 >> 4);
7
8 endmodule
```



## Chapter 7

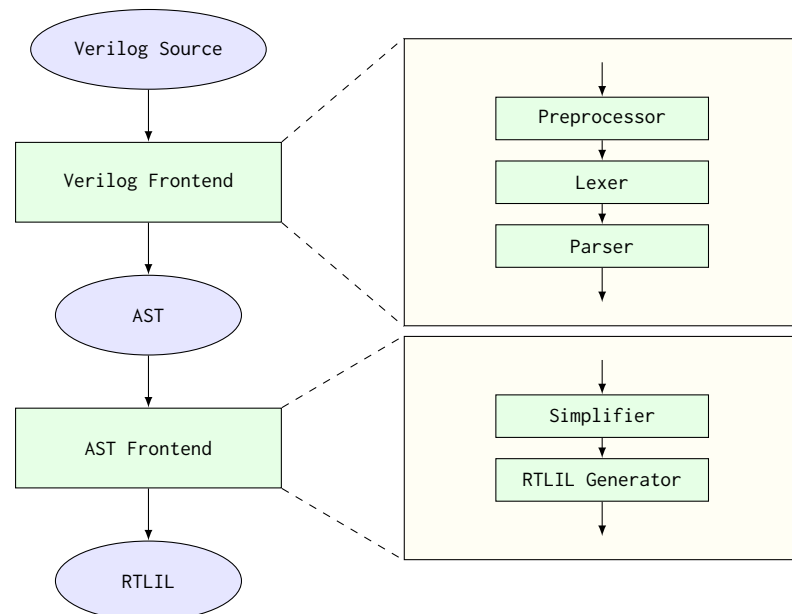
# The Verilog and AST Frontends

This chapter provides an overview of the implementation of the Yosys Verilog and AST frontends. The Verilog frontend reads Verilog-2005 code and creates an abstract syntax tree (AST) representation of the input. This AST representation is then passed to the AST frontend that converts it to RTLIL data, as illustrated in Fig. 7.1.

### 7.1 Transforming Verilog to AST

The *Verilog frontend* converts the Verilog sources to an internal AST representation that closely resembles the structure of the original Verilog code. The Verilog frontend consists of three components, the *Preprocessor*, the *Lexer* and the *Parser*.

The source code to the Verilog frontend can be found in `frontends/verilog/` in the Yosys source tree.



**Figure 7.1:** Simplified Verilog to RTLIL data flow

### 7.1.1 The Verilog Preprocessor

The Verilog preprocessor scans over the Verilog source code and interprets some of the Verilog compiler directives such as ``include`, ``define` and ``ifdef`.

It is implemented as a C++ function that is passed a file descriptor as input and returns the pre-processed Verilog code as a `std::string`.

The source code to the Verilog Preprocessor can be found in `frontends/verilog/preproc.cc` in the Yosys source tree.

### 7.1.2 The Verilog Lexer

The Verilog Lexer is written using the lexer generator *flex* [17]. Its source code can be found in `frontends/verilog/verilog_lexer.l` in the Yosys source tree. The lexer does little more than identifying all keywords and literals recognised by the Yosys Verilog frontend.

The lexer keeps track of the current location in the Verilog source code using some global variables. These variables are used by the constructor of AST nodes to annotate each node with the source code location it originated from.

Finally the lexer identifies and handles special comments such as `“// synopsys translate_off”` and `“// synopsys full_case”`. (It is recommended to use ``ifdef` constructs instead of the Synopsys `translate_on/off` comments and attributes such as `(* full_case *)` over `“// synopsys full_case”` whenever possible.)

### 7.1.3 The Verilog Parser

The Verilog Parser is written using the parser generator *bison* [18]. Its source code can be found in `frontends/verilog/verilog_parser.y` in the Yosys source tree.

It generates an AST using the `AST::AstNode` data structure defined in `frontends/ast/ast.h`. An `AST::AstNode` object has the following properties:

- **The node type**  
This enum (`AST::AstNodeType`) specifies the role of the node. Table 7.1 contains a list of all node types.
- **The child nodes**  
This is a list of pointers to all children in the abstract syntax tree.
- **Attributes**  
As almost every AST node might have Verilog attributes assigned to it, the `AST::AstNode` has direct support for attributes. Note that the attribute values are again AST nodes.
- **Node content**  
Each node might have additional content data. A series of member variables exist to hold such data. For example the member `std::string str` can hold a string value and is used e.g. in the `AST_IDENTIFIER` node type to store the identifier name.
- **Source code location**  
Each `AST::AstNode` is automatically annotated with the current source code location by the `AST::AstNode` constructor. It is stored in the `std::string filename` and `int linenum` member variables.

The `AST::AstNode` constructor can be called with up to two child nodes that are automatically added to the list of child nodes for the new object. This simplifies the creation of AST nodes for simple expressions a bit. For example the bison code for parsing multiplications:

```

1      basic_expr '*' attr basic_expr {
2          $$ = new AstNode(AST_MUL, $1, $4);
3          append_attr($$, $3);
4      } |

```

The generated AST data structure is then passed directly to the AST frontend that performs the actual conversion to RTLIL.

Note that the Yosys command `read_verilog` provides the options `-yydebug` and `-dump_ast` that can be used to print the parse tree or abstract syntax tree respectively.

## 7.2 Transforming AST to RTLIL

The *AST Frontend* converts a set of modules in AST representation to modules in RTLIL representation and adds them to the current design. This is done in two steps: *simplification* and *RTLIL generation*.

The source code to the AST frontend can be found in `frontends/ast/` in the Yosys source tree.

### 7.2.1 AST Simplification

A full-featured AST is too complex to be transformed into RTLIL directly. Therefore it must first be brought into a simpler form. This is done by calling the `AST::AstNode::simplify()` method of all `AST_MODULE` nodes in the AST. This initiates a recursive process that performs the following transformations on the AST data structure:

AST Node Type	Corresponding Verilog Construct
AST_NONE	This Node type should never be used.
AST_DESIGN	This node type is used for the top node of the AST tree. It has no corresponding Verilog construct.
AST_MODULE, AST_TASK, AST_FUNCTION	<b>module</b> , <b>task</b> and <b>function</b>
AST_WIRE	<b>input</b> , <b>output</b> , <b>wire</b> , <b>reg</b> and <b>integer</b>
AST_MEMORY	Verilog Arrays
AST_AUTOWIRE	Created by the simplifier when an undeclared signal name is used.
AST_PARAMETER, AST_LOCALPARAM	<b>parameter</b> and <b>localparam</b>
AST_PARASET	Parameter set in cell instantiation
AST_ARGUMENT	Port connection in cell instantiation
AST_RANGE	Bit-Index in a signal or element index in array
AST_CONSTANT	A literal value
AST_CELLTYPE	The type of cell in cell instantiation
AST_IDENTIFIER	An Identifier (signal name in expression or cell/task/etc. name in other contexts)
AST_PREFIX	Construct an identifier in the form <prefix>[<index>].<suffix> (used only in advanced generate constructs)
AST_FCALL, AST_TCALL	Call to function or task
AST_TO_SIGNED, AST_TO_UNSIGNED	The <code>\$signed()</code> and <code>\$unsigned()</code> functions

**Table 7.1:** AST node types with their corresponding Verilog constructs.  
(continued on next page)

## CHAPTER 7. THE VERILOG AND AST FRONTENDS

AST Node Type	Corresponding Verilog Construct
AST_CONCAT AST_REPLICATE	The <code>{...}</code> and <code>{...{...}}</code> operators
AST_BIT_NOT, AST_BIT_AND, AST_BIT_OR, AST_BIT_XOR, AST_BIT_XNOR	The bitwise operators <code>~, &amp;,  , ^</code> and <code>~^</code>
AST_REDUCE_AND, AST_REDUCE_OR, AST_REDUCE_XOR, AST_REDUCE_XNOR	The unary reduction operators <code>~, &amp;,  , ^</code> and <code>~^</code>
AST_REDUCE_BOOL	Conversion from multi-bit value to boolean value (equivalent to <code>AST_REDUCE_OR</code> )
AST_SHIFT_LEFT, AST_SHIFT_RIGHT, AST_SHIFT_SLEFT, AST_SHIFT_SRIGHT	The shift operators <code>&lt;&lt;, &gt;&gt;, &lt;&lt;&lt; and &gt;&gt;&gt;</code>
AST_LT, AST_LE, AST_EQ, AST_NE, AST_GE, AST_GT	The relational operators <code>&lt;, &lt;=, ==, !=, &gt;= and &gt;</code>
AST_ADD, AST_SUB, AST_MUL, AST_DIV, AST_MOD, AST_POW	The binary operators <code>+, -, *, /, %</code> and <code>**</code>
AST_POS, AST_NEG	The prefix operators <code>+</code> and <code>-</code>
AST_LOGIC_AND, AST_LOGIC_OR, AST_LOGIC_NOT	The logic operators <code>&amp;&amp;,   </code> and <code>!</code>
AST_TERNARY	The ternary <code>?:-</code> operator
AST_MEMRD AST_MEMWR	Read and write memories. These nodes are generated by the AST simplifier for writes/reads to/from Verilog arrays.
AST_ASSIGN	An <b>assign</b> statement
AST_CELL	A cell instantiation
AST_PRIMITIVE	A primitive cell ( <b>and</b> , <b>nand</b> , <b>or</b> , etc.)
AST_ALWAYS, AST_INITIAL	Verilog <b>always</b> - and <b>initial</b> -blocks
AST_BLOCK	A <b>begin-end</b> -block
AST_ASSIGN_EQ, AST_ASSIGN_LE	Blocking ( <code>=</code> ) and nonblocking ( <code>&lt;=</code> ) assignments within an <b>always</b> - or <b>initial</b> -block
AST_CASE, AST_COND, AST_DEFAULT	The <b>case</b> ( <b>if</b> ) statements, conditions within a case and the default case respectively
AST_FOR	A <b>for</b> -loop with an <b>always</b> - or <b>initial</b> -block
AST_GENVAR, AST_GENBLOCK, AST_GENFOR, AST_GENIF	The <b>genvar</b> and <b>generate</b> keywords and <b>for</b> and <b>if</b> within a generate block.
AST_POSEDGE, AST_NEGEDGE, AST_EDGE	Event conditions for <b>always</b> blocks.

**Table 7.1:** AST node types with their corresponding Verilog constructs.  
(continuation from previous page)

- Inline all task and function calls.
- Evaluate all **generate**-statements and unroll all **for**-loops.
- Perform const folding where it is necessary (e.g. in the value part of `AST_PARAMETER`, `AST_LOCALPARAM`, `AST_PARASET` and `AST_RANGE` nodes).
- Replace `AST_PRIMITIVE` nodes with appropriate `AST_ASSIGN` nodes.
- Replace dynamic bit ranges in the left-hand-side of assignments with `AST_CASE` nodes with `AST_COND` children for each possible case.
- Detect array access patterns that are too complicated for the `RTLIL::Memory` abstraction and replace them with a set of signals and cases for all reads and/or writes.
- Otherwise replace array accesses with `AST_MEMRD` and `AST_MEMWR` nodes.

In addition to these transformations, the simplifier also annotates the AST with additional information that is needed for the RTLIL generator, namely:

- All ranges (width of signals and bit selections) are not only const folded but (when a constant value is found) are also written to member variables in the `AST_RANGE` node.
- All identifiers are resolved and all `AST_IDENTIFIER` nodes are annotated with a pointer to the AST node that contains the declaration of the identifier. If no declaration has been found, an `AST_AUTOWIRE` node is created and used for the annotation.

This produces an AST that is fairly easy to convert to the RTLIL format.

### 7.2.2 Generating RTLIL

After AST simplification, the `AST::AstNode::genRTLIL()` method of each `AST_MODULE` node in the AST is called. This initiates a recursive process that generates equivalent RTLIL data for the AST data.

The `AST::AstNode::genRTLIL()` method returns an `RTLIL::SigSpec` structure. For nodes that represent expressions (operators, constants, signals, etc.), the cells needed to implement the calculation described by the expression are created and the resulting signal is returned. That way it is easy to generate the circuits for large expressions using depth-first recursion. For nodes that do not represent an expression (such as `AST_CELL`), the corresponding circuit is generated and an empty `RTLIL::SigSpec` is returned.

## 7.3 Synthesizing Verilog always Blocks

For behavioural Verilog code (code utilizing **always**- and **initial**-blocks) it is necessary to also generate `RTLIL::Process` objects. This is done in the following way:

- Whenever `AST::AstNode::genRTLIL()` encounters an **always**- or **initial**-block, it creates an instance of `AST_INTERNAL::ProcessGenerator`. This object then generates the `RTLIL::Process` object for the block. It also calls `AST::AstNode::genRTLIL()` for all right-hand-side expressions contained within the block.
- First the `AST_INTERNAL::ProcessGenerator` creates a list of all signals assigned within the block. It then creates a set of temporary signals using the naming scheme `$<number>\<original_name>` for each of the assigned signals.
- Then an `RTLIL::Process` is created that assigns all intermediate values for each left-hand-side signal to the temporary signal in its `RTLIL::CaseRule/RTLIL::SwitchRule` tree.
- Finally a `RTLIL::SyncRule` is created for the `RTLIL::Process` that assigns the temporary signals for the final values to the actual signals.
- A process may also contain memory writes. A `RTLIL::MemWriteAction` is created for each of them.
- Calls to `AST::AstNode::genRTLIL()` are generated for right hand sides as needed. When blocking assignments are used, `AST::AstNode::genRTLIL()` is configured using global variables to use the temporary signals that hold the correct intermediate values whenever one of the previously assigned signals is used in an expression.

Unfortunately the generation of a correct `RTLIL::CaseRule/RTLIL::SwitchRule` tree for behavioural code is a non-trivial task. The AST frontend solves the problem using the approach described on the following pages. The following example illustrates what the algorithm is supposed to do. Consider the following Verilog code:

## CHAPTER 7. THE VERILOG AND AST FRONTENDS

```

1  always @(posedge clock) begin
2      out1 = in1;
3      if (in2)
4          out1 = !out1;
5      out2 <= out1;
6      if (in3)
7          out2 <= out2;
8      if (in4)
9          if (in5)
10             out3 <= in6;
11             else
12                 out3 <= in7;
13      out1 = out1 ^ out2;
14  end

```

This is translated by the Verilog and AST frontends into the following RTLIL code (attributes, cell parameters and wire declarations not included):

```

1  cell $logic_not $logic_not$<input>:4$2
2      connect \A \in1
3      connect \Y $logic_not$<input>:4$2_Y
4  end
5  cell $xor $xor$<input>:13$3
6      connect \A $1\out1[0:0]
7      connect \B \out2
8      connect \Y $xor$<input>:13$3_Y
9  end
10 process $proc$<input>:1$1
11     assign $0\out3[0:0] \out3
12     assign $0\out2[0:0] $1\out1[0:0]
13     assign $0\out1[0:0] $xor$<input>:13$3_Y
14     switch \in2
15         case 1'1
16             assign $1\out1[0:0] $logic_not$<input>:4$2_Y
17         case
18             assign $1\out1[0:0] \in1
19     end
20     switch \in3
21         case 1'1
22             assign $0\out2[0:0] \out2
23         case
24     end
25     switch \in4
26         case 1'1
27             switch \in5
28                 case 1'1
29                     assign $0\out3[0:0] \in6
30                 case
31                     assign $0\out3[0:0] \in7
32             end
33         case
34     end
35     sync posedge \clock
36     update \out1 $0\out1[0:0]

```

```

37     update \out2 $0\out2[0:0]
38     update \out3 $0\out3[0:0]
39 end

```

Note that the two operators are translated into separate cells outside the generated process. The signal `out1` is assigned using blocking assignments and therefore `out1` has been replaced with a different signal in all expressions after the initial assignment. The signal `out2` is assigned using nonblocking assignments and therefore is not substituted on the right-hand-side expressions.

The `RTLIL::CaseRule/RTLIL::SwitchRule` tree must be interpreted the following way:

- On each case level (the body of the process is the *root case*), first the actions on this level are evaluated and then the switches within the case are evaluated. (Note that the last assignment on line 13 of the Verilog code has been moved to the beginning of the RTLIL process to line 13 of the RTLIL listing.)  
I.e. the special cases deeper in the switch hierarchy override the defaults on the upper levels. The assignments in lines 12 and 22 of the RTLIL code serve as an example for this.  
Note that in contrast to this, the order within the `RTLIL::SwitchRule` objects within a `RTLIL::CaseRule` is preserved with respect to the original AST and Verilog code.
- The whole `RTLIL::CaseRule/RTLIL::SwitchRule` tree describes an asynchronous circuit. I.e. the decision tree formed by the switches can be seen independently for each assigned signal. Whenever one assigned signal changes, all signals that depend on the changed signals are to be updated. For example the assignments in lines 16 and 18 in the RTLIL code in fact influence the assignment in line 12, even though they are in the “wrong order”.

The only synchronous part of the process is in the `RTLIL::SyncRule` object generated at line 35 in the RTLIL code. The sync rule is the only part of the process where the original signals are assigned. The synchronization event from the original Verilog code has been translated into the synchronization type (`posedge`) and signal (`\clock`) for the `RTLIL::SyncRule` object. In the case of this simple example the `RTLIL::SyncRule` object is later simply transformed into a set of d-type flip-flops and the `RTLIL::CaseRule/RTLIL::SwitchRule` tree to a decision tree using multiplexers.

In more complex examples (e.g. asynchronous resets) the part of the `RTLIL::CaseRule/RTLIL::SwitchRule` tree that describes the asynchronous reset must first be transformed to the correct `RTLIL::SyncRule` objects. This is done by the `proc_adff` pass.

### 7.3.1 The ProcessGenerator Algorithm

The `AST_INTERNAL::ProcessGenerator` uses the following internal state variables:

- `subst_rvalue_from` and `subst_rvalue_to`  
These two variables hold the replacement pattern that should be used by `AST::AstNode::genRTLIL()` for signals with blocking assignments. After initialization of `AST_INTERNAL::ProcessGenerator` these two variables are empty.
- `subst_lvalue_from` and `subst_lvalue_to`  
These two variables contain the mapping from left-hand-side signals (`\<name>`) to the current temporary signal for the same thing (initially `$0\<name>`).
- `current_case`  
A pointer to a `RTLIL::CaseRule` object. Initially this is the root case of the generated `RTLIL::Process`.

## CHAPTER 7. THE VERILOG AND AST FRONTENDS

As the algorithm runs these variables are continuously modified as well as pushed to the stack and later restored to their earlier values by popping from the stack.

On startup the ProcessGenerator generates a new `RTLIL::Process` object with an empty root case and initializes its state variables as described above. Then the `RTLIL::SyncRule` objects are created using the synchronization events from the `AST_ALWAYS` node and the initial values of `subst_lvalue_from` and `subst_lvalue_to`. Then the AST for this process is evaluated recursively.

During this recursive evaluation, three different relevant types of AST nodes can be discovered: `AST_ASSIGN_LE` (nonblocking assignments), `AST_ASSIGN_EQ` (blocking assignments) and `AST_CASE` (**if** or **case** statement).

### 7.3.1.1 Handling of Nonblocking Assignments

When an `AST_ASSIGN_LE` node is discovered, the following actions are performed by the ProcessGenerator:

- The left-hand-side is evaluated using `AST::AstNode::genRTLIL()` and mapped to a temporary signal name using `subst_lvalue_from` and `subst_lvalue_to`.
- The right-hand-side is evaluated using `AST::AstNode::genRTLIL()`. For this call, the values of `subst_rvalue_from` and `subst_rvalue_to` are used to map blocking-assigned signals correctly.
- Remove all assignments to the same left-hand-side as this assignment from the `current_case` and all cases within it.
- Add the new assignment to the `current_case`.

### 7.3.1.2 Handling of Blocking Assignments

When an `AST_ASSIGN_EQ` node is discovered, the following actions are performed by the ProcessGenerator:

- Perform all the steps that would be performed for a nonblocking assignment (see above).
- Remove the found left-hand-side (before lvalue mapping) from `subst_rvalue_from` and also remove the respective bits from `subst_rvalue_to`.
- Append the found left-hand-side (before lvalue mapping) to `subst_rvalue_from` and append the found right-hand-side to `subst_rvalue_to`.

### 7.3.1.3 Handling of Cases and if-Statements

When an `AST_CASE` node is discovered, the following actions are performed by the ProcessGenerator:

- The values of `subst_rvalue_from`, `subst_rvalue_to`, `subst_lvalue_from` and `subst_lvalue_to` are pushed to the stack.
- A new `RTLIL::SwitchRule` object is generated, the selection expression is evaluated using `AST::AstNode::genRTLIL()` (with the use of `subst_rvalue_from` and `subst_rvalue_to`) and added to the `RTLIL::SwitchRule` object and the object is added to the `current_case`.
- All lvalues assigned to within the `AST_CASE` node using blocking assignments are collected and saved in the local variable `this_case_eq_lvalue`.
- New temporary signals are generated for all signals in `this_case_eq_lvalue` and stored in `this_case_eq_ltemp`.



## CHAPTER 7. THE VERILOG AND AST FRONTENDS

- The signals in `this_case_eq_lvalue` are mapped using `subst_rvalue_from` and `subst_rvalue_to` and the resulting set of signals is stored in `this_case_eq_rvalue`.

Then the following steps are performed for each `AST_COND` node within the `AST_CASE` node:

- Set `subst_rvalue_from`, `subst_rvalue_to`, `subst_lvalue_from` and `subst_lvalue_to` to the values that have been pushed to the stack.
- Remove `this_case_eq_lvalue` from `subst_lvalue_from/subst_lvalue_to`.
- Append `this_case_eq_lvalue` to `subst_lvalue_from` and append `this_case_eq_ltemp` to `subst_lvalue_to`.
- Push the value of `current_case`.
- Create a new `RTLIL::CaseRule`. Set `current_case` to the new object and add the new object to the `RTLIL::SwitchRule` created above.
- Add an assignment from `this_case_eq_rvalue` to `this_case_eq_ltemp` to the new `current_case`.
- Evaluate the compare value for this case using `AST::AstNode::genRTLIL()` (with the use of `subst_rvalue_from` and `subst_rvalue_to`) modify the new `current_case` accordingly.
- Recursion into the children of the `AST_COND` node.
- Restore `current_case` by popping the old value from the stack.

Finally the following steps are performed:

- The values of `subst_rvalue_from`, `subst_rvalue_to`, `subst_lvalue_from` and `subst_lvalue_to` are popped from the stack.
- The signals from `this_case_eq_lvalue` are removed from the `subst_rvalue_from/subst_rvalue_to`-pair.
- The value of `this_case_eq_lvalue` is appended to `subst_rvalue_from` and the value of `this_case_eq_ltemp` is appended to `subst_rvalue_to`.
- Map the signals in `this_case_eq_lvalue` using `subst_lvalue_from/subst_lvalue_to`.
- Remove all assignments to signals in `this_case_eq_lvalue` in `current_case` and all cases within it.
- Add an assignment from `this_case_eq_ltemp` to `this_case_eq_lvalue` to `current_case`.

### 7.3.1.4 Further Analysis of the Algorithm for Cases and if-Statements

With respect to nonblocking assignments the algorithm is easy: later assignments invalidate earlier assignments. For each signal assigned using nonblocking assignments exactly one temporary variable is generated (with the `$0`-prefix) and this variable is used for all assignments of the variable.

Note how all the `_eq`-variables become empty when no blocking assignments are used and many of the steps in the algorithm can then be ignored as a result of this.

For a variable with blocking assignments the algorithm shows the following behaviour: First a new temporary variable is created. This new temporary variable is then registered as the assignment target for all assignments for this variable within the cases for this `AST_CASE` node. Then for each case the new temporary variable is first assigned the old temporary variable. This assignment is overwritten if the variable is actually assigned in this case and is kept as a default value otherwise.

This yields an `RTLIL::CaseRule` that assigns the new temporary variable in all branches. So when all cases have been processed a final assignment is added to the containing block that assigns the new temporary variable to the old one. Note how this step always overrides a previous assignment to the old temporary variable. Other than nonblocking assignments, the old assignment could still have an effect somewhere in the design, as there have been calls to `AST::AstNode::genRTLIL()` with a `subst_rvalue_from/subst_rvalue_to`-tuple that contained the right-hand-side of the old assignment.

### 7.3.2 The proc pass

The ProcessGenerator converts a behavioural model in AST representation to a behavioural model in RTLIL::Process representation. The actual conversion from a behavioural model to an RTL representation is performed by the proc pass and the passes it launches:

- `proc_clean` and `proc_rmdead`  
These two passes just clean up the RTLIL::Process structure. The `proc_clean` pass removes empty parts (eg. empty assignments) from the process and `proc_rmdead` detects and removes unreachable branches from the process's decision trees.
- `proc_arst`  
This pass detects processes that describe d-type flip-flops with asynchronous resets and rewrites the process to better reflect what they are modelling: Before this pass, an asynchronous reset has two edge-sensitive sync rules and one top-level RTLIL::SwitchRule for the reset path. After this pass the sync rule for the reset is level-sensitive and the top-level RTLIL::SwitchRule has been removed.
- `proc_mux`  
This pass converts the RTLIL::CaseRule/RTLIL::SwitchRule-tree to a tree of multiplexers per written signal. After this, the RTLIL::Process structure only contains the RTLIL::SyncRules that describe the output registers.
- `proc_dff`  
This pass replaces the RTLIL::SyncRules to d-type flip-flops (with asynchronous resets if necessary).
- `proc_dff`  
This pass replaces the RTLIL::MemWriteActionss with `$memwr` cells.
- `proc_clean`  
A final call to `proc_clean` removes the now empty RTLIL::Process objects.

Performing these last processing steps in passes instead of in the Verilog frontend has two important benefits:

First it improves the transparency of the process. Everything that happens in a separate pass is easier to debug, as the RTLIL data structures can be easily investigated before and after each of the steps.

Second it improves flexibility. This scheme can easily be extended to support other types of storage-elements, such as sr-latches or d-latches, without having to extend the actual Verilog frontend.

## 7.4 Synthesizing Verilog Arrays

### FIXME:

Add some information on the generation of `$memrd` and `$memwr` cells and how they are processed in the memory pass.

## 7.5 Synthesizing Parametric Designs

### FIXME:

Add some information on the RTLIL::Module::derive() method and how it is used to synthesize parametric modules via the hierarchy pass.

## Chapter 8

# Optimizations

Yosys employs a number of optimizations to generate better and cleaner results. This chapter outlines these optimizations.

### 8.1 Simple Optimizations

The Yosys pass `opt` runs a number of simple optimizations. This includes removing unused signals and cells and const folding. It is recommended to run this pass after each major step in the synthesis script. At the time of this writing the `opt` pass executes the following passes that each perform a simple optimization:

- Once at the beginning of `opt`:
  - `opt_expr`
  - `opt_merge -nomux`
- Repeat until result is stable:
  - `opt_muxtree`
  - `opt_reduce`
  - `opt_merge`
  - `opt_rmdff`
  - `opt_clean`
  - `opt_expr`

The following section describes each of the `opt_*` passes.

#### 8.1.1 The `opt_expr` pass

This pass performs const folding on the internal combinational cell types described in Chap. 5. This means a cell with all constant inputs is replaced with the constant value this cell drives. In some cases this pass can also optimize cells with some constant inputs.

Table 8.1 shows the replacement rules used for optimizing an `$AND` gate. The first three rules implement the obvious const folding rules. Note that ‘any’ might include dynamic values calculated by other parts of the circuit. The following three lines propagate undef (X) states. These are the only three cases in which it is allowed to propagate an undef according to Sec. 5.1.10 of IEEE Std. 1364-2005 [Ver06].

## CHAPTER 8. OPTIMIZATIONS

A-Input	B-Input	Replacement
any	0	0
0	any	0
1	1	1
X/Z	X/Z	X
1	X/Z	X
X/Z	1	X
any	X/Z	0
X/Z	any	0
$a$	1	$a$
1	$b$	$b$

**Table 8.1:** Const folding rules for `$AND_` cells as used in `opt_expr`.

The next two lines assume the value 0 for undef states. These two rules are only used if no other substitutions are possible in the current module. If other substitutions are possible they are performed first, in the hope that the ‘any’ will change to an undef value or a 1 and therefore the output can be set to undef.

The last two lines simply replace an `$AND_` gate with one constant-1 input with a buffer.

Besides this basic const folding the `opt_expr` pass can replace 1-bit wide `$eq` and `$ne` cells with buffers or not-gates if one input is constant.

The `opt_expr` pass is very conservative regarding optimizing `$mux` cells, as these cells are often used to model decision-trees and breaking these trees can interfere with other optimizations.

### 8.1.2 The `opt_muxtree` pass

This pass optimizes trees of multiplexer cells by analyzing the select inputs. Consider the following simple example:

```

1 module uut(a, y);
2   input a;
3   output [1:0] y = a ? (a ? 1 : 2) : 3;
4 endmodule

```

The output can never be 2, as this would require `a` to be 1 for the outer multiplexer and 0 for the inner multiplexer. The `opt_muxtree` pass detects this contradiction and replaces the inner multiplexer with a constant 1, yielding the logic for `y = a ? 1 : 3`.

### 8.1.3 The `opt_reduce` pass

This is a simple optimization pass that identifies and consolidates identical input bits to `$reduce_and` and `$reduce_or` cells. It also sorts the input bits to ease identification of shareable `$reduce_and` and `$reduce_or` cells in other passes.

This pass also identifies and consolidates identical inputs to multiplexer cells. In this case the new shared select bit is driven using a `$reduce_or` cell that combines the original select bits.

Lastly this pass consolidates trees of `$reduce_and` cells and trees of `$reduce_or` cells to single large `$reduce_and` or `$reduce_or` cells.

These three simple optimizations are performed in a loop until a stable result is produced.

### 8.1.4 The `opt_rmdff` pass

This pass identifies single-bit d-type flip-flops (`$_DFF_*`, `$dff`, and `$adff` cells) with a constant data input and replaces them with a constant driver.

### 8.1.5 The `opt_clean` pass

This pass identifies unused signals and cells and removes them from the design. It also creates an `\unused_bits` attribute on wires with unused bits. This attribute can be used for debugging or by other optimization passes.

### 8.1.6 The `opt_merge` pass

This pass performs trivial resource sharing. This means that this pass identifies cells with identical inputs and replaces them with a single instance of the cell.

The option `-nomux` can be used to disable resource sharing for multiplexer cells (`$mux` and `$pmux`). This can be useful as it prevents multiplexer trees to be merged, which might prevent `opt_muxtree` to identify possible optimizations.

## 8.2 FSM Extraction and Encoding

The `fsm` pass performs finite-state-machine (FSM) extraction and recoding. The `fsm` pass simply executes the following other passes:

- Identify and extract FSMs:
  - `fsm_detect`
  - `fsm_extract`
- Basic optimizations:
  - `fsm_opt`
  - `opt_clean`
  - `fsm_opt`
- Expanding to nearby gate-logic (if called with `-expand`):
  - `fsm_expand`
  - `opt_clean`
  - `fsm_opt`
- Re-code FSM states (unless called with `-norecode`):
  - `fsm_recode`
- Print information about FSMs:
  - `fsm_info`
- Export FSMs in KISS2 file format (if called with `-export`):
  - `fsm_export`

- Map FSMs to RTL cells (unless called with `-nomap`):

— fsm\_map

The `fsm_detect` pass identifies FSM state registers and marks them using the `\fsm_encoding= "auto"` attribute. The `fsm_extract` extracts all FSMs marked using the `\fsm_encoding` attribute (unless `\fsm_encoding` is set to "none") and replaces the corresponding RTL cells with a `$fsm` cell. All other `fsm_*` passes operate on these `$fsm` cells. The `fsm_map` call finally replaces the `$fsm` cells with RTL cells.

Note that these optimizations operate on an RTL netlist. I.e. the `fsm` pass should be executed after the `proc` pass has transformed all `RTLIL::Process` objects to RTL cells.

The algorithms used for FSM detection and extraction are influenced by a more general reported technique [STGR10].

### 8.2.1 FSM Detection

The `fsm_detect` pass identifies FSM state registers. It sets the `\fsm_encoding= "auto"` attribute on any (multi-bit) wire that matches the following description:

- Does not already have the `\fsm_encoding` attribute.
- Is not an output of the containing module.
- Is driven by single `$dff` or `$adff` cell.
- The `\D`-Input of this `$dff` or `$adff` cell is driven by a multiplexer tree that only has constants or the old state value on its leaves.
- The state value is only used in the said multiplexer tree or by simple relational cells that compare the state value to a constant (usually `$eq` cells).

This heuristic has proven to work very well. It is possible to overwrite it by setting `\fsm_encoding= "auto"` on registers that should be considered FSM state registers and setting `\fsm_encoding= "none"` on registers that match the above criteria but should not be considered FSM state registers.

Note however that marking state registers with `\fsm_encoding` that are not suitable for FSM recoding can cause synthesis to fail or produce invalid results.

### 8.2.2 FSM Extraction

The `fsm_extract` pass operates on all state signals marked with the `\fsm_encoding (!= "none")` attribute. For each state signal the following information is determined:

- The state registers
- The asynchronous reset state if the state registers use asynchronous reset
- All states and the control input signals used in the state transition functions
- The control output signals calculated from the state signals and control inputs
- A table of all state transitions and corresponding control inputs- and outputs

## CHAPTER 8. OPTIMIZATIONS

The state registers (and asynchronous reset state, if applicable) is simply determined by identifying the driver for the state signal.

From there the `$mux-tree` driving the state register inputs is recursively traversed. All select inputs are control signals and the leaves of the `$mux-tree` are the states. The algorithm fails if a non-constant leaf that is not the state signal itself is found.

The list of control outputs is initialized with the bits from the state signal. It is then extended by adding all values that are calculated by cells that compare the state signal with a constant value.

In most cases this will cover all uses of the state register, thus rendering the state encoding arbitrary. If however a design uses e.g. a single bit of the state value to drive a control output directly, this bit of the state signal will be transformed to a control output of the same value.

Finally, a transition table for the FSM is generated. This is done by using the `ConstEval` C++ helper class (defined in `kernel/consteval.h`) that can be used to evaluate parts of the design. The `ConstEval` class can be asked to calculate a given set of result signals using a set of signal-value assignments. It can also be passed a list of stop-signals that abort the `ConstEval` algorithm if the value of a stop-signal is needed in order to calculate the result signals.

The `fsm_extract` pass uses the `ConstEval` class in the following way to create a transition table. For each state:

1. Create a `ConstEval` object for the module containing the FSM
2. Add all control inputs to the list of stop signals
3. Set the state signal to the current state
4. Try to evaluate the next state and control output
5. If step 4 was not successful:
  - Recursively goto step 4 with the offending stop-signal set to 0.
  - Recursively goto step 4 with the offending stop-signal set to 1.
6. If step 4 was successful: Emit transition

Finally a `$fsm` cell is created with the generated transition table and added to the module. This new cell is connected to the control signals and the old drivers for the control outputs are disconnected.

### 8.2.3 FSM Optimization

The `fsm_opt` pass performs basic optimizations on `$fsm` cells (not including state recoding). The following optimizations are performed (in this order):

- Unused control outputs are removed from the `$fsm` cell. The attribute `\unused_bits` (that is usually set by the `opt_clean` pass) is used to determine which control outputs are unused.
- Control inputs that are connected to the same driver are merged.
- When a control input is driven by a control output, the control input is removed and the transition table altered to give the same performance without the external feedback path.
- Entries in the transition table that yield the same output and only differ in the value of a single control input bit are merged and the different bit is removed from the sensitivity list (turned into a don't-care bit).
- Constant inputs are removed and the transition table is altered to give an unchanged behaviour.
- Unused inputs are removed.

### 8.2.4 FSM Recoding

The `fsm_recode` pass assigns new bit pattern to the states. Usually this also implies a change in the width of the state signal. At the moment of this writing only one-hot encoding with all-zero for the reset state is supported.

The `fsm_recode` pass can also write a text file with the changes performed by it that can be used when verifying designs synthesized by Yosys using Synopsys Formality [24].

## 8.3 Logic Optimization

Yosys can perform multi-level combinational logic optimization on gate-level netlists using the external program ABC [27]. The `abc` pass extracts the combinational gate-level parts of the design, passes it through ABC, and re-integrates the results. The `abc` pass can also be used to perform other operations using ABC, such as technology mapping (see Sec. 9.3 for details).



## Chapter 9

# Technology Mapping

Previous chapters outlined how HDL code is transformed into an RTL netlist. The RTL netlist is still based on abstract coarse-grain cell types like arbitrary width adders and even multipliers. This chapter covers how an RTL netlist is transformed into a functionally equivalent netlist utilizing the cell types available in the target architecture.

Technology mapping is often performed in two phases. In the first phase RTL cells are mapped to an internal library of single-bit cells (see Sec. 5.2). In the second phase this netlist of internal gate types is transformed to a netlist of gates from the target technology library.

When the target architecture provides coarse-grain cells (such as block ram or ALUs), these must be mapped to directly form the RTL netlist, as information on the coarse-grain structure of the design is lost when it is mapped to bit-width gate types.

### 9.1 Cell Substitution

The simplest form of technology mapping is cell substitution, as performed by the `techmap` pass. This pass, when provided with a Verilog file that implements the RTL cell types using simpler cells, simply replaces the RTL cells with the provided implementation.

When no map file is provided, `techmap` uses a built-in map file that maps the Yosys RTL cell types to the internal gate library used by Yosys. The curious reader may find this map file as `techlibs/common/techmap.v` in the Yosys source tree.

Additional features have been added to `techmap` to allow for conditional mapping of cells (see `help techmap` or Sec. C.209). This can for example be useful if the target architecture supports hardware multipliers for certain bit-widths but not for others.

A usual synthesis flow would first use the `techmap` pass to directly map some RTL cells to coarse-grain cells provided by the target architecture (if any) and then use `techmap` with the built-in default file to map the remaining RTL cells to gate logic.

### 9.2 Subcircuit Substitution

Sometimes the target architecture provides cells that are more powerful than the RTL cells used by Yosys. For example a cell in the target architecture that can calculate the absolute-difference of two numbers does not match any single RTL cell type but only combinations of cells.

For these cases Yosys provides the `extract` pass that can match a given set of modules against a design and identify the portions of the design that are identical (i.e. isomorphic subcircuits) to any of the given modules. These matched subcircuits are then replaced by instances of the given modules.

The `extract` pass also finds basic variations of the given modules, such as swapped inputs on commutative cell types.

In addition to this the `extract` pass also has limited support for frequent subcircuit mining, i.e. the process of finding recurring subcircuits in the design. This has a few applications, including the design of new coarse-grain architectures [GW13].

The hard algorithmic work done by the `extract` pass (solving the isomorphic subcircuit problem and frequent subcircuit mining) is performed using the SubCircuit library that can also be used stand-alone without Yosys (see Sec. A.3).

### 9.3 Gate-Level Technology Mapping

On the gate-level the target architecture is usually described by a “Liberty file”. The Liberty file format is an industry standard format that can be used to describe the behaviour and other properties of standard library cells [25].

Mapping a design utilizing the Yosys internal gate library (e.g. as a result of mapping it to this representation using the `techmap` pass) is performed in two phases.

First the register cells must be mapped to the registers that are available on the target architectures. The target architecture might not provide all variations of d-type flip-flops with positive and negative clock edge, high-active and low-active asynchronous set and/or reset, etc. Therefore the process of mapping the registers might add additional inverters to the design and thus it is important to map the register cells first.

Mapping of the register cells may be performed by using the `dfflibmap` pass. This pass expects a Liberty file as argument (using the `-liberty` option) and only uses the register cells from the Liberty file.

Secondly the combinational logic must be mapped to the target architecture. This is done using the external program ABC [27] via the `abc` pass by using the `-liberty` option to the pass. Note that in this case only the combinatorial cells are used from the cell library.

Occasionally Liberty files contain trade secrets (such as sensitive timing information) that cannot be shared freely. This complicates processes such as reporting bugs in the tools involved. When the information in the Liberty file used by Yosys and ABC are not part of the sensitive information, the additional tool `yosys-filterlib` (see Sec. B.2) can be used to strip the sensitive information from the Liberty file.

# Appendix A

## Auxiliary Libraries

The Yosys source distribution contains some auxiliary libraries that are bundled with Yosys.

### A.1 SHA1

The files in `libs/sha1/` provide a public domain SHA1 implementation written by Steve Reid, Bruce Guenter, and Volker Grabsch. It is used for generating unique names when specializing parameterized modules.

### A.2 BigInt

The files in `libs/bigint/` provide a library for performing arithmetic with arbitrary length integers. It is written by Matt McCutchen [29].

The BigInt library is used for evaluating constant expressions, e.g. using the `ConstEval` class provided in `kernel/consteval.h`.

### A.3 SubCircuit

The files in `libs/subcircuit` provide a library for solving the subcircuit isomorphism problem. It is written by C. Wolf and based on the Ullmann Subgraph Isomorphism Algorithm [Ull76]. It is used by the `extract` pass (see `help extract` or Sec. C.63).

### A.4 ezSAT

The files in `libs/ezsat` provide a library for simplifying generating CNF formulas for SAT solvers. It also contains bindings of MiniSAT. The ezSAT library is written by C. Wolf. It is used by the `sat` pass (see `help sat` or Sec. C.169).

## Appendix B

# Auxiliary Programs

Besides the main `yosys` executable, the Yosys distribution contains a set of additional helper programs.

### B.1 `yosys-config`

The `yosys-config` tool (an auto-generated shell-script) can be used to query compiler options and other information needed for building loadable modules for Yosys. FIXME: See Sec. 6 for details.

### B.2 `yosys-filterlib`

The `yosys-filterlib` tool is a small utility that can be used to strip or extract information from a Liberty file. See Sec. 9.3 for details.

### B.3 `yosys-abc`

This is a fork of ABC [27] with a small set of custom modifications that have not yet been accepted upstream. Not all versions of Yosys work with all versions of ABC. So Yosys comes with its own `yosys-abc` to avoid compatibility issues between the two.

## Appendix C

# Command Reference Manual

### C.1 abc – use ABC for technology mapping

```
1  abc [options] [selection]
2
3  This pass uses the ABC tool [1] for technology mapping of yosys's internal gate
4  library to a target architecture.
5
6  -exe <command>
7      use the specified command instead of "<yosys-bindir>/yosys-abc" to execute ABC.
8      This can e.g. be used to call a specific version of ABC or a wrapper.
9
10 -script <file>
11     use the specified ABC script file instead of the default script.
12
13     if <file> starts with a plus sign (+), then the rest of the filename
14     string is interpreted as the command string to be passed to ABC. The
15     leading plus sign is removed and all commas (,) in the string are
16     replaced with blanks before the string is passed to ABC.
17
18     if no -script parameter is given, the following scripts are used:
19
20     for -liberty/-genlib without -constr:
21         strash; ifraig; scorr; dc2; dretime; strash; &get -n; &dch -f;
22         &nf {D}; &put
23
24     for -liberty/-genlib with -constr:
25         strash; ifraig; scorr; dc2; dretime; strash; &get -n; &dch -f;
26         &nf {D}; &put; buffer; upsize {D}; dnsiz {D}; stime -p
27
28     for -lut/-luts (only one LUT size):
29         strash; ifraig; scorr; dc2; dretime; strash; dch -f; if; mfs2;
30         lutpack {S}
31
32     for -lut/-luts (different LUT sizes):
33         strash; ifraig; scorr; dc2; dretime; strash; dch -f; if; mfs2
34
35     for -sop:
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

36     strash; ifraig; scorr; dc2; dretime; strash; dch -f;
37     cover {I} {P}
38
39 otherwise:
40     strash; ifraig; scorr; dc2; dretime; strash; &get -n; &dch -f;
41     &nf {D}; &put
42
43 -fast
44     use different default scripts that are slightly faster (at the cost
45     of output quality):
46
47     for -liberty/-genlib without -constr:
48         strash; dretime; map {D}
49
50     for -liberty/-genlib with -constr:
51         strash; dretime; map {D}; buffer; upsize {D}; dnsiz {D};
52         stime -p
53
54     for -lut/-luts:
55         strash; dretime; if
56
57     for -sop:
58         strash; dretime; cover {I} {P}
59
60 otherwise:
61     strash; dretime; map
62
63 -liberty <file>
64     generate netlists for the specified cell library (using the liberty
65     file format).
66
67 -genlib <file>
68     generate netlists for the specified cell library (using the SIS Genlib
69     file format).
70
71 -constr <file>
72     pass this file with timing constraints to ABC.
73     use with -liberty/-genlib.
74
75     a constr file contains two lines:
76         set_driving_cell <cell_name>
77         set_load <floating_point_number>
78
79     the set_driving_cell statement defines which cell type is assumed to
80     drive the primary inputs and the set_load statement sets the load in
81     femtofarads for each primary output.
82
83 -D <picoseconds>
84     set delay target. the string {D} in the default scripts above is
85     replaced by this option when used, and an empty string otherwise.
86     this also replaces 'dretime' with 'dretime; retime -o {D}' in the
87     default scripts above.
88
89 -I <num>

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

maximum number of SOP inputs.
(replaces {I} in the default scripts above)

-P <num>
maximum number of SOP products.
(replaces {P} in the default scripts above)

-S <num>
maximum number of LUT inputs shared.
(replaces {S} in the default scripts above, default: -S 1)

-lut <width>
generate netlist using luts of (max) the specified width.

-lut <w1>:<w2>
generate netlist using luts of (max) the specified width <w2>. All
luts with width <= <w1> have constant cost. for luts larger than <w1>
the area cost doubles with each additional input bit. the delay cost
is still constant for all lut widths.

-luts <cost1>,<cost2>,<cost3>,<sizeN>:<cost4-N>,...
generate netlist using luts. Use the specified costs for luts with 1,
2, 3, .. inputs.

-sop
map to sum-of-product cells and inverters

-g type1,type2,...
Map to the specified list of gate types. Supported gates types are:
AND, NAND, OR, NOR, XOR, XNOR, ANDNOT, ORNOT, MUX,
NMUX, AOI3, OAI3, AOI4, OAI4.
(The NOT gate is always added to this list automatically.)

The following aliases can be used to reference common sets of gate types:
simple: AND OR XOR MUX
cmos2:  NAND NOR
cmos3:  NAND NOR AOI3 OAI3
cmos4:  NAND NOR AOI3 OAI3 AOI4 OAI4
cmos:   NAND NOR AOI3 OAI3 AOI4 OAI4 NMUX MUX XOR XNOR
gates:  AND NAND OR NOR XOR XNOR ANDNOT ORNOT
aig:    AND NAND OR NOR ANDNOT ORNOT

The alias 'all' represent the full set of all gate types.

Prefix a gate type with a '-' to remove it from the list. For example
the arguments 'AND,OR,XOR' and 'simple,-MUX' are equivalent.

The default is 'all,-NMUX,-AOI3,-OAI3,-AOI4,-OAI4'.

-dff
also pass $_DFF_?_ and $_DFFE_??_ cells through ABC. modules with many
clock domains are automatically partitioned in clock domains and each
domain is passed through ABC independently.

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
144 -clk [!]<clock-signal-name>[,!]<enable-signal-name>]
145     use only the specified clock domain. this is like -dff, but only FF
146     cells that belong to the specified clock domain are used.
147
148 -keepff
149     set the "keep" attribute on flip-flop output wires. (and thus preserve
150     them, for example for equivalence checking.)
151
152 -nocleanup
153     when this option is used, the temporary files created by this pass
154     are not removed. this is useful for debugging.
155
156 -showtmp
157     print the temp dir name in log. usually this is suppressed so that the
158     command output is identical across runs.
159
160 -markgroups
161     set a 'abcgroupp' attribute on all objects created by ABC. The value of
162     this attribute is a unique integer for each ABC process started. This
163     is useful for debugging the partitioning of clock domains.
164
165 -dress
166     run the 'dress' command after all other ABC commands. This aims to
167     preserve naming by an equivalence check between the original and post-ABC
168     netlists (experimental).
169
170 When no target cell library is specified the Yosys standard cell library is
171 loaded into ABC before the ABC script is executed.
172
173 Note that this is a logic optimization pass within Yosys that is calling ABC
174 internally. This is not going to "run ABC on your design". It will instead run
175 ABC on logic snippets extracted from your design. You will not get any useful
176 output when passing an ABC script that writes a file. Instead write your full
177 design as BLIF file with write_blif and then load that into ABC externally if
178 you want to use ABC to convert your design into another format.
179
180 [1] http://www.eecs.berkeley.edu/~alanmi/abc/
```

### C.2 abc9 – use ABC9 for technology mapping

```
1     abc9 [options] [selection]
2
3     This script pass performs a sequence of commands to facilitate the use of the ABC
4     tool [1] for technology mapping of the current design to a target FPGA
5     architecture. Only fully-selected modules are supported.
6
7     -run <from_label>:<to_label>
8         only run the commands between the labels (see below). an empty
9         from label is synonymous to 'begin', and empty to label is
10        synonymous to the end of the command list.
11
12     -exe <command>
```



## APPENDIX C. COMMAND REFERENCE MANUAL

13        use the specified command instead of "<yosys-bindir>/yosys-abc" to execute ABC.  
14        This can e.g. be used to call a specific version of ABC or a wrapper.  
15  
16        -script <file>  
17            use the specified ABC script file instead of the default script.  
18  
19            if <file> starts with a plus sign (+), then the rest of the filename  
20            string is interpreted as the command string to be passed to ABC. The  
21            leading plus sign is removed and all commas (,) in the string are  
22            replaced with blanks before the string is passed to ABC.  
23  
24            if no -script parameter is given, the following scripts are used:  
25                &scorr; &sweep; &dc2; &dch -f; &ps; &if {C} {W} {D} {R} -v; &mfs  
26  
27        -fast  
28            use different default scripts that are slightly faster (at the cost  
29            of output quality):  
30                &if {C} {W} {D} {R} -v  
31  
32        -D <picoseconds>  
33            set delay target. the string {D} in the default scripts above is  
34            replaced by this option when used, and an empty string otherwise  
35            (indicating best possible delay).  
36  
37        -lut <width>  
38            generate netlist using luts of (max) the specified width.  
39  
40        -lut <w1>:<w2>  
41            generate netlist using luts of (max) the specified width <w2>. All  
42            luts with width <= <w1> have constant cost. for luts larger than <w1>  
43            the area cost doubles with each additional input bit. the delay cost  
44            is still constant for all lut widths.  
45  
46        -lut <file>  
47            pass this file with lut library to ABC.  
48  
49        -luts <cost1>,<cost2>,<cost3>,<sizeN>:<cost4-N>,...  
50            generate netlist using luts. Use the specified costs for luts with 1,  
51            2, 3, .. inputs.  
52  
53        -maxlut <width>  
54            when auto-generating the lut library, discard all luts equal to or  
55            greater than this size (applicable when neither -lut nor -luts is  
56            specified).  
57  
58        -dff  
59            also pass \$\_DFF\_[NP]\_ cells through to ABC. modules with many clock  
60            domains are supported and automatically partitioned by ABC.  
61  
62        -nocleanup  
63            when this option is used, the temporary files created by this pass  
64            are not removed. this is useful for debugging.  
65  
66        -showtmp

## APPENDIX C. COMMAND REFERENCE MANUAL

```

67     print the temp dir name in log. usually this is suppressed so that the
68     command output is identical across runs.
69
70     -box <file>
71         pass this file with box library to ABC.
72
73 Note that this is a logic optimization pass within Yosys that is calling ABC
74 internally. This is not going to "run ABC on your design". It will instead run
75 ABC on logic snippets extracted from your design. You will not get any useful
76 output when passing an ABC script that writes a file. Instead write your full
77 design as an XAIGER file with `write_xaiger' and then load that into ABC
78 externally if you want to use ABC to convert your design into another format.
79
80 [1] http://www.eecs.berkeley.edu/~alanmi/abc/
81
82
83     check:
84         abc9_ops -check [-dff]      (option if -dff)
85
86     map:
87         abc9_ops -prep_hier [-dff]   (option if -dff)
88         scc -specify -set_attr abc9_scc_id {}
89         abc9_ops -prep_bypass [-prep_dff] (option if -dff)
90         design -stash $abc9
91         design -load $abc9_map
92         proc
93         wbflip
94         techmap -wb -map %$abc9 -map +/techmap.v A:abc9_flop
95         opt -nodffe -nosdff
96         abc9_ops -prep_dff_submod                                     (only if -dff)
97         setattr -set submod "$abc9_flop" t:$_DFF?_ %ci* %co* t:$_DFF?_ %d (only if -dff)
98         submod                                                       (only if -dff)
99         setattr -mod -set whitebox 1 -set abc9_flop 1 -set abc9_box 1 *$_abc9_flop (only if -dff)
100        foreach module in design
101            rename <module-name>_$_abc9_flop _TECHMAP_REPLACE_ (only if -dff)
102        abc9_ops -prep_dff_unmap (only if -dff)
103        design -copy-to $abc9 =*_$_abc9_flop (only if -dff)
104        delete =*_$_abc9_flop (only if -dff)
105        design -stash $abc9_map
106        design -load $abc9
107        design -delete $abc9
108        techmap -wb -max_iter 1 -map %$abc9_map -map +/abc9_map.v [-D DFF] (option if -dff)
109        design -delete $abc9_map
110
111     pre:
112         read_verilog -icells -lib -specify +/abc9_model.v
113         abc9_ops -break_scc -prep_delays -prep_xaiger [-dff] (option for -dff)
114         abc9_ops -prep_lut <maxlut> (skip if -lut or -luts)
115         abc9_ops -prep_box (skip if -box)
116         design -stash $abc9
117         design -load $abc9_holes
118         techmap -wb -map %$abc9 -map +/techmap.v
119         opt -purge
120         aigmap

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
121     design -stash $abc9_holes
122     design -load $abc9
123     design -delete $abc9
124
125     exe:
126     aigmap
127     foreach module in selection
128         abc9_ops -write_lut <abc-temp-dir>/input.lut      (skip if '-lut' or '-luts')
129         abc9_ops -write_box <abc-temp-dir>/input.box      (skip if '-box')
130         write_xaiger -map <abc-temp-dir>/input.sym [-dff] <abc-temp-dir>/input.xaig
131         abc9_exe [options] -cwd <abc-temp-dir> -lut [<abc-temp-dir>/input.lut] -box [<abc-temp-dir>/inp
132         read_aiger -xaiger -wideports -module_name <module-name>$abc9 -map <abc-temp-dir>/input.sym <ab
133         abc9_ops -reintegrate [-dff]
134
135     unmap:
136         techmap -wb -map %$abc9_unmap -map +/abc9_unmap.v
137         design -delete $abc9_unmap
138         design -delete $abc9_holes
139         delete =*_$abc9_byp
140         setattr -mod -unset abc9_box_id
```

### C.3 abc9\_exe – use ABC9 for technology mapping

```
1     abc9_exe [options]
2
3
4     This pass uses the ABC tool [1] for technology mapping of the top module
5     (according to the (* top *) attribute or if only one module is currently selected)
6     to a target FPGA architecture.
7
8     -exe <command>
9         use the specified command instead of "<yosys-bindir>/yosys-abc" to execute ABC.
10        This can e.g. be used to call a specific version of ABC or a wrapper.
11
12     -script <file>
13         use the specified ABC script file instead of the default script.
14
15         if <file> starts with a plus sign (+), then the rest of the filename
16         string is interpreted as the command string to be passed to ABC. The
17         leading plus sign is removed and all commas (,) in the string are
18         replaced with blanks before the string is passed to ABC.
19
20         if no -script parameter is given, the following scripts are used:
21         &scorr; &sweep; &dc2; &dch -f; &ps; &if {C} {W} {D} {R} -v; &mfs
22
23     -fast
24         use different default scripts that are slightly faster (at the cost
25         of output quality):
26         &if {C} {W} {D} {R} -v
27
28     -D <picoseconds>
29         set delay target. the string {D} in the default scripts above is
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
30     replaced by this option when used, and an empty string otherwise
31     (indicating best possible delay).
32
33     -lut <width>
34         generate netlist using luts of (max) the specified width.
35
36     -lut <w1>:<w2>
37         generate netlist using luts of (max) the specified width <w2>. All
38         luts with width <= <w1> have constant cost. for luts larger than <w1>
39         the area cost doubles with each additional input bit. the delay cost
40         is still constant for all lut widths.
41
42     -lut <file>
43         pass this file with lut library to ABC.
44
45     -luts <cost1>,<cost2>,<cost3>,<sizeN>:<cost4-N>,...
46         generate netlist using luts. Use the specified costs for luts with 1,
47         2, 3, .. inputs.
48
49     -showtmp
50         print the temp dir name in log. usually this is suppressed so that the
51         command output is identical across runs.
52
53     -box <file>
54         pass this file with box library to ABC.
55
56     -cwd <dir>
57         use this as the current working directory, inside which the 'input.xaig'
58         file is expected. temporary files will be created in this directory, and
59         the mapped result will be written to 'output.aig'.
60
61 Note that this is a logic optimization pass within Yosys that is calling ABC
62 internally. This is not going to "run ABC on your design". It will instead run
63 ABC on logic snippets extracted from your design. You will not get any useful
64 output when passing an ABC script that writes a file. Instead write your full
65 design as BLIF file with write_blif and then load that into ABC externally if
66 you want to use ABC to convert your design into another format.
67
68 [1] http://www.eecs.berkeley.edu/~alanmi/abc/
```

### C.4 abc9\_ops – helper functions for ABC9

```
1     abc9_ops [options] [selection]
2
3 This pass contains a set of supporting operations for use during ABC technology
4 mapping, and is expected to be called in conjunction with other operations from
5 the `abc9' script pass. Only fully-selected modules are supported.
6
7     -check
8         check that the design is valid, e.g. (* abc9_box_id *) values are unique,
9         (* abc9_carry *) is only given for one input/output port, etc.
10
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

11 -prep_hier
12     derive all used (* abc9_box *) or (* abc9_flop *) (if -dff option)
13     whitebox modules. with (* abc9_flop *) modules, only those containing
14     $dff/$_DFF_[NP]_ cells with zero initial state -- due to an ABC limitation
15     -- will be derived.
16
17 -prep_bypass
18     create techmap rules in the '$abc9_map' and '$abc9_unmap' designs for
19     bypassing sequential (* abc9_box *) modules using a combinatorial box
20     (named *_abc9_byp). bypassing is necessary if sequential elements (e.g.
21     $dff, $mem, etc.) are discovered inside so that any combinatorial paths
22     will be correctly captured. this bypass box will only contain ports that
23     are referenced by a simple path declaration ($specify2 cell) inside a
24     specify block.
25
26 -prep_dff
27     select all (* abc9_flop *) modules instantiated in the design and store
28     in the named selection '$abc9_flops'.
29
30 -prep_dff_submod
31     within (* abc9_flop *) modules, rewrite all edge-sensitive path
32     declarations and $setup() timing checks ($specify3 and $specrule cells)
33     that share a 'DST' port with the $_DFF_[NP]_.Q port from this 'Q' port to
34     the DFF's 'D' port. this is to prepare such specify cells to be moved
35     into the flop box.
36
37 -prep_dff_unmap
38     populate the '$abc9_unmap' design with techmap rules for mapping *_abc9_flop
39     cells back into their derived cell types (where the rules created by
40     -prep_hier will then map back to the original cell with parameters).
41
42 -prep_delays
43     insert `$_ABC9_DELAY' blackbox cells into the design to account for
44     certain required times.
45
46 -break_scc
47     for an arbitrarily chosen cell in each unique SCC of each selected module
48     (tagged with an (* abc9_scc_id = <int> *) attribute) interrupt all wires
49     driven by this cell's outputs with a temporary $_ABC9_SCC_BREAKER cell
50     to break the SCC.
51
52 -prep_xaiger
53     prepare the design for XAIGER output. this includes computing the
54     topological ordering of ABC9 boxes, as well as preparing the '$abc9_holes'
55     design that contains the logic behaviour of ABC9 whiteboxes.
56
57 -dff
58     consider flop cells (those instantiating modules marked with (* abc9_flop *))
59     during -prep_{delays,xaiger,box}.
60
61 -prep_lut <maxlut>
62     pre-compute the lut library by analysing all modules marked with
63     (* abc9_lut=<area> *).
64

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
65 -write_lut <dst>
66     write the pre-computed lut library to <dst>.
67
68 -prep_box
69     pre-compute the box library by analysing all modules marked with
70     (* abc9_box *).
71
72 -write_box <dst>
73     write the pre-computed box library to <dst>.
74
75 -reintegrate
76     for each selected module, re-intergrate the module '<module-name>$abc9'
77     by first recovering ABC9 boxes, and then stitching in the remaining primary
78     inputs and outputs.
```

### C.5 add – add objects to the design

```
1  add <command> [selection]
2
3  This command adds objects to the design. It operates on all fully selected
4  modules. So e.g. 'add -wire foo' will add a wire foo to all selected modules.
5
6
7  add {-wire|-input|-inout|-output} <name> <width> [selection]
8
9  Add a wire (input, inout, output port) with the given name and width. The
10 command will fail if the object exists already and has different properties
11 than the object to be created.
12
13
14 add -global_input <name> <width> [selection]
15
16 Like 'add -input', but also connect the signal between instances of the
17 selected modules.
18
19
20 add {-assert|-assume|-live|-fair|-cover} <name1> [-if <name2>]
21
22 Add an $assert, $assume, etc. cell connected to a wire named name1, with its
23 enable signal optionally connected to a wire named name2 (default: 1'b1).
24
25
26 add -mod <name[s]>
27
28 Add module[s] with the specified name[s].
```

### C.6 aigmap – map logic to and-inverter-graph circuit

```
1  aigmap [options] [selection]
2
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
3 Replace all logic cells with circuits made of only $AND_ and
4 $NOT_ cells.
5
6     -nand
7         Enable creation of $NAND_ cells
8
9     -select
10        Overwrite replaced cells in the current selection with new $AND_,
11        $NOT_, and $NAND_ cells
```

### C.7 alumacc – extract ALU and MACC cells

```
1     alumacc [selection]
2
3 This pass translates arithmetic operations like $add, $mul, $lt, etc. to $alu
4 and $macc cells.
```

### C.8 anlogic\_eqn – Anlogic: Calculate equations for luts

```
1     anlogic_eqn [selection]
2
3 Calculate equations for luts since bitstream generator depends on it.
```

### C.9 anlogic\_fixcarry – Anlogic: fix carry chain

```
1     anlogic_fixcarry [options] [selection]
2
3 Add Anlogic adders to fix carry chain if needed.
```

### C.10 assertpmux – adds asserts for parallel muxes

```
1     assertpmux [options] [selection]
2
3 This command adds asserts to the design that assert that all parallel muxes
4 ($pmux cells) have a maximum of one of their inputs enable at any time.
5
6     -noinit
7         do not enforce the pmux condition during the init state
8
9     -always
10        usually the $pmux condition is only checked when the $pmux output
11        is used by the mux tree it drives. this option will deactivate this
12        additional constraint and check the $pmux condition always.
```

## C.11 `async2sync` – convert async FF inputs to sync circuits

```

1  async2sync [options] [selection]
2
3  This command replaces async FF inputs with sync circuits emulating the same
4  behavior for when the async signals are actually synchronized to the clock.
5
6  This pass assumes negative hold time for the async FF inputs. For example when
7  a reset deasserts with the clock edge, then the FF output will still drive the
8  reset value in the next cycle regardless of the data-in value at the time of
9  the clock edge.

```

## C.12 `attrmap` – renaming attributes

```

1  attrmap [options] [selection]
2
3  This command renames attributes and/or maps key/value pairs to
4  other key/value pairs.
5
6  -tocase <name>
7      Match attribute names case-insensitively and set it to the specified
8      name.
9
10 -rename <old_name> <new_name>
11     Rename attributes as specified
12
13 -map <old_name>=<old_value> <new_name>=<new_value>
14     Map key/value pairs as indicated.
15
16 -imap <old_name>=<old_value> <new_name>=<new_value>
17     Like -map, but use case-insensitive match for <old_value> when
18     it is a string value.
19
20 -remove <name>=<value>
21     Remove attributes matching this pattern.
22
23 -modattr
24     Operate on module attributes instead of attributes on wires and cells.
25
26 For example, mapping Xilinx-style "keep" attributes to Yosys-style:
27
28 attrmap -tocase keep -imap keep="true" keep=1 \
29         -imap keep="false" keep=0 -remove keep=0

```

## C.13 `attrmvcp` – move or copy attributes from wires to driving cells

```

1  attrmvcp [options] [selection]
2

```



## APPENDIX C. COMMAND REFERENCE MANUAL

```
3 Move or copy attributes on wires to the cells driving them.
4
5     -copy
6         By default, attributes are moved. This will only add
7         the attribute to the cell, without removing it from
8         the wire.
9
10    -purge
11        If no selected cell consumes the attribute, then it is
12        left on the wire by default. This option will cause the
13        attribute to be removed from the wire, even if no selected
14        cell takes it.
15
16    -driven
17        By default, attributes are moved to the cell driving the
18        wire. With this option set it will be moved to the cell
19        driven by the wire instead.
20
21    -attr <attrname>
22        Move or copy this attribute. This option can be used
23        multiple times.
```

### C.14 autaname – automatically assign names to objects

```
1     autaname [selection]
2
3 Assign auto-generated public names to objects with private names (the ones
4 with $-prefix).
```

### C.15 blackbox – convert modules into blackbox modules

```
1     blackbox [options] [selection]
2
3 Convert modules into blackbox modules (remove contents and set the blackbox
4 module attribute).
```

### C.16 bmuxmap – transform \$bmux cells to trees of \$mux cells

```
1     bmuxmap [selection]
2
3 This pass transforms $bmux cells to trees of $mux cells.
```

### C.17 bugpoint – minimize testcases

## APPENDIX C. COMMAND REFERENCE MANUAL

```
1  bugpoint [options] [-script <filename> | -command "<command>"]
2
3  This command minimizes the current design that is known to crash Yosys with the
4  given script into a smaller testcase. It does this by removing an arbitrary part
5  of the design and recursively invokes a new Yosys process with this modified design
6  and the same script, repeating these steps while it can find a smaller design that
7  still causes a crash. Once this command finishes, it replaces the current design
8  with the smallest testcase it was able to produce.
9  In order to save the reduced testcase you must write this out to a file with
10 another command after `bugpoint` like `write_rtlil` or `write_verilog`.
11
12  -script <filename> | -command "<command>"
13      use this script file or command to crash Yosys. required.
14
15  -yosys <filename>
16      use this Yosys binary. if not specified, `yosys` is used.
17
18  -grep "<string>"
19      only consider crashes that place this string in the log file.
20
21  -fast
22      run `proc_clean; clean -purge` after each minimization step. converges
23      faster, but produces larger testcases, and may fail to produce any
24      testcase at all if the crash is related to dangling wires.
25
26  -clean
27      run `proc_clean; clean -purge` before checking testcase and after
28      finishing. produces smaller and more useful testcases, but may fail to
29      produce any testcase at all if the crash is related to dangling wires.
30
31  It is possible to constrain which parts of the design will be considered for
32  removal. Unless one or more of the following options are specified, all parts
33  will be considered.
34
35  -modules
36      try to remove modules. modules with a (* bugpoint_keep *) attribute
37      will be skipped.
38
39  -ports
40      try to remove module ports. ports with a (* bugpoint_keep *) attribute
41      will be skipped (useful for clocks, resets, etc.)
42
43  -cells
44      try to remove cells. cells with a (* bugpoint_keep *) attribute will
45      be skipped.
46
47  -connections
48      try to reconnect ports to 'x'.
49
50  -processes
51      try to remove processes. processes with a (* bugpoint_keep *) attribute
52      will be skipped.
53
54  -assigns
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
55         try to remove process assigns from cases.
56
57     -updates
58         try to remove process updates from syncs.
59
60     -runner "<prefix>"
61         child process wrapping command, e.g., "timeout 30", or valgrind.
```

### C.18 cd – a shortcut for 'select -module <name>'

```
1         cd <modname>
2
3     This is just a shortcut for 'select -module <modname>'.
4
5
6         cd <cellname>
7
8     When no module with the specified name is found, but there is a cell
9     with the specified name in the current module, then this is equivalent
10    to 'cd <celltype>'.
11
12        cd ..
13
14    Remove trailing substrings that start with '.' in current module name until
15    the name of a module in the current design is generated, then switch to that
16    module. Otherwise clear the current selection.
17
18        cd
19
20    This is just a shortcut for 'select -clear'.
```

### C.19 check – check for obvious problems in the design

```
1         check [options] [selection]
2
3     This pass identifies the following problems in the current design:
4
5     - combinatorial loops
6     - two or more conflicting drivers for one wire
7     - used wires that do not have a driver
8
9     Options:
10
11     -noinit
12         also check for wires which have the 'init' attribute set
13
14     -initdrv
15         also check for wires that have the 'init' attribute set and are not
16         driven by an FF cell type
17
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
18  -mapped
19      also check for internal cells that have not been mapped to cells of the
20      target architecture
21
22  -allow-tbuf
23      modify the -mapped behavior to still allow $_TBUF_ cells
24
25  -assert
26      produce a runtime error if any problems are found in the current design
```

### C.20 chformal – change formal constraints of the design

```
1      chformal [types] [mode] [options] [selection]
2
3  Make changes to the formal constraints of the design. The [types] options
4  the type of constraint to operate on. If none of the following options are given,
5  the command will operate on all constraint types:
6
7  -assert      $assert cells, representing assert(...) constraints
8  -assume      $assume cells, representing assume(...) constraints
9  -live        $live cells, representing assert(s_eventually ...)
10 -fair         $fair cells, representing assume(s_eventually ...)
11 -cover        $cover cells, representing cover() statements
12
13 Exactly one of the following modes must be specified:
14
15 -remove
16     remove the cells and thus constraints from the design
17
18 -early
19     bypass FFs that only delay the activation of a constraint
20
21 -delay <N>
22     delay activation of the constraint by <N> clock cycles
23
24 -skip <N>
25     ignore activation of the constraint in the first <N> clock cycles
26
27 -assert2assume
28 -assume2assert
29 -live2fair
30 -fair2live
31     change the roles of cells as indicated. these options can be combined
```

### C.21 chparam – re-evaluate modules with new parameters

```
1      chparam [ -set name value ]... [selection]
2
3  Re-evaluate the selected modules with new parameters. String values must be
4  passed in double quotes (").
```

```

5
6
7     chparam -list [selection]
8
9 List the available parameters of the selected modules.

```

## C.22 chtype – change type of cells in the design

```

1     chtype [options] [selection]
2
3 Change the types of cells in the design.
4
5     -set <type>
6         set the cell type to the given type
7
8     -map <old_type> <new_type>
9         change cells types that match <old_type> to <new_type>

```

## C.23 clean – remove unused cells and wires

```

1     clean [options] [selection]
2
3 This is identical to 'opt_clean', but less verbose.
4
5 When commands are separated using the ';' token, this command will be executed
6 between the commands.
7
8 When commands are separated using the ';;' token, this command will be executed
9 in -purge mode between the commands.

```

## C.24 clean\_zerowidth – clean zero-width connections from the design

```

1     clean_zerowidth [selection]
2
3 Fixes the selected cells and processes to contain no zero-width connections.
4 Depending on the cell type, this may be implemented by removing the connection,
5 widening it to 1-bit, or removing the cell altogether.

```

## C.25 clk2fflogic – convert clocked FFs to generic \$ff cells

```

1     clk2fflogic [options] [selection]
2
3 This command replaces clocked flip-flops with generic $ff cells that use the

```

```

4 implicit global clock. This is useful for formal verification of designs with
5 multiple clocks.

```

## C.26 clkbufmap – insert clock buffers on clock networks

```

1      clkbufmap [options] [selection]
2
3 Inserts clock buffers between nets connected to clock inputs and their drivers.
4
5 In the absence of any selection, all wires without the 'clkbuf_inhibit'
6 attribute will be considered for clock buffer insertion.
7 Alternatively, to consider all wires without the 'buffer_type' attribute set to
8 'none' or 'bufr' one would specify:
9     'w:* a:buffer_type=none a:buffer_type=bufr %u %d'
10 as the selection.
11
12     -buf <celltype> <portname_out>:<portname_in>
13         Specifies the cell type to use for the clock buffers
14         and its port names. The first port will be connected to
15         the clock network sinks, and the second will be connected
16         to the actual clock source.
17
18     -inpad <celltype> <portname_out>:<portname_in>
19         If specified, a PAD cell of the given type is inserted on
20         clock nets that are also top module's inputs (in addition
21         to the clock buffer, if any).
22
23 At least one of -buf or -inpad should be specified.

```

## C.27 connect – create or remove connections

```

1      connect [-nomap] [-nounset] -set <lhs-expr> <rhs-expr>
2
3 Create a connection. This is equivalent to adding the statement 'assign
4 <lhs-expr> = <rhs-expr>;' to the Verilog input. Per default, all existing
5 drivers for <lhs-expr> are unconnected. This can be overwritten by using
6 the -nounset option.
7
8
9      connect [-nomap] -unset <expr>
10
11 Unconnect all existing drivers for the specified expression.
12
13
14      connect [-nomap] [-assert] -port <cell> <port> <expr>
15
16 Connect the specified cell port to the specified cell port.
17
18
19 Per default signal alias names are resolved and all signal names are mapped

```

## APPENDIX C. COMMAND REFERENCE MANUAL

the the signal name of the primary driver. Using the `-nomap` option deactivates this behavior.

The `connect` command operates in one module only. Either only one module must be selected or an active module must be set using the `'cd'` command.

The `-assert` option verifies that the connection already exists, instead of making it.

This command does not operate on module with processes.

### C.28 `connect_rpc` – connect to RPC frontend

```
1 connect_rpc -exec <command> [args...]
2 connect_rpc -path <path>
3
4 Load modules using an out-of-process frontend.
5
6 -exec <command> [args...]
7     run <command> with arguments [args...]. send requests on stdin, read
8     responses from stdout.
9
10 -path <path>
11     connect to Unix domain socket at <path>. (Unix)
12     connect to bidirectional byte-type named pipe at <path>. (Windows)
13
14 A simple JSON-based, newline-delimited protocol is used for communicating with
15 the frontend. Yosys requests data from the frontend by sending exactly 1 line
16 of JSON. Frontend responds with data or error message by replying with exactly
17 1 line of JSON as well.
18
19 -> {"method": "modules"}
20 <- {"modules": ["<module-name>", ...]}
21 <- {"error": "<error-message>"}
22     request for the list of modules that can be derived by this frontend.
23     the 'hierarchy' command will call back into this frontend if a cell
24     with type <module-name> is instantiated in the design.
25
26 -> {"method": "derive", "module": "<module-name>", "parameters": {
27     "<param-name>": {"type": "[unsigned|signed|string|real]",
28         "value": "<param-value>"}, ...}}
29 <- {"frontend": "[rtlil|verilog|...]", "source": "<source>"}
30 <- {"error": "<error-message>"}
31     request for the module <module-name> to be derived for a specific set of
32     parameters. <param-name> starts with \ for named parameters, and with $
33     for unnamed parameters, which are numbered starting at 1.<param-value>
34     for integer parameters is always specified as a binary string of unlimited
35     precision. the <source> returned by the frontend is hygienically parsed
36     by a built-in Yosys <frontend>, allowing the RPC frontend to return any
37     convenient representation of the module. the derived module is cached,
38     so the response should be the same whenever the same set of parameters
39     is provided.
```

**C.29 connwrappers – match width of input-output port pairs**

```

1      connwrappers [options] [selection]
2
3  Wrappers are used in coarse-grain synthesis to wrap cells with smaller ports
4  in wrapper cells with a (larger) constant port size. I.e. the upper bits
5  of the wrapper output are signed/unsigned bit extended. This command uses this
6  knowledge to rewire the inputs of the driven cells to match the output of
7  the driving cell.
8
9      -signed <cell_type> <port_name> <width_param>
10     -unsigned <cell_type> <port_name> <width_param>
11         consider the specified signed/unsigned wrapper output
12
13     -port <cell_type> <port_name> <width_param> <sign_param>
14         use the specified parameter to decide if signed or unsigned
15
16 The options -signed, -unsigned, and -port can be specified multiple times.

```

**C.30 coolrunner2\_fixup – insert necessary buffer cells for CoolRunner-II architecture**

```

1      coolrunner2_fixup [options] [selection]
2
3  Insert necessary buffer cells for CoolRunner-II architecture.

```

**C.31 coolrunner2\_sop – break \$sop cells into ANDTERM/ORTERM cells**

```

1      coolrunner2_sop [options] [selection]
2
3  Break $sop cells into ANDTERM/ORTERM cells.

```

**C.32 copy – copy modules in the design**

```

1      copy old_name new_name
2
3  Copy the specified module. Note that selection patterns are not supported
4  by this command.

```

**C.33 cover – print code coverage counters**



## APPENDIX C. COMMAND REFERENCE MANUAL

```
1  cover [options] [pattern]
2
3  Print the code coverage counters collected using the cover() macro in the Yosys
4  C++ code. This is useful to figure out what parts of Yosys are utilized by a
5  test bench.
6
7  -q
8      Do not print output to the normal destination (console and/or log file)
9
10 -o file
11     Write output to this file, truncate if exists.
12
13 -a file
14     Write output to this file, append if exists.
15
16 -d dir
17     Write output to a newly created file in the specified directory.
18
19 When one or more pattern (shell wildcards) are specified, then only counters
20 matching at least one pattern are printed.
21
22
23 It is also possible to instruct Yosys to print the coverage counters on program
24 exit to a file using environment variables:
25
26 YOSYS_COVER_DIR="{dir-name}" yosys {args}
27
28     This will create a file (with an auto-generated name) in this
29     directory and write the coverage counters to it.
30
31 YOSYS_COVER_FILE="{file-name}" yosys {args}
32
33     This will append the coverage counters to the specified file.
34
35
36 Hint: Use the following AWK command to consolidate Yosys coverage files:
37
38 gawk '{ p[$3] = $1; c[$3] += $2; } END { for (i in p)
39     printf "%-60s %10d %s\n", p[i], c[i], i; }' {files} | sort -k3
40
41
42 Coverage counters are only available in Yosys for Linux.
```

### C.34 cutpoint – adds formal cut points to the design

```
1  cutpoint [options] [selection]
2
3  This command adds formal cut points to the design.
4
5  -undef
6      set cupoint nets to undef (x). the default behavior is to create a
```

```
7 | $anyseq cell and drive the cutpoint net from that
```

### C.35 debug – run command with debug log messages enabled

```
1 | debug cmd
2 |
3 | Execute the specified command with debug log messages enabled
```

### C.36 delete – delete objects in the design

```
1 | delete [selection]
2 |
3 | Deletes the selected objects. This will also remove entire modules, if the
4 | whole module is selected.
5 |
6 |
7 | delete {-input|-output|-port} [selection]
8 |
9 | Does not delete any object but removes the input and/or output flag on the
10 | selected wires, thus 'deleting' module ports.
```

### C.37 deminout – demote inout ports to input or output

```
1 | deminout [options] [selection]
2 |
3 | "Demote" inout ports to input or output ports, if possible.
```

### C.38 demuxmap – transform \$demux cells to \$eq + \$mux cells

```
1 | demuxmap [selection]
2 |
3 | This pass transforms $demux cells to a bunch of equality comparisons.
```

### C.39 design – save, restore and reset current design

```
1 | design -reset
2 |
3 | Clear the current design.
4 |
5 |
6 | design -save <name>
7 |
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
8 | Save the current design under the given name.
9 |
10 |
11 |     design -stash <name>
12 |
13 | Save the current design under the given name and then clear the current design.
14 |
15 |
16 |     design -push
17 |
18 | Push the current design to the stack and then clear the current design.
19 |
20 |
21 |     design -push-copy
22 |
23 | Push the current design to the stack without clearing the current design.
24 |
25 |
26 |     design -pop
27 |
28 | Reset the current design and pop the last design from the stack.
29 |
30 |
31 |     design -load <name>
32 |
33 | Reset the current design and load the design previously saved under the given
34 | name.
35 |
36 |
37 |     design -copy-from <name> [-as <new_mod_name>] <selection>
38 |
39 | Copy modules from the specified design into the current one. The selection is
40 | evaluated in the other design.
41 |
42 |
43 |     design -copy-to <name> [-as <new_mod_name>] [selection]
44 |
45 | Copy modules from the current design into the specified one.
46 |
47 |
48 |     design -import <name> [-as <new_top_name>] [selection]
49 |
50 | Import the specified design into the current design. The source design must
51 | either have a selected top module or the selection must contain exactly one
52 | module that is then used as top module for this command.
53 |
54 |
55 |     design -reset-vlog
56 |
57 | The Verilog front-end remembers defined macros and top-level declarations
58 | between calls to 'read_verilog'. This command resets this memory.
59 |
60 |     design -delete <name>
61 |
```

62 Delete the design previously saved under the given name.

## C.40 dffinit – set INIT param on FF cells

```

1  dffinit [options] [selection]
2
3  This pass sets an FF cell parameter to the the initial value of the net it
4  drives. (This is primarily used in FPGA flows.)
5
6  -ff <cell_name> <output_port> <init_param>
7      operate on the specified cell type. this option can be used
8      multiple times.
9
10 -highlow
11     use the string values "high" and "low" to represent a single-bit
12     initial value of 1 or 0. (multi-bit values are not supported in this
13     mode.)
14
15 -strinit <string for high> <string for low>
16     use string values in the command line to represent a single-bit
17     initial value of 1 or 0. (multi-bit values are not supported in this
18     mode.)
19
20 -noreinit
21     fail if the FF cell has already a defined initial value set in other
22     passes and the initial value of the net it drives is not equal to
23     the already defined initial value.
```

## C.41 dfflegalize – convert FFs to types supported by the target

```

1  dfflegalize [options] [selection]
2
3  Converts FFs to types supported by the target.
4
5  -cell <cell_type_pattern> <init_values>
6      specifies a supported group of FF cells. <cell_type_pattern>
7      is a yosys internal fine cell name, where ? characters can be
8      as a wildcard matching any character. <init_values> specifies
9      which initialization values these FF cells can support, and can
10     be one of:
11
12     - x (no init value supported)
13     - 0
14     - 1
15     - r (init value has to match reset value, only for some FF types)
16     - 01 (both 0 and 1 supported).
17
18 -mince <num>
19     specifies a minimum number of FFs that should be using any given
20     clock enable signal. If a clock enable signal doesn't meet this
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

21     threshold, it is unmapped into soft logic.
22
23     -minsrst <num>
24         specifies a minimum number of FFs that should be using any given
25         sync set/reset signal. If a sync set/reset signal doesn't meet this
26         threshold, it is unmapped into soft logic.
27
28 The following cells are supported by this pass (ie. will be ingested,
29 and can be specified as allowed targets):
30
31 - $_DFF_[NP]_
32 - $_DFFE_[NP][NP]_
33 - $_DFF_[NP][NP][01]_
34 - $_DFFE_[NP][NP][01][NP]_
35 - $_ALDFF_[NP][NP]_
36 - $_ALDFFE_[NP][NP][NP]_
37 - $_DFFSR_[NP][NP][NP]_
38 - $_DFFSRE_[NP][NP][NP][NP]_
39 - $_SDFF_[NP][NP][01]_
40 - $_SDFFE_[NP][NP][01][NP]_
41 - $_SDFFCE_[NP][NP][01][NP]_
42 - $_SR_[NP][NP]_
43 - $_DLATCH_[NP]_
44 - $_DLATCH_[NP][NP][01]_
45 - $_DLATCHSR_[NP][NP][NP]_
46
47 The following transformations are performed by this pass:
48 - upconversion from a less capable cell to a more capable cell, if the less capable cell is not supported
49 - unmapping FFs with clock enable (due to unsupported cell type or -mince)
50 - unmapping FFs with sync reset (due to unsupported cell type or -minsrst)
51 - adding inverters on the control pins (due to unsupported polarity)
52 - adding inverters on the D and Q pins and inverting the init/reset values
53   (due to unsupported init or reset value)
54 - converting sr into adlatch (by tying D to 1 and using E as set input)
55 - emulating unsupported dffsr cell by adff + adff + sr + mux
56 - emulating unsupported dlatchsr cell by adlatch + adlatch + sr + mux
57 - emulating adff when the (reset, init) value combination is unsupported by
58   dff + adff + dlatch + mux
59 - emulating adlatch when the (reset, init) value combination is unsupported by
60   dlatch + adlatch + dlatch + mux
61 If the pass is unable to realize a given cell type (eg. adff when only plain dffis available), an error is

```

### C.42 dfflibmap – technology mapping of flip-flops

```

1     dfflibmap [-prepare] [-map-only] [-info] -liberty <file> [selection]
2
3 Map internal flip-flop cells to the flip-flop cells in the technology
4 library specified in the given liberty file.
5
6 This pass may add inverters as needed. Therefore it is recommended to
7 first run this pass and then map the logic paths to the target technology.
8

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
9 | When called with -prepare, this command will convert the internal FF cells
10 | to the internal cell types that best match the cells found in the given
11 | liberty file, but won't actually map them to the target cells.
12 |
13 | When called with -map-only, this command will only map internal cell
14 | types that are already of exactly the right type to match the target
15 | cells, leaving remaining internal cells untouched.
16 |
17 | When called with -info, this command will only print the target cell
18 | list, along with their associated internal cell types, and the arguments that would be passed to the dfflega
19 | The design will not be
    changed.
```

### C.43 dffunmap – unmap clock enable and synchronous reset from FFs

```
1 | dffunmap [options] [selection]
2 |
3 | This pass transforms FF types with clock enable and/or synchronous reset into
4 | their base type (with neither clock enable nor sync reset) by emulating the clock
5 | enable and synchronous reset with multiplexers on the cell input.
6 |
7 | -ce-only
8 |     unmap only clock enables, leave synchronous resets alone.
9 |
10 | -srst-only
11 |     unmap only synchronous resets, leave clock enables alone.
```

### C.44 dump – print parts of the design in RTLIL format

```
1 | dump [options] [selection]
2 |
3 | Write the selected parts of the design to the console or specified file in
4 | RTLIL format.
5 |
6 | -m
7 |     also dump the module headers, even if only parts of a single
8 |     module is selected
9 |
10 | -n
11 |     only dump the module headers if the entire module is selected
12 |
13 | -o <filename>
14 |     write to the specified file.
15 |
16 | -a <filename>
17 |     like -outfile but append instead of overwrite
```

## C.45 echo – turning echoing back of commands on and off

```

1      echo on
2
3  Print all commands to log before executing them.
4
5
6      echo off
7
8  Do not print all commands to log before executing them. (default)

```

## C.46 ecp5\_gsr – ECP5: handle GSR

```

1      ecp5_gsr [options] [selection]
2
3  Trim active low async resets connected to GSR and resolve GSR parameter,
4  if a GSR or SGR primitive is used in the design.
5
6  If any cell has the GSR parameter set to "AUTO", this will be resolved
7  to "ENABLED" if a GSR primitive is present and the (* nogsr *) attribute
8  is not set, otherwise it will be resolved to "DISABLED".

```

## C.47 edgetypes – list all types of edges in selection

```

1      edgetypes [options] [selection]
2
3  This command lists all unique types of 'edges' found in the selection. An 'edge'
4  is a 4-tuple of source and sink cell type and port name.

```

## C.48 efinix\_fixcarry – Efinix: fix carry chain

```

1      efinix_fixcarry [options] [selection]
2
3  Add Efinix adders to fix carry chain if needed.

```

## C.49 equiv\_add – add a \$equiv cell

```

1      equiv_add [-try] gold_sig gate_sig
2
3  This command adds an $equiv cell for the specified signals.
4
5
6      equiv_add [-try] -cell gold_cell gate_cell
7
8  This command adds $equiv cells for the ports of the specified cells.

```

## C.50 equiv\_induct – proving \$equiv cells using temporal induction

```

1  equiv_induct [options] [selection]
2
3  Uses a version of temporal induction to prove $equiv cells.
4
5  Only selected $equiv cells are proven and only selected cells are used to
6  perform the proof.
7
8  -undef
9      enable modelling of undef states
10
11  -seq <N>
12      the max. number of time steps to be considered (default = 4)
13
14  This command is very effective in proving complex sequential circuits, when
15  the internal state of the circuit quickly propagates to $equiv cells.
16
17  However, this command uses a weak definition of 'equivalence': This command
18  proves that the two circuits will not diverge after they produce equal
19  outputs (observable points via $equiv) for at least <N> cycles (the <N>
20  specified via -seq).
21
22  Combined with simulation this is very powerful because simulation can give
23  you confidence that the circuits start out synced for at least <N> cycles
24  after reset.
```

## C.51 equiv\_make – prepare a circuit for equivalence checking

```

1  equiv_make [options] gold_module gate_module equiv_module
2
3  This creates a module annotated with $equiv cells from two presumably
4  equivalent modules. Use commands such as 'equiv_simple' and 'equiv_status'
5  to work with the created equivalent checking module.
6
7  -inames
8      Also match cells and wires with $... names.
9
10  -blacklist <file>
11      Do not match cells or signals that match the names in the file.
12
13  -encfile <file>
14      Match FSM encodings using the description from the file.
15      See 'help fsm_recode' for details.
16
17  Note: The circuit created by this command is not a miter (with something like
18  a trigger output), but instead uses $equiv cells to encode the equivalence
19  checking problem. Use 'miter -equiv' if you want to create a miter circuit.
```



## C.52 equiv\_mark – mark equivalence checking regions

```

1  equiv_mark [options] [selection]
2
3  This command marks the regions in an equivalence checking module. Region 0 is
4  the proven part of the circuit. Regions with higher numbers are connected
5  unproven subcircuits. The integer attribute 'equiv_region' is set on all
6  wires and cells.

```

## C.53 equiv\_miter – extract miter from equiv circuit

```

1  equiv_miter [options] miter_module [selection]
2
3  This creates a miter module for further analysis of the selected $equiv cells.
4
5  -trigger
6      Create a trigger output
7
8  -cmp
9      Create cmp_* outputs for individual unproven $equiv cells
10
11 -assert
12     Create a $assert cell for each unproven $equiv cell
13
14 -undef
15     Create compare logic that handles undefs correctly

```

## C.54 equiv\_opt – prove equivalence for optimized circuit

```

1  equiv_opt [options] [command]
2
3  This command uses temporal induction to check circuit equivalence before and
4  after an optimization pass.
5
6  -run <from_label>:<to_label>
7      only run the commands between the labels (see below). an empty
8      from label is synonymous to the start of the command list, and empty to
9      label is synonymous to the end of the command list.
10
11 -map <filename>
12     expand the modules in this file before proving equivalence. this is
13     useful for handling architecture-specific primitives.
14
15 -blacklist <file>
16     Do not match cells or signals that match the names in the file
17     (passed to equiv_make).
18
19 -assert
20     produce an error if the circuits are not equivalent.

```

```

21
22     -multiclock
23         run clk2fflogic before equivalence checking.
24
25     -async2sync
26         run async2sync before equivalence checking.
27
28     -undef
29         enable modelling of undef states during equiv_induct.
30
31 The following commands are executed by this verification command:
32
33     run_pass:
34         hierarchy -auto-top
35         design -save preopt
36         [command]
37         design -stash postopt
38
39     prepare:
40         design -copy-from preopt  -as gold A:top
41         design -copy-from postopt -as gate A:top
42
43     techmap:    (only with -map)
44         techmap -wb -D EQUIV -autoproc -map <filename> ...
45
46     prove:
47         clk2fflogic    (only with -multiclock)
48         async2sync    (only with -async2sync)
49         equiv_make -blacklist <filename> ... gold gate equiv
50         equiv_induct [-undef] equiv
51         equiv_status [-assert] equiv
52
53     restore:
54         design -load preopt

```

## C.55 equiv\_purge – purge equivalence checking module

```

1     equiv_purge [options] [selection]
2
3 This command removes the proven part of an equivalence checking module, leaving
4 only the unproven segments in the design. This will also remove and add module
5 ports as needed.

```

## C.56 equiv\_remove – remove \$equiv cells

```

1     equiv_remove [options] [selection]
2
3 This command removes the selected $equiv cells. If neither -gold nor -gate is
4 used then only proven cells are removed.
5

```

```

6      -gold
7          keep gold circuit
8
9      -gate
10         keep gate circuit

```

## C.57 equiv\_simple – try proving simple \$equiv instances

```

1      equiv_simple [options] [selection]
2
3      This command tries to prove $equiv cells using a simple direct SAT approach.
4
5      -v
6          verbose output
7
8      -undef
9          enable modelling of undef states
10
11     -short
12         create shorter input cones that stop at shared nodes. This yields
13         simpler SAT problems but sometimes fails to prove equivalence.
14
15     -nogroup
16         disabling grouping of $equiv cells by output wire
17
18     -seq <N>
19         the max. number of time steps to be considered (default = 1)

```

## C.58 equiv\_status – print status of equivalent checking module

```

1      equiv_status [options] [selection]
2
3      This command prints status information for all selected $equiv cells.
4
5      -assert
6          produce an error if any unproven $equiv cell is found

```

## C.59 equiv\_struct – structural equivalence checking

```

1      equiv_struct [options] [selection]
2
3      This command adds additional $equiv cells based on the assumption that the
4      gold and gate circuit are structurally equivalent. Note that this can introduce
5      bad $equiv cells in cases where the netlists are not structurally equivalent,
6      for example when analyzing circuits with cells with commutative inputs. This
7      command will also de-duplicate gates.
8

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
9      -fwd
10         by default this command performans forward sweeps until nothing can
11         be merged by forwards sweeps, then backward sweeps until forward
12         sweeps are effective again. with this option set only forward sweeps
13         are performed.
14
15      -fwnonly <cell_type>
16         add the specified cell type to the list of cell types that are only
17         merged in forward sweeps and never in backward sweeps. $equiv is in
18         this list automatically.
19
20      -icells
21         by default, the internal RTL and gate cell types are ignored. add
22         this option to also process those cell types with this command.
23
24      -maxiter <N>
25         maximum number of iterations to run before aborting
```

### C.60 eval – evaluate the circuit given an input

```
1      eval [options] [selection]
2
3      This command evaluates the value of a signal given the value of all required
4      inputs.
5
6      -set <signal> <value>
7         set the specified signal to the specified value.
8
9      -set-undef
10         set all unspecified source signals to undef (x)
11
12      -table <signal>
13         create a truth table using the specified input signals
14
15      -show <signal>
16         show the value for the specified signal. if no -show option is passed
17         then all output ports of the current module are used.
```

### C.61 exec – execute commands in the operating system shell

```
1      exec [options] -- [command]
2
3      Execute a command in the operating system shell. All supplied arguments are
4      concatenated and passed as a command to popen(3). Whitespace is not guaranteed
5      to be preserved, even if quoted. stdin and stderr are not connected, while stdout is
6      logged unless the "-q" option is specified.
7
8
9      -q
10         Suppress stdout and stderr from subprocess
```

```

11
12 -expect-return <int>
13     Generate an error if popen() does not return specified value.
14     May only be specified once; the final specified value is controlling
15     if specified multiple times.
16
17 -expect-stdout <regex>
18     Generate an error if the specified regex does not match any line
19     in subprocess's stdout. May be specified multiple times.
20
21 -not-expect-stdout <regex>
22     Generate an error if the specified regex matches any line
23     in subprocess's stdout. May be specified multiple times.
24
25
26 Example: exec -q -expect-return 0 -- echo "bananapie" | grep "nana"

```

## C.62 expose – convert internal signals to module ports

```

1  expose [options] [selection]
2
3  This command exposes all selected internal signals of a module as additional
4  outputs.
5
6  -dff
7      only consider wires that are directly driven by register cell.
8
9  -cut
10     when exposing a wire, create an input/output pair and cut the internal
11     signal path at that wire.
12
13  -input
14     when exposing a wire, create an input port and disconnect the internal
15     driver.
16
17  -shared
18     only expose those signals that are shared among the selected modules.
19     this is useful for preparing modules for equivalence checking.
20
21  -evert
22     also turn connections to instances of other modules to additional
23     inputs and outputs and remove the module instances.
24
25  -evert-dff
26     turn flip-flops to sets of inputs and outputs.
27
28  -sep <separator>
29     when creating new wire/port names, the original object name is suffixed
30     with this separator (default: '.') and the port name or a type
31     designator for the exposed signal.

```

**C.63 extract – find subcircuits and replace them with cells**

```

1  extract -map <map_file> [options] [selection]
2  extract -mine <out_file> [options] [selection]
3
4  This pass looks for subcircuits that are isomorphic to any of the modules
5  in the given map file and replaces them with instances of this modules. The
6  map file can be a Verilog source file (*.v) or an RTLIL source file (*.il).
7
8  -map <map_file>
9      use the modules in this file as reference. This option can be used
10     multiple times.
11
12  -map %<design-name>
13      use the modules in this in-memory design as reference. This option can
14      be used multiple times.
15
16  -verbose
17      print debug output while analyzing
18
19  -constports
20      also find instances with constant drivers. this may be much
21      slower than the normal operation.
22
23  -nodefaultswaps
24      normally builtin port swapping rules for internal cells are used per
25      default. This turns that off, so e.g. 'a^b' does not match 'b^a'
26      when this option is used.
27
28  -compat <needle_type> <haystack_type>
29      Per default, the cells in the map file (needle) must have the
30      type as the cells in the active design (haystack). This option
31      can be used to register additional pairs of types that should
32      match. This option can be used multiple times.
33
34  -swap <needle_type> <port1>,<port2>[,...]
35      Register a set of swappable ports for a needle cell type.
36      This option can be used multiple times.
37
38  -perm <needle_type> <port1>,<port2>[,...] <portA>,<portB>[,...]
39      Register a valid permutation of swappable ports for a needle
40      cell type. This option can be used multiple times.
41
42  -cell_attr <attribute_name>
43      Attributes on cells with the given name must match.
44
45  -wire_attr <attribute_name>
46      Attributes on wires with the given name must match.
47
48  -ignore_parameters
49      Do not use parameters when matching cells.
50
51  -ignore_param <cell_type> <parameter_name>
52      Do not use this parameter when matching cells.

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
53
54 This pass does not operate on modules with unprocessed processes in it.
55 (I.e. the 'proc' pass should be used first to convert processes to netlists.)
56
57 This pass can also be used for mining for frequent subcircuits. In this mode
58 the following options are to be used instead of the -map option.
59
60 -mine <out_file>
61     mine for frequent subcircuits and write them to the given RTLIL file
62
63 -mine_cells_span <min> <max>
64     only mine for subcircuits with the specified number of cells
65     default value: 3 5
66
67 -mine_min_freq <num>
68     only mine for subcircuits with at least the specified number of matches
69     default value: 10
70
71 -mine_limit_matches_per_module <num>
72     when calculating the number of matches for a subcircuit, don't count
73     more than the specified number of matches per module
74
75 -mine_max_fanout <num>
76     don't consider internal signals with more than <num> connections
77
78 The modules in the map file may have the attribute 'extract_order' set to an
79 integer value. Then this value is used to determine the order in which the pass
80 tries to map the modules to the design (ascending, default value is 0).
81
82 See 'help techmap' for a pass that does the opposite thing.
```

### C.64 extract\_counter – Extract GreenPak4 counter cells

```
1     extract_counter [options] [selection]
2
3 This pass converts non-resettable or async resettable down counters to
4 counter cells. Use a target-specific 'techmap' map file to convert those cells
5 to the actual target cells.
6
7 -maxwidth N
8     Only extract counters up to N bits wide (default 64)
9
10 -minwidth N
11     Only extract counters at least N bits wide (default 2)
12
13 -allow_arst yes|no
14     Allow counters to have async reset (default yes)
15
16 -dir up|down|both
17     Look for up-counters, down-counters, or both (default down)
18
19 -pout X,Y,...
```

```

20     Only allow parallel output from the counter to the listed cell types
21     (if not specified, parallel outputs are not restricted)

```

## C.65 `extract_fa` – find and extract full/half adders

```

1     extract_fa [options] [selection]
2
3     This pass extracts full/half adders from a gate-level design.
4
5     -fa, -ha
6         Enable cell types (fa=full adder, ha=half adder)
7         All types are enabled if none of this options is used
8
9     -d <int>
10        Set maximum depth for extracted logic cones (default=20)
11
12     -b <int>
13        Set maximum breadth for extracted logic cones (default=6)
14
15     -v
16        Verbose output

```

## C.66 `extract_reduce` – converts gate chains into `$reduce_*` cells

```

1     extract_reduce [options] [selection]
2
3     converts gate chains into $reduce_* cells
4
5     This command finds chains of $_AND_, $_OR_, and $_XOR_ cells and replaces them
6     with their corresponding $reduce_* cells. Because this command only operates on
7     these cell types, it is recommended to map the design to only these cell types
8     using the `abc -g` command. Note that, in some cases, it may be more effective
9     to map the design to only $_AND_ cells, run extract_reduce, map the remaining
10    parts of the design to AND/OR/XOR cells, and run extract_reduce a second time.
11
12    -allow-off-chain
13        Allows matching of cells that have loads outside the chain. These cells
14        will be replicated and folded into the $reduce_* cell, but the original
15        cell will remain, driving its original loads.

```

## C.67 `extractinv` – extract explicit inverter cells for invertible cell pins

```

1     extractinv [options] [selection]
2
3     Searches the design for all cells with invertible pins controlled by a cell
4     parameter (eg. IS_CLK_INVERTED on many Xilinx cells) and removes the parameter.

```



## APPENDIX C. COMMAND REFERENCE MANUAL

```
5 If the parameter was set to 1, inserts an explicit inverter cell in front of
6 the pin instead. Normally used for output to ISE, which does not support the
7 inversion parameters.
8
9 To mark a cell port as invertible, use (* invertible_pin = "param_name" *)
10 on the wire in the blackbox module. The parameter value should have
11 the same width as the port, and will be effectively XORed with it.
12
13     -inv <celltype> <portname_out>:<portname_in>
14         Specifies the cell type to use for the inverters and its port names.
15         This option is required.
```

### C.68 flatten – flatten design

```
1     flatten [options] [selection]
2
3 This pass flattens the design by replacing cells by their implementation. This
4 pass is very similar to the 'techmap' pass. The only difference is that this
5 pass is using the current design as mapping library.
6
7 Cells and/or modules with the 'keep_hierarchy' attribute set will not be
8 flattened by this command.
9
10     -wb
11         Ignore the 'whitebox' attribute on cell implementations.
```

### C.69 flowmap – pack LUTs with FlowMap

```
1     flowmap [options] [selection]
2
3 This pass uses the FlowMap technology mapping algorithm to pack logic gates
4 into k-LUTs with optimal depth. It allows mapping any circuit elements that can
5 be evaluated with the `eval` pass, including cells with multiple output ports
6 and multi-bit input and output ports.
7
8     -maxlut k
9         perform technology mapping for a k-LUT architecture. if not specified,
10         defaults to 3.
11
12     -minlut n
13         only produce n-input or larger LUTs. if not specified, defaults to 1.
14
15     -cells <cell>[,<cell>,...]
16         map specified cells. if not specified, maps $_NOT_, $_AND_, $_OR_,
17         $_XOR_ and $_MUX_, which are the outputs of the `simplemap` pass.
18
19     -relax
20         perform depth relaxation and area minimization.
21
22     -r-alpha n, -r-beta n, -r-gamma n
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
23     parameters of depth relaxation heuristic potential function.
24     if not specified, alpha=8, beta=2, gamma=1.
25
26     -optarea n
27         optimize for area by trading off at most n logic levels for fewer LUTs.
28         n may be zero, to optimize for area without increasing depth.
29         implies -relax.
30
31     -debug
32         dump intermediate graphs.
33
34     -debug-relax
35         explain decisions performed during depth relaxation.
```

### C.70 fmcombine – combine two instances of a cell into one

```
1     fmcombine [options] module_name gold_cell gate_cell
2
3     This pass takes two cells, which are instances of the same module, and replaces
4     them with one instance of a special 'combined' module, that effectively
5     contains two copies of the original module, plus some formal properties.
6
7     This is useful for formal test benches that check what differences in behavior
8     a slight difference in input causes in a module.
9
10    -initeq
11        Insert assumptions that initially all FFs in both circuits have the
12        same initial values.
13
14    -anyeq
15        Do not duplicate $anyseq/$anyconst cells.
16
17    -fwd
18        Insert forward hint assumptions into the combined module.
19
20    -bwd
21        Insert backward hint assumptions into the combined module.
22        (Backward hints are logically equivalent to forward hits, but
23        some solvers are faster with bwd hints, or even both -bwd and -fwd.)
24
25    -nop
26        Don't insert hint assumptions into the combined module.
27        (This should not provide any speedup over the original design, but
28        strangely sometimes it does.)
29
30    If none of -fwd, -bwd, and -nop is given, then -fwd is used as default.
```

### C.71 fminit – set init values/sequences for formal

## APPENDIX C. COMMAND REFERENCE MANUAL

```
1  fminit [options] <selection>
2
3  This pass creates init constraints (for example for reset sequences) in a formal
4  model.
5
6  -seq <signal> <sequence>
7      Set sequence using comma-separated list of values, use 'z' for
8      unconstrained bits. The last value is used for the remainder of the
9      trace.
10
11  -set <signal> <value>
12      Add constant value constraint
13
14  -posedge <signal>
15  -negedge <signal>
16      Set clock for init sequences
```

### C.72 freduce – perform functional reduction

```
1  freduce [options] [selection]
2
3  This pass performs functional reduction in the circuit. I.e. if two nodes are
4  equivalent, they are merged to one node and one of the redundant drivers is
5  disconnected. A subsequent call to 'clean' will remove the redundant drivers.
6
7  -v, -vv
8      enable verbose or very verbose output
9
10  -inv
11      enable explicit handling of inverted signals
12
13  -stop <n>
14      stop after <n> reduction operations. this is mostly used for
15      debugging the freduce command itself.
16
17  -dump <prefix>
18      dump the design to <prefix>_<module>_<num>.il after each reduction
19      operation. this is mostly used for debugging the freduce command.
20
21  This pass is undef-aware, i.e. it considers don't-care values for detecting
22  equivalent nodes.
23
24  All selected wires are considered for rewiring. The selected cells cover the
25  circuit that is analyzed.
```

### C.73 fsm – extract and optimize finite state machines

```
1  fsm [options] [selection]
2
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
3 This pass calls all the other fsm_* passes in a useful order. This performs
4 FSM extraction and optimization. It also calls opt_clean as needed:
5
6     fsm_detect          unless got option -nodetect
7     fsm_extract
8
9     fsm_opt
10    opt_clean
11    fsm_opt
12
13    fsm_expand          if got option -expand
14    opt_clean           if got option -expand
15    fsm_opt             if got option -expand
16
17    fsm_recode          unless got option -norecode
18
19    fsm_info
20
21    fsm_export          if got option -export
22    fsm_map             unless got option -nomap
23
24 Options:
25
26     -expand, -norecode, -export, -nomap
27         enable or disable passes as indicated above
28
29     -fullexpand
30         call expand with -full option
31
32     -encoding type
33     -fm_set_fsm_file file
34     -encfile file
35         passed through to fsm_recode pass
```

### C.74 fsm\_detect – finding FSMs in design

```
1     fsm_detect [selection]
2
3 This pass detects finite state machines by identifying the state signal.
4 The state signal is then marked by setting the attribute 'fsm_encoding'
5 on the state signal to "auto".
6
7 Existing 'fsm_encoding' attributes are not changed by this pass.
8
9 Signals can be protected from being detected by this pass by setting the
10 'fsm_encoding' attribute to "none".
```

### C.75 fsm\_expand – expand FSM cells by merging logic into it

```

1      fsm_expand [-full] [selection]
2
3  The fsm_extract pass is conservative about the cells that belong to a finite
4  state machine. This pass can be used to merge additional auxiliary gates into
5  the finite state machine.
6
7  By default, fsm_expand is still a bit conservative regarding merging larger
8  word-wide cells. Call with -full to consider all cells for merging.

```

## C.76 fsm\_export – exporting FSMs to KISS2 files

```

1      fsm_export [-noauto] [-o filename] [-origenc] [selection]
2
3  This pass creates a KISS2 file for every selected FSM. For FSMs with the
4  'fsm_export' attribute set, the attribute value is used as filename, otherwise
5  the module and cell name is used as filename. If the parameter '-o' is given,
6  the first exported FSM is written to the specified filename. This overwrites
7  the setting as specified with the 'fsm_export' attribute. All other FSMs are
8  exported to the default name as mentioned above.
9
10     -noauto
11         only export FSMs that have the 'fsm_export' attribute set
12
13     -o filename
14         filename of the first exported FSM
15
16     -origenc
17         use binary state encoding as state names instead of s0, s1, ...

```

## C.77 fsm\_extract – extracting FSMs in design

```

1      fsm_extract [selection]
2
3  This pass operates on all signals marked as FSM state signals using the
4  'fsm_encoding' attribute. It consumes the logic that creates the state signal
5  and uses the state signal to generate control signal and replaces it with an
6  FSM cell.
7
8  The generated FSM cell still generates the original state signal with its
9  original encoding. The 'fsm_opt' pass can be used in combination with the
10 'opt_clean' pass to eliminate this signal.

```

## C.78 fsm\_info – print information on finite state machines

```

1      fsm_info [selection]
2

```

```

3 | This pass dumps all internal information on FSM cells. It can be useful for
4 | analyzing the synthesis process and is called automatically by the 'fsm'
5 | pass so that this information is included in the synthesis log file.

```

## C.79 fsm\_map – mapping FSMs to basic logic

```

1 | fsm_map [selection]
2 |
3 | This pass translates FSM cells to flip-flops and logic.

```

## C.80 fsm\_opt – optimize finite state machines

```

1 | fsm_opt [selection]
2 |
3 | This pass optimizes FSM cells. It detects which output signals are actually
4 | not used and removes them from the FSM. This pass is usually used in
5 | combination with the 'opt_clean' pass (see also 'help fsm').

```

## C.81 fsm\_recode – recoding finite state machines

```

1 | fsm_recode [options] [selection]
2 |
3 | This pass reassign the state encodings for FSM cells. At the moment only
4 | one-hot encoding and binary encoding is supported.
5 |   -encoding <type>
6 |       specify the encoding scheme used for FSMs without the
7 |       'fsm_encoding' attribute or with the attribute set to `auto'.
8 |
9 |   -fm_set_fsm_file <file>
10 |       generate a file containing the mapping from old to new FSM encoding
11 |       in form of Synopsys Formality set_fsm_* commands.
12 |
13 |   -encfile <file>
14 |       write the mappings from old to new FSM encoding to a file in the
15 |       following format:
16 |
17 |           .fsm <module_name> <state_signal>
18 |           .map <old_bitpattern> <new_bitpattern>

```

## C.82 fst2tb – generate testbench out of fst file

```

1 | fst2tb [options] [top-level]
2 |
3 | This command generates testbench for the circuit using the given top-level module
4 | and stimulus signal from FST file

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
5
6  -tb <name>
7      generated testbench name.
8      files <name>.v and <name>.txt are created as result.
9
10 -r <filename>
11     read simulation FST file
12
13 -clock <portname>
14     name of top-level clock input
15
16 -clockn <portname>
17     name of top-level clock input (inverse polarity)
18
19 -scope <name>
20     scope of simulation top model
21
22 -start <time>
23     start co-simulation in arbitrary time (default 0)
24
25 -stop <time>
26     stop co-simulation in arbitrary time (default END)
27
28 -n <integer>
29     number of clock cycles to simulate (default: 20)
```

### C.83 `gatemate__foldinv` – fold inverters into Gatemate LUT trees

```
1  gatemate_foldinv [selection]
2
3
4  This pass searches for $__CC_NOT cells and folds them into CC_LUT2, CC_L2T4
5  and CC_L2T5 cells as created by LUT tree mapping.
```

### C.84 `glift` – create GLIFT models and optimization problems

```
1  glift <command> [options] [selection]
2
3  Augments the current or specified module with gate-level information flow tracking
4  (GLIFT) logic using the "constructive mapping" approach. Also can set up QBF-SAT
5  optimization problems in order to optimize GLIFT models or trade off precision and
6  complexity.
7
8
9  Commands:
10
11  -create-precise-model
12      Replaces the current or specified module with one that has corresponding "taint"
13      inputs, outputs, and internal nets along with precise taint tracking logic.
14      For example, precise taint tracking logic for an AND gate is:
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

15
16     y_t = a & b_t | b & a_t | a_t & b_t
17
18
19 -create-imprecise-model
20     Replaces the current or specified module with one that has corresponding "taint"
21     inputs, outputs, and internal nets along with imprecise "All OR" taint tracking
22     logic:
23
24     y_t = a_t | b_t
25
26
27 -create-instrumented-model
28     Replaces the current or specified module with one that has corresponding "taint"
29     inputs, outputs, and internal nets along with 4 varying-precision versions of taint
30     tracking logic. Which version of taint tracking logic is used for a given gate is
31     determined by a MUX selected by an $anyconst cell. By default, unless the
32     `no-cost-model` option is provided, an additional wire named `__glift_weight` with
33     the `keep` and `minimize` attributes is added to the module along with pmuxes and
34     adders to calculate a rough estimate of the number of logic gates in the GLIFT model
35     given an assignment for the $anyconst cells. The four versions of taint tracking logic
36     for an AND gate are:
37     y_t = a & b_t | b & a_t | a_t & b_t                (like `create-precise-model`)
38     y_t = a_t | a & b_t
39     y_t = b_t | b & a_t
40     y_t = a_t | b_t                (like `create-imprecise-model`)
41
42
43 Options:
44
45 -taint-constants
46     Constant values in the design are labeled as tainted.
47     (default: label constants as un-tainted)
48
49 -keep-outputs
50     Do not remove module outputs. Taint tracking outputs will appear in the module ports
51     alongside the original outputs.
52     (default: original module outputs are removed)
53
54 -simple-cost-model
55     Do not model logic area. Instead model the number of non-zero assignments to $anyconsts.
56     Taint tracking logic versions vary in their size, but all reduced-precision versions are
57     significantly smaller than the fully-precise version. A non-zero $anyconst assignment means
58     that reduced-precision taint tracking logic was chosen for some gate.
59     Only applicable in combination with `create-instrumented-model`.
60     (default: use a complex model and give that wire the "keep" and "minimize" attributes)
61
62 -no-cost-model
63     Do not model taint tracking logic area and do not create a `__glift_weight` wire.
64     Only applicable in combination with `create-instrumented-model`.
65     (default: model area and give that wire the "keep" and "minimize" attributes)
66
67 -instrument-more
68     Allow choice from more versions of (even simpler) taint tracking logic. A total

```



```

69 of 8 versions of taint tracking logic will be added per gate, including the 4
70 versions from `-create-instrumented-model` and these additional versions:
71
72     y_t = a_t
73     y_t = b_t
74     y_t = 1
75     y_t = 0
76
77 Only applicable in combination with `-create-instrumented-model`.
78 (default: do not add more versions of taint tracking logic.

```

## C.85 greenpak4\_dffinv – merge greenpak4 inverters and DF-F/latches

```

1  greenpak4_dffinv [options] [selection]
2
3  Merge GP_INV cells with GP_DFF* and GP_DLATCH* cells.

```

## C.86 help – display help messages

```

1  help ..... list all commands
2  help <command> ..... print help message for given command
3  help -all ..... print complete command reference
4
5  help -cells ..... list all cell types
6  help <celltype> ..... print help message for given cell type
7  help <celltype>+ .... print verilog code for given cell type

```

## C.87 hierarchy – check, expand and clean up design hierarchy

```

1  hierarchy [-check] [-top <module>]
2  hierarchy -generate <cell-types> <port-decls>
3
4  In parametric designs, a module might exists in several variations with
5  different parameter values. This pass looks at all modules in the current
6  design and re-runs the language frontends for the parametric modules as
7  needed. It also resolves assignments to wired logic data types (wand/wor),
8  resolves positional module parameters, unrolls array instances, and more.
9
10  -check
11      also check the design hierarchy. this generates an error when
12      an unknown module is used as cell type.
13
14  -simcheck
15      like -check, but also throw an error if blackbox modules are
16      instantiated, and throw an error if the design has no top module.
17

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
18  -smtcheck
19      like -simcheck, but allow smtlib2_module modules.
20
21  -purge_lib
22      by default the hierarchy command will not remove library (blackbox)
23      modules. use this option to also remove unused blackbox modules.
24
25  -libdir <directory>
26      search for files named <module_name>.v in the specified directory
27      for unknown modules and automatically run read_verilog for each
28      unknown module.
29
30  -keep_positionals
31      per default this pass also converts positional arguments in cells
32      to arguments using port names. This option disables this behavior.
33
34  -keep_portwidths
35      per default this pass adjusts the port width on cells that are
36      module instances when the width does not match the module port. This
37      option disables this behavior.
38
39  -nodefaults
40      do not resolve input port default values
41
42  -nokeep_asserts
43      per default this pass sets the "keep" attribute on all modules
44      that directly or indirectly contain one or more formal properties.
45      This option disables this behavior.
46
47  -top <module>
48      use the specified top module to build the design hierarchy. Modules
49      outside this tree (unused modules) are removed.
50
51      when the -top option is used, the 'top' attribute will be set on the
52      specified top module. otherwise a module with the 'top' attribute set
53      will implicitly be used as top module, if such a module exists.
54
55  -auto-top
56      automatically determine the top of the design hierarchy and mark it.
57
58  -chparam name value
59      elaborate the top module using this parameter value. Modules on which
60      this parameter does not exist may cause a warning message to be output.
61      This option can be specified multiple times to override multiple
62      parameters. String values must be passed in double quotes ("").
63
64  In -generate mode this pass generates blackbox modules for the given cell
65  types (wildcards supported). For this the design is searched for cells that
66  match the given types and then the given port declarations are used to
67  determine the direction of the ports. The syntax for a port declaration is:
68
69      {i|o|io}[@<num>]:<portname>
70
71  Input ports are specified with the 'i' prefix, output ports with the 'o'
```

```

72 | prefix and inout ports with the 'io' prefix. The optional <num> specifies
73 | the position of the port in the parameter list (needed when instantiated
74 | using positional arguments). When <num> is not specified, the <portname> can
75 | also contain wildcard characters.
76 |
77 | This pass ignores the current selection and always operates on all modules
78 | in the current design.

```

## C.88 hilomap – technology mapping of constant hi- and/or lo-drivers

```

1 |   hilomap [options] [selection]
2 |
3 | Map constants to 'tielo' and 'tiehi' driver cells.
4 |
5 |   -hicell <celltype> <portname>
6 |       Replace constant hi bits with this cell.
7 |
8 |   -locell <celltype> <portname>
9 |       Replace constant lo bits with this cell.
10 |
11 |   -singleton
12 |       Create only one hi/lo cell and connect all constant bits
13 |       to that cell. Per default a separate cell is created for
14 |       each constant bit.

```

## C.89 history – show last interactive commands

```

1 |   history
2 |
3 | This command prints all commands in the shell history buffer. This are
4 | all commands executed in an interactive session, but not the commands
5 | from executed scripts.

```

## C.90 ice40\_braminit – iCE40: perform SB\_RAM40\_4K initialization from file

```

1 |   ice40_braminit
2 |
3 | This command processes all SB_RAM40_4K blocks with a non-empty INIT_FILE
4 | parameter and converts it into the required INIT_x attributes

```

## C.91 ice40\_dsp – iCE40: map multipliers

```

1      ice40_dsp [options] [selection]
2
3      Map multipliers ($mul/SB_MAC16) and multiply-accumulate ($mul/SB_MAC16 + $add)
4      cells into iCE40 DSP resources.
5      Currently, only the 16x16 multiply mode is supported and not the 2 x 8x8 mode.
6
7      Pack input registers (A, B, {C,D}; with optional hold), pipeline registers
8      ({F,J,K,G}, H), output registers (O -- full 32-bits or lower 16-bits only; with
9      optional hold), and post-adder into the SB_MAC16 resource.
10
11     Multiply-accumulate operations using the post-adder with feedback on the {C,D}
12     input will be folded into the DSP. In this scenario only, resetting the
13     the accumulator to an arbitrary value can be inferred to use the {C,D} input.

```

## C.92 ice40\_opt – iCE40: perform simple optimizations

```

1      ice40_opt [options] [selection]
2
3      This command executes the following script:
4
5      do
6          <ice40 specific optimizations>
7          opt_expr -mux_undef -undriven [-full]
8          opt_merge
9          opt_dff
10         opt_clean
11     while <changed design>

```

## C.93 ice40\_wrapcarry – iCE40: wrap carries

```

1      ice40_wrapcarry [selection]
2
3      Wrap manually instantiated SB_CARRY cells, along with their associated SB_LUT4s,
4      into an internal $__ICE40_CARRY_WRAPPER cell for preservation across technology
5      mapping.
6
7      Attributes on both cells will have their names prefixed with 'SB_CARRY.' or
8      'SB_LUT4.' and attached to the wrapping cell.
9      A (* keep *) attribute on either cell will be logically OR-ed together.
10
11     -unwrap
12         unwrap $__ICE40_CARRY_WRAPPER cells back into SB_CARRYs and SB_LUT4s,
13         including restoring their attributes.

```

## C.94 insbuf – insert buffer cells for connected wires

```

1      insbuf [options] [selection]
2
3      Insert buffer cells into the design for directly connected wires.
4
5      -buf <celltype> <in-portname> <out-portname>
6          Use the given cell type instead of $_BUF_. (Notice that the next
7          call to "clean" will remove all $_BUF_ in the design.)

```

## C.95 iopadmap – technology mapping of i/o pads (or buffers)

```

1      iopadmap [options] [selection]
2
3      Map module inputs/outputs to PAD cells from a library. This pass
4      can only map to very simple PAD cells. Use 'techmap' to further map
5      the resulting cells to more sophisticated PAD cells.
6
7      -inpad <celltype> <in_port>[:<ext_port>]
8          Map module input ports to the given cell type with the
9          given output port name. if a 2nd portname is given, the
10         signal is passed through the pad cell, using the 2nd
11         portname as the port facing the module port.
12
13      -outpad <celltype> <out_port>[:<ext_port>]
14      -inoutpad <celltype> <io_port>[:<ext_port>]
15          Similar to -inpad, but for output and inout ports.
16
17      -toutpad <celltype> <oe_port>:<out_port>[:<ext_port>]
18          Merges $_TBUF_ cells into the output pad cell. This takes precedence
19          over the other -outpad cell. The first portname is the enable input
20          of the tristate driver, which can be prefixed with `~` for negative
21          polarity enable.
22
23      -tinoutpad <celltype> <oe_port>:<in_port>:<out_port>[:<ext_port>]
24          Merges $_TBUF_ cells into the inout pad cell. This takes precedence
25          over the other -inoutpad cell. The first portname is the enable input
26          of the tristate driver and the 2nd portname is the internal output
27          buffering the external signal. Like with `~`-toutpad`, the enable can
28          be marked as negative polarity by prefixing the name with `~`.
29
30      -ignore <celltype> <portname>[:<portname>]*
31          Skips mapping inputs/outputs that are already connected to given
32          ports of the given cell. Can be used multiple times. This is in
33          addition to the cells specified as mapping targets.
34
35      -widthparam <param_name>
36          Use the specified parameter name to set the port width.
37
38      -nameparam <param_name>
39          Use the specified parameter to set the port name.
40

```

```

41     -bits
42         create individual bit-wide buffers even for ports that
43         are wider. (the default behavior is to create word-wide
44         buffers using -widthparam to set the word size on the cell.)
45
46 Tristate PADS (-toutpad, -tinoutpad) always operate in -bits mode.

```

## C.96 jny – write design and metadata

```

1     jny [options] [selection]
2
3 Write a JSON netlist metadata for the current design
4
5     -o <filename>
6         write to the specified file.
7
8     -no-connections
9         Don't include connection information in the netlist output.
10
11     -no-attributes
12         Don't include attributed information in the netlist output.
13
14     -no-properties
15         Don't include property information in the netlist output.
16
17 See 'help write_jny' for a description of the JSON format used.

```

## C.97 json – write design in JSON format

```

1     json [options] [selection]
2
3 Write a JSON netlist of all selected objects.
4
5     -o <filename>
6         write to the specified file.
7
8     -aig
9         also include AIG models for the different gate types
10
11     -compat-int
12         emit 32-bit or smaller fully-defined parameter values directly
13         as JSON numbers (for compatibility with old parsers)
14
15 See 'help write_json' for a description of the JSON format used.

```

## C.98 log – print text and log files

```

1      log string
2
3      Print the given string to the screen and/or the log file. This is useful for TCL
4      scripts, because the TCL command "puts" only goes to stdout but not to
5      logfiles.
6
7      -stdout
8          Print the output to stdout too. This is useful when all Yosys is executed
9          with a script and the -q (quiet operation) argument to notify the user.
10
11     -stderr
12         Print the output to stderr too.
13
14     -nolog
15         Don't use the internal log() command. Use either -stdout or -stderr,
16         otherwise no output will be generated at all.
17
18     -n
19         do not append a newline

```

## C.99 logger – set logger properties

```

1      logger [options]
2
3      This command sets global logger properties, also available using command line
4      options.
5
6      -[no]time
7          enable/disable display of timestamp in log output.
8
9      -[no]stderr
10         enable/disable logging errors to stderr.
11
12     -warn regex
13         print a warning for all log messages matching the regex.
14
15     -nowarn regex
16         if a warning message matches the regex, it is printed as regular
17         message instead.
18
19     -werror regex
20         if a warning message matches the regex, it is printed as error
21         message instead and the tool terminates with a nonzero return code.
22
23     -[no]debug
24         globally enable/disable debug log messages.
25
26     -experimental <feature>
27         do not print warnings for the specified experimental feature
28
29     -expect <type> <regex> <expected_count>

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
30     expect log, warning or error to appear. matched errors will terminate
31     with exit code 0.
32
33     -expect-no-warnings
34         gives error in case there is at least one warning that is not expected.
35
36     -check-expected
37         verifies that the patterns previously set up by -expect have actually
38         been met, then clears the expected log list. If this is not called
39         manually, the check will happen at yosys exist time instead.
```

### C.100 ls – list modules or objects in modules

```
1     ls [selection]
2
3     When no active module is selected, this prints a list of modules.
4
5     When an active module is selected, this prints a list of objects in the module.
```

### C.101 ltp – print longest topological path

```
1     ltp [options] [selection]
2
3     This command prints the longest topological path in the design. (Only considers
4     paths within a single module, so the design must be flattened.)
5
6     -noff
7         automatically exclude FF cell types
```

### C.102 lut2mux – convert \$lut to \$\_\_MUX\_\_

```
1     lut2mux [options] [selection]
2
3     This pass converts $lut cells to $__MUX_ gates.
```

### C.103 maccmap – mapping macc cells

```
1     maccmap [-unmap] [selection]
2
3     This pass maps $macc cells to yosys $fa and $alu cells. When the -unmap option
4     is used then the $macc cell is mapped to $add, $sub, etc. cells instead.
```



## C.104 memory – translate memories to basic cells

```

1      memory [-norom] [-nomap] [-nordff] [-nowiden] [-nosat] [-memx] [-no-rw-check] [-bram <bram_rules>] [sel
2
3  This pass calls all the other memory_* passes in a useful order:
4
5      opt_mem
6      opt_mem_priority
7      opt_mem_feedback
8      memory_bmux2rom                      (skipped if called with -norom)
9      memory_dff [-no-rw-check]            (skipped if called with -nordff or -memx)
10     opt_clean
11     memory_share [-nowiden] [-nosat]
12     opt_mem_widen
13     memory_memx                          (when called with -memx)
14     opt_clean
15     memory_collect
16     memory_bram -rules <bram_rules>      (when called with -bram)
17     memory_map                          (skipped if called with -nomap)
18
19  This converts memories to word-wide DFFs and address decoders
20  or multiport memory blocks if called with the -nomap option.

```

## C.105 memory\_\_bmux2rom – convert muxes to ROMs

```

1      memory_bmux2rom [options] [selection]
2
3  This pass converts $bmux cells with constant A input to ROMs.

```

## C.106 memory\_\_bram – map memories to block rams

```

1      memory_bram -rules <rule_file> [selection]
2
3  This pass converts the multi-port $mem memory cells into block ram instances.
4  The given rules file describes the available resources and how they should be
5  used.
6
7  The rules file contains configuration options, a set of block ram description
8  and a sequence of match rules.
9
10 The option 'attr_icode' configures how attribute values are matched. The value 0
11 means case-sensitive, 1 means case-insensitive.
12
13 A block ram description looks like this:
14
15     bram RAMB1024X32      # name of BRAM cell
16     init 1                # set to '1' if BRAM can be initialized
17     abits 10              # number of address bits
18     dbits 32              # number of data bits

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

19     groups 2          # number of port groups
20     ports 1 1         # number of ports in each group
21     wrmode 1 0        # set to '1' if this groups is write ports
22     enable 4 1        # number of enable bits
23     transp 0 2        # transparent (for read ports)
24     clocks 1 2        # clock configuration
25     clkpol 2 2        # clock polarity configuration
26     endbram
27
28 For the option 'transp' the value 0 means non-transparent, 1 means transparent
29 and a value greater than 1 means configurable. All groups with the same
30 value greater than 1 share the same configuration bit.
31
32 For the option 'clocks' the value 0 means non-clocked, and a value greater
33 than 0 means clocked. All groups with the same value share the same clock
34 signal.
35
36 For the option 'clkpol' the value 0 means negative edge, 1 means positive edge
37 and a value greater than 1 means configurable. All groups with the same value
38 greater than 1 share the same configuration bit.
39
40 Using the same bram name in different bram blocks will create different variants
41 of the bram. Verilog configuration parameters for the bram are created as needed.
42
43 It is also possible to create variants by repeating statements in the bram block
44 and appending '@<label>' to the individual statements.
45
46 A match rule looks like this:
47
48     match RAMB1024X32
49         max waste 16384    # only use this bram if <= 16k ram bits are unused
50         min efficiency 80  # only use this bram if efficiency is at least 80%
51     endmatch
52
53 It is possible to match against the following values with min/max rules:
54
55     words ..... number of words in memory in design
56     abits ..... number of address bits on memory in design
57     dbits ..... number of data bits on memory in design
58     wports ..... number of write ports on memory in design
59     rports ..... number of read ports on memory in design
60     ports ..... number of ports on memory in design
61     bits ..... number of bits in memory in design
62     dups ..... number of duplications for more read ports
63
64     awaste ..... number of unused address slots for this match
65     dwaste ..... number of unused data bits for this match
66     bwaste ..... number of unused bram bits for this match
67     waste ..... total number of unused bram bits (bwaste*dups)
68     efficiency ... total percentage of used and non-duplicated bits
69
70     acells ..... number of cells in 'address-direction'
71     dcells ..... number of cells in 'data-direction'
72     cells ..... total number of cells (acells*dcells*dups)

```

```

73
74 A match containing the command 'attribute' followed by a list of space
75 separated 'name [=string_value]' values requires that the memory contains any
76 one of the given attribute name and string values (where specified), or name
77 and integer 1 value (if no string_value given, since Verilog will interpret
78 '(* attr *)' as '(* attr=1 *)').
79 A name prefixed with '!' indicates that the attribute must not exist.
80
81 The interface for the created bram instances is derived from the bram
82 description. Use 'techmap' to convert the created bram instances into
83 instances of the actual bram cells of your target architecture.
84
85 A match containing the command 'or_next_if_better' is only used if it
86 has a higher efficiency than the next match (and the one after that if
87 the next also has 'or_next_if_better' set, and so forth).
88
89 A match containing the command 'make_transp' will add external circuitry
90 to simulate 'transparent read', if necessary.
91
92 A match containing the command 'make_outreg' will add external flip-flops
93 to implement synchronous read ports, if necessary.
94
95 A match containing the command 'shuffle_enable A' will re-organize
96 the data bits to accommodate the enable pattern of port A.

```

## C.107 memory\_collect – creating multi-port memory cells

```

1 memory_collect [selection]
2
3 This pass collects memories and memory ports and creates generic multiport
4 memory cells.

```

## C.108 memory\_dff – merge input/output DFFs into memory read ports

```

1 memory_dff [-no-rw-check] [selection]
2
3 This pass detects DFFs at memory read ports and merges them into the memory port.
4 I.e. it consumes an asynchronous memory port and the flip-flops at its
5 interface and yields a synchronous memory port.
6
7 -no-rw-check
8 marks all recognized read ports as "return don't-care value on
9 read/write collision" (same result as setting the no_rw_check
10 attribute on all memories).

```

**C.109 memory\_libmap – map memories to cells**

```

1  memory_libmap -lib <library_file> [-D <condition>] [selection]
2
3  This pass takes a description of available RAM cell types and maps
4  all selected memories to one of them, or leaves them to be mapped to FFs.
5
6  -lib <library_file>
7      Selects a library file containing RAM cell definitions. This option
8      can be passed more than once to select multiple libraries.
9      See passes/memory/memlib.md for description of the library format.
10
11  -D <condition>
12      Enables a condition that can be checked within the library file
13      to eg. select between slightly different hardware variants.
14      This option can be passed any number of times.
15
16  -logic-cost-rom <num>
17  -logic-cost-ram <num>
18      Sets the cost of a single bit for memory lowered to soft logic.
19
20  -no-auto-distributed
21  -no-auto-block
22  -no-auto-huge
23      Disables automatic mapping of given kind of RAMs. Manual mapping
24      (using ram_style or other attributes) is still supported.

```

**C.110 memory\_map – translate multiport memories to basic cells**

```

1  memory_map [options] [selection]
2
3  This pass converts multiport memory cells as generated by the memory_collect
4  pass to word-wide DFFs and address decoders.
5
6  -attr !<name>
7      do not map memories that have attribute <name> set.
8
9  -attr <name>[=<value>]
10      for memories that have attribute <name> set, only map them if its value
11      is a string <value> (if specified), or an integer 1 (otherwise). if this
12      option is specified multiple times, map the memory if the attribute is
13      to any of the values.
14
15  -iattr
16      for -attr, ignore case of <value>.
17
18  -rom-only
19      only perform conversion for ROMs (memories with no write ports).
20
21  -keepdc

```

22 when mapping ROMs, keep x-bits shared across read ports.

### C.111 `memory_memx` – emulate vlog sim behavior for mem ports

```
1  memory_memx [selection]
2
3  This pass adds additional circuitry that emulates the Verilog simulation
4  behavior for out-of-bounds memory reads and writes.
```

### C.112 `memory_narrow` – split up wide memory ports

```
1  memory_narrow [options] [selection]
2
3  This pass splits up wide memory ports into several narrow ports.
```

### C.113 `memory_nordff` – extract read port FFs from memories

```
1  memory_nordff [options] [selection]
2
3  This pass extracts FFs from memory read ports. This results in a netlist
4  similar to what one would get from not calling memory_dff.
```

### C.114 `memory_share` – consolidate memory ports

```
1  memory_share [-nosat] [-nowiden] [selection]
2
3  This pass merges share-able memory ports into single memory ports.
4
5  The following methods are used to consolidate the number of memory ports:
6
7  - When multiple write ports access the same address then this is converted
8    to a single write port with a more complex data and/or enable logic path.
9
10 - When multiple read or write ports access adjacent aligned addresses, they are
11    merged to a single wide read or write port. This transformation can be
12    disabled with the "-nowiden" option.
13
14 - When multiple write ports are never accessed at the same time (a SAT
15    solver is used to determine this), then the ports are merged into a single
16    write port. This transformation can be disabled with the "-nosat" option.
17
18 Note that in addition to the algorithms implemented in this pass, the $memrd
19 and $memwr cells are also subject to generic resource sharing passes (and other
20 optimizations) such as "share" and "opt_merge".
```

## C.115 `memory_unpack` – unpack multi-port memory cells

```

1  memory_unpack [selection]
2
3  This pass converts the multi-port $mem memory cells into individual $memrd and
4  $memwr cells. It is the counterpart to the memory_collect pass.
```

## C.116 `miter` – automatically create a miter circuit

```

1  miter -equiv [options] gold_name gate_name miter_name
2
3  Creates a miter circuit for equivalence checking. The gold- and gate- modules
4  must have the same interfaces. The miter circuit will have all inputs of the
5  two source modules, prefixed with 'in_'. The miter circuit has a 'trigger'
6  output that goes high if an output mismatch between the two source modules is
7  detected.
8
9  -ignore_gold_x
10     a undef (x) bit in the gold module output will match any value in
11     the gate module output.
12
13  -make_outputs
14     also route the gold- and gate-outputs to 'gold_*' and 'gate_*' outputs
15     on the miter circuit.
16
17  -make_outcmp
18     also create a cmp_* output for each gold/gate output pair.
19
20  -make_assert
21     also create an 'assert' cell that checks if trigger is always low.
22
23  -flatten
24     call 'flatten -wb; opt_expr -keepdc -undriven;;' on the miter circuit.
25
26
27  miter -assert [options] module [miter_name]
28
29  Creates a miter circuit for property checking. All input ports are kept,
30  output ports are discarded. An additional output 'trigger' is created that
31  goes high when an assert is violated. Without a miter_name, the existing
32  module is modified.
33
34  -make_outputs
35     keep module output ports.
36
37  -flatten
38     call 'flatten -wb; opt_expr -keepdc -undriven;;' on the miter circuit.
```

## C.117 `mutate` – generate or apply design mutations

## APPENDIX C. COMMAND REFERENCE MANUAL

```
1      mutate -list N [options] [selection]
2
3  Create a list of N mutations using an even sampling.
4
5      -o filename
6          Write list to this file instead of console output
7
8      -s filename
9          Write a list of all src tags found in the design to the specified file
10
11     -seed N
12         RNG seed for selecting mutations
13
14     -none
15         Include a "none" mutation in the output
16
17     -ctrl name width value
18         Add -ctrl options to the output. Use 'value' for first mutation, then
19         simply count up from there.
20
21     -mode name
22     -module name
23     -cell name
24     -port name
25     -portbit int
26     -ctrlbit int
27     -wire name
28     -wirebit int
29     -src string
30         Filter list of mutation candidates to those matching
31         the given parameters.
32
33     -cfg option int
34         Set a configuration option. Options available:
35         weight_pq_w weight_pq_b weight_pq_c weight_pq_s
36         weight_pq_mw weight_pq_mb weight_pq_mc weight_pq_ms
37         weight_cover pick_cover_prcnt
38
39
40     mutate -mode MODE [options]
41
42  Apply the given mutation.
43
44     -ctrl name width value
45         Add a control signal with the given name and width. The mutation is
46         activated if the control signal equals the given value.
47
48     -module name
49     -cell name
50     -port name
51     -portbit int
52     -ctrlbit int
53         Mutation parameters, as generated by 'mutate -list N'.
54
```

```

55 -wire name
56 -wirebit int
57 -src string
58 Ignored. (They are generated by -list for documentation purposes.)

```

## C.118 muxcover – cover trees of MUX cells with wider MUXes

```

1  muxcover [options] [selection]
2
3  Cover trees of $_MUX_ cells with $_MUX{4,8,16}_ cells
4
5  -mux4[=cost], -mux8[=cost], -mux16[=cost]
6      Cover $_MUX_ trees using the specified types of MUXes (with optional
7      integer costs). If none of these options are given, the effect is the
8      same as if all of them are.
9      Default costs: $_MUX4_ = 220, $_MUX8_ = 460,
10                     $_MUX16_ = 940
11
12  -mux2=cost
13      Use the specified cost for $_MUX_ cells when making covering decisions.
14      Default cost: $_MUX_ = 100
15
16  -dmux=cost
17      Use the specified cost for $_MUX_ cells used in decoders.
18      Default cost: 90
19
20  -nodecode
21      Do not insert decoder logic. This reduces the number of possible
22      substitutions, but guarantees that the resulting circuit is not
23      less efficient than the original circuit.
24
25  -nopartial
26      Do not consider mappings that use $_MUX<N>_ to select from less
27      than <N> different signals.

```

## C.119 muxpack – \$mux/\$pmux cascades to \$pmux

```

1  muxpack [selection]
2
3  This pass converts cascaded chains of $pmux cells (e.g. those create from case
4  constructs) and $mux cells (e.g. those created by if-else constructs) into
5  $pmux cells.
6
7  This optimisation is conservative --- it will only pack $mux or $pmux cells
8  whose select lines are driven by '$eq' cells with other such cells if it can be
9  certain that their select inputs are mutually exclusive.

```



## C.120 nlutmap – map to LUTs of different sizes

```

1      nlutmap [options] [selection]
2
3      This pass uses successive calls to 'abc' to map to an architecture. That
4      provides a small number of differently sized LUTs.
5
6      -luts N_1,N_2,N_3,...
7          The number of LUTs with 1, 2, 3, ... inputs that are
8          available in the target architecture.
9
10     -assert
11         Create an error if not all logic can be mapped
12
13     Excess logic that does not fit into the specified LUTs is mapped back
14     to generic logic gates ($_AND_, etc.).

```

## C.121 onehot – optimize \$eq cells for onehot signals

```

1      onehot [options] [selection]
2
3      This pass optimizes $eq cells that compare one-hot signals against constants
4
5      -v, -vv
6          verbose output

```

## C.122 opt – perform simple optimizations

```

1      opt [options] [selection]
2
3      This pass calls all the other opt_* passes in a useful order. This performs
4      a series of trivial optimizations and cleanups. This pass executes the other
5      passes in the following order:
6
7      opt_expr [-mux_undef] [-mux_bool] [-undriven] [-noclkinv] [-fine] [-full] [-keepdc]
8      opt_merge [-share_all] -nomux
9
10     do
11         opt_muxtree
12         opt_reduce [-fine] [-full]
13         opt_merge [-share_all]
14         opt_share (-full only)
15         opt_dff [-nodffe] [-nosdff] [-keepdc] [-sat] (except when called with -noff)
16         opt_clean [-purge]
17         opt_expr [-mux_undef] [-mux_bool] [-undriven] [-noclkinv] [-fine] [-full] [-keepdc]
18     while <changed design>
19
20     When called with -fast the following script is used instead:
21

```

```

22     do
23         opt_expr [-mux_undef] [-mux_bool] [-undriven] [-noclkinv] [-fine] [-full] [-keepdc]
24         opt_merge [-share_all]
25         opt_dff [-nodffe] [-nosdff] [-keepdc] [-sat] (except when called with -noff)
26         opt_clean [-purge]
27     while <changed design in opt_dff>
28
29 Note: Options in square brackets (such as [-keepdc]) are passed through to
30 the opt_* commands when given to 'opt'.

```

### C.123 opt\_clean – remove unused cells and wires

```

1     opt_clean [options] [selection]
2
3 This pass identifies wires and cells that are unused and removes them. Other
4 passes often remove cells but leave the wires in the design or reconnect the
5 wires but leave the old cells in the design. This pass can be used to clean up
6 after the passes that do the actual work.
7
8 This pass only operates on completely selected modules without processes.
9
10    -purge
11        also remove internal nets if they have a public name

```

### C.124 opt\_demorgan – Optimize reductions with DeMorgan equivalents

```

1     opt_demorgan [selection]
2
3 This pass pushes inverters through $reduce_* cells if this will reduce the
4 overall gate count of the circuit

```

### C.125 opt\_dff – perform DFF optimizations

```

1     opt_dff [-nodffe] [-nosdff] [-keepdc] [-sat] [selection]
2
3 This pass converts flip-flops to a more suitable type by merging clock enables
4 and synchronous reset multiplexers, removing unused control inputs, or potentially
5 removes the flip-flop altogether, converting it to a constant driver.
6
7    -nodffe
8        disables dff -> dffe conversion, and other transforms recognizing clock enable
9
10    -nosdff
11        disables dff -> sdff conversion, and other transforms recognizing sync resets
12
13    -simple-dffe

```

```

14         only enables clock enable recognition transform for obvious cases
15
16     -sat
17         additionally invoke SAT solver to detect and remove flip-flops (with
18         non-constant inputs) that can also be replaced with a constant driver
19
20     -keepdc
21         some optimizations change the behavior of the circuit with respect to
22         don't-care bits. for example in 'a+0' a single x-bit in 'a' will cause
23         all result bits to be set to x. this behavior changes when 'a+0' is
24         replaced by 'a'. the -keepdc option disables all such optimizations.

```

## C.126 `opt_expr` – perform const folding and simple expression rewriting

```

1     opt_expr [options] [selection]
2
3     This pass performs const folding on internal cell types with constant inputs.
4     It also performs some simple expression rewriting.
5
6     -mux_undef
7         remove 'undef' inputs from $mux, $pmux and $_MUX_ cells
8
9     -mux_bool
10        replace $mux cells with inverters or buffers when possible
11
12    -undriven
13        replace undriven nets with undef (x) constants
14
15    -noclkinv
16        do not optimize clock inverters by changing FF types
17
18    -fine
19        perform fine-grain optimizations
20
21    -full
22        alias for -mux_undef -mux_bool -undriven -fine
23
24    -keepdc
25        some optimizations change the behavior of the circuit with respect to
26        don't-care bits. for example in 'a+0' a single x-bit in 'a' will cause
27        all result bits to be set to x. this behavior changes when 'a+0' is
28        replaced by 'a'. the -keepdc option disables all such optimizations.

```

## C.127 `opt_ffinv` – push inverters through FFs

```

1     opt_ffinv [selection]
2
3     This pass pushes inverters to the other side of a FF when they can be merged

```

4 into LUTs on the other side.

## C.128 `opt_lut` – optimize LUT cells

```

1  opt_lut [options] [selection]
2
3  This pass combines cascaded $lut cells with unused inputs.
4
5  -dlogic <type>:<cell-port>=<LUT-input>[:<cell-port>=<LUT-input>...]
6      preserve connections to dedicated logic cell <type> that has ports
7      <cell-port> connected to LUT inputs <LUT-input>. this includes
8      the case where both LUT and dedicated logic input are connected to
9      the same constant.
10
11  -limit N
12      only perform the first N combines, then stop. useful for debugging.
```

## C.129 `opt_lut_ins` – discard unused LUT inputs

```

1  opt_lut_ins [options] [selection]
2
3  This pass removes unused inputs from LUT cells (that is, inputs that can not
4  influence the output signal given this LUT's value). While such LUTs cannot
5  be directly emitted by ABC, they can be a result of various post-ABC
6  transformations, such as mapping wide LUTs (not all sub-LUTs will use the
7  full set of inputs) or optimizations such as xilinx_dffopt.
8
9  -tech <technology>
10      Instead of generic $lut cells, operate on LUT cells specific
11      to the given technology. Valid values are: xilinx, ecp5, gowin.
```

## C.130 `opt_mem` – optimize memories

```

1  opt_mem [options] [selection]
2
3  This pass performs various optimizations on memories in the design.
```

## C.131 `opt_mem_feedback` – convert memory read-to-write port feedback paths to write enables

```

1  opt_mem_feedback [selection]
2
3  This pass detects cases where an asynchronous read port is only connected via
4  a mux tree to a write port with the same address. When such a connection is
```

```

5 found, it is replaced with a new condition on an enable signal, allowing
6 for removal of the read port.

```

### C.132 `opt_mem_priority` – remove priority relations between write ports that can never collide

```

1  opt_mem_priority [selection]
2
3  This pass detects cases where one memory write port has priority over another
4  even though they can never collide with each other -- ie. there can never be
5  a situation where a given memory bit is written by both ports at the same
6  time, for example because of always-different addresses, or mutually exclusive
7  enable signals. In such cases, the priority relation is removed.

```

### C.133 `opt_mem_widen` – optimize memories where all ports are wide

```

1  opt_mem_widen [options] [selection]
2
3  This pass looks for memories where all ports are wide and adjusts the base
4  memory width up until that stops being the case.

```

### C.134 `opt_merge` – consolidate identical cells

```

1  opt_merge [options] [selection]
2
3  This pass identifies cells with identical type and input signals. Such cells
4  are then merged to one cell.
5
6  -nomux
7      Do not merge MUX cells.
8
9  -share_all
10     Operate on all cell types, not just built-in types.
11
12  -keepdc
13     Do not merge flipflops with don't-care bits in their initial value.

```

### C.135 `opt_muxtree` – eliminate dead trees in multiplexer trees

```

1  opt_muxtree [selection]
2
3  This pass analyzes the control signals for the multiplexer trees in the design
4  and identifies inputs that can never be active. It then removes this dead

```

```

5 | branches from the multiplexer trees.
6 |
7 | This pass only operates on completely selected modules without processes.

```

### C.136 `opt_reduce` – simplify large MUXes and AND/OR gates

```

1 |   opt_reduce [options] [selection]
2 |
3 | This pass performs two interlinked optimizations:
4 |
5 | 1. it consolidates trees of large AND gates or OR gates and eliminates
6 |    duplicated inputs.
7 |
8 | 2. it identifies duplicated inputs to MUXes and replaces them with a single
9 |    input with the original control signals OR'ed together.
10 |
11 |   -fine
12 |       perform fine-grain optimizations
13 |
14 |   -full
15 |       alias for -fine

```

### C.137 `opt_share` – merge mutually exclusive cells of the same type that share an input signal

```

1 |   opt_share [selection]
2 |
3 | This pass identifies mutually exclusive cells of the same type that:
4 |   (a) share an input signal,
5 |   (b) drive the same $mux, $_MUX_, or $pmux multiplexing cell,
6 |
7 | allowing the cell to be merged and the multiplexer to be moved from
8 | multiplexing its output to multiplexing the non-shared input signals.

```

### C.138 `paramap` – renaming cell parameters

```

1 |   paramap [options] [selection]
2 |
3 | This command renames cell parameters and/or maps key/value pairs to
4 | other key/value pairs.
5 |
6 |   -tocase <name>
7 |       Match attribute names case-insensitively and set it to the specified
8 |       name.
9 |
10 |   -rename <old_name> <new_name>
11 |       Rename attributes as specified

```

```

12
13     -map <old_name>=<old_value> <new_name>=<new_value>
14         Map key/value pairs as indicated.
15
16     -imap <old_name>=<old_value> <new_name>=<new_value>
17         Like -map, but use case-insensitive match for <old_value> when
18         it is a string value.
19
20     -remove <name>=<value>
21         Remove attributes matching this pattern.
22
23 For example, mapping Diamond-style ECP5 "init" attributes to Yosys-style:
24
25     paramap -tocase INIT t:LUT4

```

### C.139 peepopt – collection of peephole optimizers

```

1     peepopt [options] [selection]
2
3 This pass applies a collection of peephole optimizers to the current design.

```

### C.140 plugin – load and list loaded plugins

```

1     plugin [options]
2
3 Load and list loaded plugins.
4
5     -i <plugin_filename>
6         Load (install) the specified plugin.
7
8     -a <alias_name>
9         Register the specified alias name for the loaded plugin
10
11     -l
12         List loaded plugins

```

### C.141 pmux2shiftx – transform \$pmux cells to \$shiftx cells

```

1     pmux2shiftx [options] [selection]
2
3 This pass transforms $pmux cells to $shiftx cells.
4
5     -v, -vv
6         verbose output
7
8     -min_density <percentage>
9         specifies the minimum density for the shifter

```

```

10     default: 50
11
12     -min_choices <int>
13         specified the minimum number of choices for a control signal
14         default: 3
15
16     -onehot ignore|pmux|shiftx
17         select strategy for one-hot encoded control signals
18         default: pmux
19
20     -norange
21         disable $sub inference for "range decoders"

```

## C.142 pmuxtree – transform \$pmux cells to trees of \$mux cells

```

1     pmuxtree [selection]
2
3     This pass transforms $pmux cells to trees of $mux cells.

```

## C.143 portlist – list (top-level) ports

```

1     portlist [options] [selection]
2
3     This command lists all module ports found in the selected modules.
4
5     If no selection is provided then it lists the ports on the top module.
6
7     -m
8         print verilog blackbox module definitions instead of port lists

```

## C.144 prep – generic synthesis script

```

1     prep [options]
2
3     This command runs a conservative RTL synthesis. A typical application for this
4     is the preparation stage of a verification flow. This command does not operate
5     on partly selected designs.
6
7     -top <module>
8         use the specified module as top module (default='top')
9
10    -auto-top
11        automatically determine the top of the design hierarchy
12
13    -flatten
14        flatten the design before synthesis. this will pass '-auto-top' to
15        'hierarchy' if no top module is specified.

```



## APPENDIX C. COMMAND REFERENCE MANUAL

```
16
17 -ifx
18     passed to 'proc'. uses verilog simulation behavior for verilog if/case
19     undef handling. this also prevents 'wreduce' from being run.
20
21 -memx
22     simulate verilog simulation behavior for out-of-bounds memory accesses
23     using the 'memory_memx' pass.
24
25 -nomem
26     do not run any of the memory_* passes
27
28 -rdff
29     call 'memory_dff'. This enables merging of FFs into
30     memory read ports.
31
32 -nokeepdc
33     do not call opt_* with -keepdc
34
35 -run <from_label>[:<to_label>]
36     only run the commands between the labels (see below). an empty
37     from label is synonymous to 'begin', and empty to label is
38     synonymous to the end of the command list.
39
40
```

The following commands are executed by this synthesis command:

```
41
42
43 begin:
44     hierarchy -check [-top <top> | -auto-top]
45
46 coarse:
47     proc [-ifx]
48     flatten      (if -flatten)
49     opt_expr -keepdc
50     opt_clean
51     check
52     opt -noff -keepdc
53     wreduce -keepdc [-memx]
54     memory_dff      (if -rdff)
55     memory_memx      (if -memx)
56     opt_clean
57     memory_collect
58     opt -noff -keepdc -fast
59
60 check:
61     stat
62     check
```

### C.145 printattrs – print attributes of selected objects

```
1 printattrs [selection]
2
```

3 | Print all attributes of the selected objects.

## C.146 `proc` – translate processes to netlists

```

1   proc [options] [selection]
2
3   This pass calls all the other proc_* passes in the most common order.
4
5   proc_clean
6   proc_rmdead
7   proc_prune
8   proc_init
9   proc_arst
10  proc_rom
11  proc_mux
12  proc_dlatch
13  proc_dff
14  proc_memwr
15  proc_clean
16  opt_expr -keepdc
17
18  This replaces the processes in the design with multiplexers,
19  flip-flops and latches.
20
21  The following options are supported:
22
23  -nomux
24      Will omit the proc_mux pass.
25
26  -norom
27      Will omit the proc_rom pass.
28
29  -global_arst [!]<netname>
30      This option is passed through to proc_arst.
31
32  -ifx
33      This option is passed through to proc_mux. proc_rmdead is not
34      executed in -ifx mode.
35
36  -noopt
37      Will omit the opt_expr pass.
```

## C.147 `proc__arst` – detect asynchronous resets

```

1   proc_arst [-global_arst [!]<netname>] [selection]
2
3   This pass identifies asynchronous resets in the processes and converts them
4   to a different internal representation that is suitable for generating
5   flip-flop cells with asynchronous resets.
6
```

```

7  -global_arst [!]<netname>
8      In modules that have a net with the given name, use this net as async
9      reset for registers that have been assign initial values in their
10     declaration ('reg foobar = constant_value;'). Use the '!' modifier for
11     active low reset signals. Note: the frontend stores the default value
12     in the 'init' attribute on the net.

```

### C.148 `proc_clean` – remove empty parts of processes

```

1  proc_clean [options] [selection]
2
3  -quiet
4      do not print any messages.
5
6  This pass removes empty parts of processes and ultimately removes a process
7  if it contains only empty structures.

```

### C.149 `proc_dff` – extract flip-flops from processes

```

1  proc_dff [selection]
2
3  This pass identifies flip-flops in the processes and converts them to
4  d-type flip-flop cells.

```

### C.150 `proc_dlatch` – extract latches from processes

```

1  proc_dlatch [selection]
2
3  This pass identifies latches in the processes and converts them to
4  d-type latches.

```

### C.151 `proc_init` – convert initial block to init attributes

```

1  proc_init [selection]
2
3  This pass extracts the 'init' actions from processes (generated from Verilog
4  'initial' blocks) and sets the initial value to the 'init' attribute on the
5  respective wire.

```

### C.152 `proc_memwr` – extract memory writes from processes

```

1  proc_memwr [selection]
2
3  This pass converts memory writes in processes into $memwr cells.

```

**C.153 proc\_mux – convert decision trees to multiplexers**

```

1  proc_mux [options] [selection]
2
3  This pass converts the decision trees in processes (originating from if-else
4  and case statements) to trees of multiplexer cells.
5
6  -ifx
7      Use Verilog simulation behavior with respect to undef values in
8      'case' expressions and 'if' conditions.
```

**C.154 proc\_prune – remove redundant assignments**

```

1  proc_prune [selection]
2
3  This pass identifies assignments in processes that are always overwritten by
4  a later assignment to the same signal and removes them.
```

**C.155 proc\_rmdead – eliminate dead trees in decision trees**

```

1  proc_rmdead [selection]
2
3  This pass identifies unreachable branches in decision trees and removes them.
```

**C.156 proc\_rom – convert switches to ROMs**

```

1  proc_rom [selection]
2
3  This pass converts switches into read-only memories when appropriate.
```

**C.157 qbfsat – solve a 2QBF-SAT problem in the circuit**

```

1  qbfsat [options] [selection]
2
3  This command solves an "exists-forall" 2QBF-SAT problem defined over the currently
4  selected module. Existentially-quantified variables are declared by assigning a wire
5  "$anyconst". Universally-quantified variables may be explicitly declared by assigning
6  a wire "$allconst", but module inputs will be treated as universally-quantified
7  variables by default.
8
9  -nocleanup
10     Do not delete temporary files and directories. Useful for debugging.
11
12  -dump-final-smt2 <file>
```

## APPENDIX C. COMMAND REFERENCE MANUAL

13       Pass the `--dump-smt2` option to `yosys-smtbmc`.  
14  
15       `-assume-outputs`  
16       Add an `"$assume"` cell for the conjunction of all one-bit module output wires.  
17  
18       `-assume-negative-polarity`  
19       When adding `$assume` cells for one-bit module output wires, assume they are  
20       negative polarity signals and should always be low, for example like the  
21       miters created with the ``miter`` command.  
22  
23       `-nootimize`  
24       Ignore `"\minimize"` and `"\maximize"` attributes, do not emit `"(maximize)"` or  
25       `"(minimize)"` in the SMT-LIBv2, and generally make no attempt to optimize anything.  
26  
27       `-nobisection`  
28       If a wire is marked with the `"\minimize"` or `"\maximize"` attribute, do not  
29       attempt to optimize that value with the default iterated solving and threshold  
30       bisection approach. Instead, have `yosys-smtbmc` emit a `"(minimize)"` or `"(maximize)"`  
31       command in the SMT-LIBv2 output and hope that the solver supports optimizing  
32       quantified bitvector problems.  
33  
34       `-solver <solver>`  
35       Use a particular solver. Choose one of: `"z3"`, `"yices"`, and `"cvc4"`.  
36       (default: `yices`)  
37  
38       `-solver-option <name> <value>`  
39       Set the specified solver option in the SMT-LIBv2 problem file.  
40  
41       `-timeout <value>`  
42       Set the per-iteration timeout in seconds.  
43       (default: no timeout)  
44  
45       `-00, -01, -02`  
46       Control the use of ABC to simplify the QBF-SAT problem before solving.  
47  
48       `-sat`  
49       Generate an error if the solver does not return `"sat"`.  
50  
51       `-unsat`  
52       Generate an error if the solver does not return `"unsat"`.  
53  
54       `-show-smtbmc`  
55       Print the output from `yosys-smtbmc`.  
56  
57       `-specialize`  
58       If the problem is satisfiable, replace each `"$anyconst"` cell with its  
59       corresponding constant value from the model produced by the solver.  
60  
61       `-specialize-from-file <solution file>`  
62       Do not run the solver, but instead only attempt to replace each `"$anyconst"`  
63       cell in the current module with a constant value provided by the specified file.  
64  
65       `-write-solution <solution file>`  
66       If the problem is satisfiable, write the corresponding constant value for each

67 | "\$anyconst" cell from the model produced by the solver to the specified file.

## C.158 qwp – quadratic wirelength placer

```

1   qwp [options] [selection]
2
3   This command runs quadratic wirelength placement on the selected modules and
4   annotates the cells in the design with 'qwp_position' attributes.
5
6   -ltr
7       Add left-to-right constraints: constrain all inputs on the left border
8       outputs to the right border.
9
10  -alpha
11      Add constraints for inputs/outputs to be placed in alphanumerical
12      order along the y-axis (top-to-bottom).
13
14  -grid N
15      Number of grid divisions in x- and y-direction. (default=16)
16
17  -dump <html_file_name>
18      Dump a protocol of the placement algorithm to the html file.
19
20  -v
21      Verbose solver output for profiling or debugging
22
23  Note: This implementation of a quadratic wirelength placer uses exact
24  dense matrix operations. It is only a toy-placer for small circuits.
```

## C.159 read – load HDL designs

```

1   read {-vlog95|-vlog2k|-sv2005|-sv2009|-sv2012|-sv|-formal} <verilog-file>..
2
3   Load the specified Verilog/SystemVerilog files. (Full SystemVerilog support
4   is only available via Verific.)
5
6   Additional -D<macro>[=<value>] options may be added after the option indicating
7   the language version (and before file names) to set additional verilog defines.
8
9
10  read {-vhd187|-vhd193|-vhd12k|-vhd12008|-vhd1} <vhdl-file>..
11
12  Load the specified VHDL files. (Requires Verific.)
13
14
15  read {-f|-F} <command-file>
16
17  Load and execute the specified command file. (Requires Verific.)
18  Check verific command for more information about supported commands in file.
19
```

```

20
21     read -define <macro>[=<value>]..
22
23 Set global Verilog/SystemVerilog defines.
24
25
26     read -undef <macro>..
27
28 Unset global Verilog/SystemVerilog defines.
29
30
31     read -incdir <directory>
32
33 Add directory to global Verilog/SystemVerilog include directories.
34
35
36     read -verific
37     read -noverific
38
39 Subsequent calls to 'read' will either use or not use Verific. Calling 'read'
40 with -verific will result in an error on Yosys binaries that are built without
41 Verific support. The default is to use Verific if it is available.

```

## C.160 read\_aiger – read AIGER file

```

1     read_aiger [options] [filename]
2
3 Load module from an AIGER file into the current design.
4
5     -module_name <module_name>
6         name of module to be created (default: <filename>)
7
8     -clk_name <wire_name>
9         if specified, AIGER latches to be transformed into $_DFF_P_ cells
10        clocked by wire of this name. otherwise, $_FF_ cells will be used
11
12     -map <filename>
13        read file with port and latch symbols
14
15     -wideports
16        merge ports that match the pattern 'name[int]' into a single
17        multi-bit port 'name'
18
19     -xaiger
20        read XAIGER extensions

```

## C.161 read\_blif – read BLIF file

```

1     read_blif [options] [filename]
2

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
3 Load modules from a BLIF file into the current design.
4
5     -sop
6         Create $sop cells instead of $lut cells
7
8     -wideports
9         Merge ports that match the pattern 'name[int]' into a single
10        multi-bit port 'name'.
```

### C.162 read\_ilang – (deprecated) alias of read\_rtlil

```
1 See `help read_rtlil`.
```

### C.163 read\_json – read JSON file

```
1     read_json [filename]
2
3 Load modules from a JSON file into the current design See "help write_json"
4 for a description of the file format.
```

### C.164 read\_liberty – read cells from liberty file

```
1     read_liberty [filename]
2
3 Read cells from liberty file as modules into current design.
4
5     -lib
6         only create empty blackbox modules
7
8     -nooverwrite
9         ignore re-definitions of modules. (the default behavior is to
10        create an error message if the existing module is not a blackbox
11        module, and overwrite the existing module if it is a blackbox module.)
12
13     -overwrite
14         overwrite existing modules with the same name
15
16     -ignore_miss_func
17         ignore cells with missing function specification of outputs
18
19     -ignore_miss_dir
20         ignore cells with a missing or invalid direction
21         specification on a pin
22
23     -ignore_miss_data_latch
24         ignore latches with missing data and/or enable pins
25
```



```

26 -setattr <attribute_name>
27     set the specified attribute (to the value 1) on all loaded modules

```

## C.165 read\_rtlil – read modules from RTLIL file

```

1  read_rtlil [filename]
2
3  Load modules from an RTLIL file to the current design. (RTLIL is a text
4  representation of a design in yosys's internal format.)
5
6  -nooverwrite
7      ignore re-definitions of modules. (the default behavior is to
8      create an error message if the existing module is not a blackbox
9      module, and overwrite the existing module if it is a blackbox module.)
10
11 -overwrite
12     overwrite existing modules with the same name
13
14 -lib
15     only create empty blackbox modules

```

## C.166 read\_verilog – read modules from Verilog file

```

1  read_verilog [options] [filename]
2
3  Load modules from a Verilog file to the current design. A large subset of
4  Verilog-2005 is supported.
5
6  -sv
7      enable support for SystemVerilog features. (only a small subset
8      of SystemVerilog is supported)
9
10 -formal
11     enable support for SystemVerilog assertions and some Yosys extensions
12     replace the implicit -D SYNTHESIS with -D FORMAL
13
14 -nosynthesis
15     don't add implicit -D SYNTHESIS
16
17 -noassert
18     ignore assert() statements
19
20 -noassume
21     ignore assume() statements
22
23 -norestrict
24     ignore restrict() statements
25
26 -assume-asserts
27     treat all assert() statements like assume() statements

```

## APPENDIX C. COMMAND REFERENCE MANUAL

28  
29     `-assert-assumes`  
30         treat all `assume()` statements like `assert()` statements  
31  
32     `-debug`  
33         alias for `-dump_ast1 -dump_ast2 -dump_vlog1 -dump_vlog2 -yydebug`  
34  
35     `-dump_ast1`  
36         dump abstract syntax tree (before simplification)  
37  
38     `-dump_ast2`  
39         dump abstract syntax tree (after simplification)  
40  
41     `-no_dump_ptr`  
42         do not include hex memory addresses in dump (easier to diff dumps)  
43  
44     `-dump_vlog1`  
45         dump ast as Verilog code (before simplification)  
46  
47     `-dump_vlog2`  
48         dump ast as Verilog code (after simplification)  
49  
50     `-dump_rtlil`  
51         dump generated RTLIL netlist  
52  
53     `-yydebug`  
54         enable parser debug output  
55  
56     `-nolatches`  
57         usually latches are synthesized into logic loops  
58         this option prohibits this and sets the output to 'x'  
59         in what would be the latches hold condition  
60  
61         this behavior can also be achieved by setting the  
62         'nolatches' attribute on the respective module or  
63         always block.  
64  
65     `-nomem2reg`  
66         under certain conditions memories are converted to registers  
67         early during simplification to ensure correct handling of  
68         complex corner cases. this option disables this behavior.  
69  
70         this can also be achieved by setting the 'nomem2reg'  
71         attribute on the respective module or register.  
72  
73         This is potentially dangerous. Usually the front-end has good  
74         reasons for converting an array to a list of registers.  
75         Prohibiting this step will likely result in incorrect synthesis  
76         results.  
77  
78     `-mem2reg`  
79         always convert memories to registers. this can also be  
80         achieved by setting the 'mem2reg' attribute on the respective  
81         module or register.

## APPENDIX C. COMMAND REFERENCE MANUAL

```
82
83 -nomeminit
84     do not infer $meminit cells and instead convert initialized
85     memories to registers directly in the front-end.
86
87 -ppdump
88     dump Verilog code after pre-processor
89
90 -nopp
91     do not run the pre-processor
92
93 -nodpi
94     disable DPI-C support
95
96 -noblackbox
97     do not automatically add a (* blackbox *) attribute to an
98     empty module.
99
100 -lib
101     only create empty blackbox modules. This implies -DBLACKBOX.
102     modules with the (* whitebox *) attribute will be preserved.
103     (* lib_whitebox *) will be treated like (* whitebox *).
104
105 -nowb
106     delete (* whitebox *) and (* lib_whitebox *) attributes from
107     all modules.
108
109 -specify
110     parse and import specify blocks
111
112 -noopt
113     don't perform basic optimizations (such as const folding) in the
114     high-level front-end.
115
116 -icells
117     interpret cell types starting with '$' as internal cell types
118
119 -pwires
120     add a wire for each module parameter
121
122 -nooverwrite
123     ignore re-definitions of modules. (the default behavior is to
124     create an error message if the existing module is not a black box
125     module, and overwrite the existing module otherwise.)
126
127 -overwrite
128     overwrite existing modules with the same name
129
130 -defer
131     only read the abstract syntax tree and defer actual compilation
132     to a later 'hierarchy' command. Useful in cases where the default
133     parameters of modules yield invalid or not synthesizable code.
134
135 -noautowire
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
136     make the default of `default_nettype be "none" instead of "wire".
137
138     -setattr <attribute_name>
139         set the specified attribute (to the value 1) on all loaded modules
140
141     -Dname[=definition]
142         define the preprocessor symbol 'name' and set its optional value
143         'definition'
144
145     -Idir
146         add 'dir' to the directories which are used when searching include
147         files
148
149 The command 'verilog_defaults' can be used to register default options for
150 subsequent calls to 'read_verilog'.
151
152 Note that the Verilog frontend does a pretty good job of processing valid
153 verilog input, but has not very good error reporting. It generally is
154 recommended to use a simulator (for example Icarus Verilog) for checking
155 the syntax of the code, rather than to rely on read_verilog for that.
156
157 Depending on if read_verilog is run in -formal mode, either the macro
158 SYNTHESIS or FORMAL is defined automatically, unless -nosynthesis is used.
159 In addition, read_verilog always defines the macro YOSYS.
160
161 See the Yosys README file for a list of non-standard Verilog features
162 supported by the Yosys Verilog front-end.
```

### C.167 rename – rename object in the design

```
1     rename old_name new_name
2
3 Rename the specified object. Note that selection patterns are not supported
4 by this command.
5
6
7
8     rename -output old_name new_name
9
10 Like above, but also make the wire an output. This will fail if the object is
11 not a wire.
12
13
14     rename -src [selection]
15
16 Assign names auto-generated from the src attribute to all selected wires and
17 cells with private names.
18
19
20     rename -wire [selection] [-suffix <suffix>]
21
22 Assign auto-generated names based on the wires they drive to all selected
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
23 cells with private names. Ignores cells driving privately named wires.
24 By default, the cell is named after the wire with the cell type as suffix.
25 The -suffix option can be used to set the suffix to the given string instead.
26
27
28     rename -enumerate [-pattern <pattern>] [selection]
29
30 Assign short auto-generated names to all selected wires and cells with private
31 names. The -pattern option can be used to set the pattern for the new names.
32 The character % in the pattern is replaced with a integer number. The default
33 pattern is '%_'.
34
35
36     rename -hide [selection]
37
38 Assign private names (the ones with $-prefix) to all selected wires and cells
39 with public names. This ignores all selected ports.
40
41
42     rename -top new_name
43
44 Rename top module.
```

### C.168 rmports – remove module ports with no connections

```
1     rmports [selection]
2
3 This pass identifies ports in the selected modules which are not used or
4 driven and removes them.
```

### C.169 sat – solve a SAT problem in the circuit

```
1     sat [options] [selection]
2
3 This command solves a SAT problem defined over the currently selected circuit
4 and additional constraints passed as parameters.
5
6     -all
7         show all solutions to the problem (this can grow exponentially, use
8         -max <N> instead to get <N> solutions)
9
10    -max <N>
11        like -all, but limit number of solutions to <N>
12
13    -enable_undef
14        enable modeling of undef value (aka 'x-bits')
15        this option is implied by -set-def, -set-undef et. cetera
16
17    -max_undef
18        maximize the number of undef bits in solutions, giving a better
```

## APPENDIX C. COMMAND REFERENCE MANUAL

19 picture of which input bits are actually vital to the solution.  
20  
21 `-set <signal> <value>`  
22 set the specified signal to the specified value.  
23  
24 `-set-def <signal>`  
25 add a constraint that all bits of the given signal must be defined  
26  
27 `-set-any-undef <signal>`  
28 add a constraint that at least one bit of the given signal is undefined  
29  
30 `-set-all-undef <signal>`  
31 add a constraint that all bits of the given signal are undefined  
32  
33 `-set-def-inputs`  
34 add `-set-def` constraints for all module inputs  
35  
36 `-show <signal>`  
37 show the model for the specified signal. if no `-show` option is  
38 passed then a set of signals to be shown is automatically selected.  
39  
40 `-show-inputs, -show-outputs, -show-ports`  
41 add all module (input/output) ports to the list of shown signals  
42  
43 `-show-regs, -show-public, -show-all`  
44 show all registers, show signals with 'public' names, show all signals  
45  
46 `-ignore_div_by_zero`  
47 ignore all solutions that involve a division by zero  
48  
49 `-ignore_unknown_cells`  
50 ignore all cells that can not be matched to a SAT model  
51  
52 The following options can be used to set up a sequential problem:  
53  
54 `-seq <N>`  
55 set up a sequential problem with `<N>` time steps. The steps will  
56 be numbered from 1 to `N`.  
57  
58 note: for large `<N>` it can be significantly faster to use  
59 `-tempinduct-baseonly -maxsteps <N>` instead of `-seq <N>`.  
60  
61 `-set-at <N> <signal> <value>`  
62 `-unset-at <N> <signal>`  
63 set or unset the specified signal to the specified value in the  
64 given timestep. this has priority over a `-set` for the same signal.  
65  
66 `-set-assumes`  
67 set all assumptions provided via `$assume` cells  
68  
69 `-set-def-at <N> <signal>`  
70 `-set-any-undef-at <N> <signal>`  
71 `-set-all-undef-at <N> <signal>`  
72 add undef constraints in the given timestep.

## APPENDIX C. COMMAND REFERENCE MANUAL

```

73
74 -set-init <signal> <value>
75     set the initial value for the register driving the signal to the value
76
77 -set-init-undef
78     set all initial states (not set using -set-init) to undef
79
80 -set-init-def
81     do not force a value for the initial state but do not allow undef
82
83 -set-init-zero
84     set all initial states (not set using -set-init) to zero
85
86 -dump_vcd <vcd-file-name>
87     dump SAT model (counter example in proof) to VCD file
88
89 -dump_json <json-file-name>
90     dump SAT model (counter example in proof) to a WaveJSON file.
91
92 -dump_cnf <cnf-file-name>
93     dump CNF of SAT problem (in DIMACS format). in temporal induction
94     proofs this is the CNF of the first induction step.
95
96 The following additional options can be used to set up a proof. If also -seq
97 is passed, a temporal induction proof is performed.
98
99 -tempinduct
100     Perform a temporal induction proof. In a temporal induction proof it is
101     proven that the condition holds forever after the number of time steps
102     specified using -seq.
103
104 -tempinduct-def
105     Perform a temporal induction proof. Assume an initial state with all
106     registers set to defined values for the induction step.
107
108 -tempinduct-baseonly
109     Run only the basecase half of temporal induction (requires -maxsteps)
110
111 -tempinduct-inductonly
112     Run only the induction half of temporal induction
113
114 -tempinduct-skip <N>
115     Skip the first <N> steps of the induction proof.
116
117     note: this will assume that the base case holds for <N> steps.
118     this must be proven independently with "-tempinduct-baseonly
119     -maxsteps <N>". Use -initsteps if you just want to set a
120     minimal induction length.
121
122 -prove <signal> <value>
123     Attempt to proof that <signal> is always <value>.
124
125 -prove-x <signal> <value>
126     Like -prove, but an undef (x) bit in the lhs matches any value on

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
127         the right hand side. Useful for equivalence checking.
128
129     -prove-asserts
130         Prove that all asserts in the design hold.
131
132     -prove-skip <N>
133         Do not enforce the prove-condition for the first <N> time steps.
134
135     -maxsteps <N>
136         Set a maximum length for the induction.
137
138     -initsteps <N>
139         Set initial length for the induction.
140         This will speed up the search of the right induction length
141         for deep induction proofs.
142
143     -stepsize <N>
144         Increase the size of the induction proof in steps of <N>.
145         This will speed up the search of the right induction length
146         for deep induction proofs.
147
148     -timeout <N>
149         Maximum number of seconds a single SAT instance may take.
150
151     -verify
152         Return an error and stop the synthesis script if the proof fails.
153
154     -verify-no-timeout
155         Like -verify but do not return an error for timeouts.
156
157     -falsify
158         Return an error and stop the synthesis script if the proof succeeds.
159
160     -falsify-no-timeout
161         Like -falsify but do not return an error for timeouts.
```

### C.170 scatter – add additional intermediate nets

```
1     scatter [selection]
2
3     This command adds additional intermediate nets on all cell ports. This is used
4     for testing the correct use of the SigMap helper in passes. If you don't know
5     what this means: don't worry -- you only need this pass when testing your own
6     extensions to Yosys.
7
8     Use the opt_clean command to get rid of the additional nets.
```

### C.171 scc – detect strongly connected components (logic loops)



## APPENDIX C. COMMAND REFERENCE MANUAL

```
1  scc [options] [selection]
2
3  This command identifies strongly connected components (aka logic loops) in the
4  design.
5
6  -expect <num>
7      expect to find exactly <num> SCCs. A different number of SCCs will
8      produce an error.
9
10 -max_depth <num>
11     limit to loops not longer than the specified number of cells. This
12     can e.g. be useful in identifying small local loops in a module that
13     implements one large SCC.
14
15 -nofeedback
16     do not count cells that have their output fed back into one of their
17     inputs as single-cell scc.
18
19 -all_cell_types
20     Usually this command only considers internal non-memory cells. With
21     this option set, all cells are considered. For unknown cells all ports
22     are assumed to be bidirectional 'inout' ports.
23
24 -set_attr <name> <value>
25     set the specified attribute on all cells that are part of a logic
26     loop. the special token {} in the value is replaced with a unique
27     identifier for the logic loop.
28
29 -select
30     replace the current selection with a selection of all cells and wires
31     that are part of a found logic loop
32
33 -specify
34     examine specify rules to detect logic loops in whitebox/blackbox cells
```

### C.172 scratchpad – get/set values in the scratchpad

```
1  scratchpad [options]
2
3  This pass allows to read and modify values from the scratchpad of the current
4  design. Options:
5
6  -get <identifier>
7      print the value saved in the scratchpad under the given identifier.
8
9  -set <identifier> <value>
10     save the given value in the scratchpad under the given identifier.
11
12 -unset <identifier>
13     remove the entry for the given identifier from the scratchpad.
14
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
15  -copy <identifier_from> <identifier_to>
16      copy the value of the first identifier to the second identifier.
17
18  -assert <identifier> <value>
19      assert that the entry for the given identifier is set to the given value.
20
21  -assert-set <identifier>
22      assert that the entry for the given identifier exists.
23
24  -assert-unset <identifier>
25      assert that the entry for the given identifier does not exist.
26
27  The identifier may not contain whitespace. By convention, it is usually prefixed
28  by the name of the pass that uses it, e.g. 'opt.did_something'. If the value
29  contains whitespace, it must be enclosed in double quotes.
```

### C.173 script – execute commands from file or wire

```
1  script <filename> [<from_label>:<to_label>]
2  script -scriptwire [selection]
3
4  This command executes the yosys commands in the specified file (default
5  behaviour), or commands embedded in the constant text value connected to the
6  selected wires.
7
8  In the default (file) case, the 2nd argument can be used to only execute the
9  section of the file between the specified labels. An empty from label is
10 synonymous with the beginning of the file and an empty to label is synonymous
11 with the end of the file.
12
13 If only one label is specified (without ':') then only the block
14 marked with that label (until the next label) is executed.
15
16 In "-scriptwire" mode, the commands on the selected wire(s) will be executed
17 in the scope of (and thus, relative to) the wires' owning module(s). This
18 '-module' mode can be exited by using the 'cd' command.
```

### C.174 select – modify and view the list of selected objects

```
1  select [ -add | -del | -set <name> ] {-read <filename> | <selection>}
2  select [ -unset <name> ]
3  select [ <assert_option> ] {-read <filename> | <selection>}
4  select [ -list | -write <filename> | -count | -clear ]
5  select -module <modname>
6
7  Most commands use the list of currently selected objects to determine which part
8  of the design to operate on. This command can be used to modify and view this
9  list of selected objects.
10
11 Note that many commands support an optional [selection] argument that can be
```

## APPENDIX C. COMMAND REFERENCE MANUAL

used to override the global selection for the command. The syntax of this optional argument is identical to the syntax of the <selection> argument described here.

**-add, -del**  
add or remove the given objects to the current selection.  
without this options the current selection is replaced.

**-set <name>**  
do not modify the current selection. instead save the new selection under the given name (see @<name> below). to save the current selection, use "select -set <name> %"

**-unset <name>**  
do not modify the current selection. instead remove a previously saved selection under the given name (see @<name> below).

**-assert-none**  
do not modify the current selection. instead assert that the given selection is empty. i.e. produce an error if any object or module matching the selection is found.

**-assert-any**  
do not modify the current selection. instead assert that the given selection is non-empty. i.e. produce an error if no object or module matching the selection is found.

**-assert-count N**  
do not modify the current selection. instead assert that the given selection contains exactly N objects.

**-assert-max N**  
do not modify the current selection. instead assert that the given selection contains less than or exactly N objects.

**-assert-min N**  
do not modify the current selection. instead assert that the given selection contains at least N objects.

**-list**  
list all objects in the current selection

**-write <filename>**  
like -list but write the output to the specified file

**-read <filename>**  
read the specified file (written by -write)

**-count**  
count all objects in the current selection

**-clear**  
clear the current selection. this effectively selects the whole design. it also resets the selected module (see -module). use the

## APPENDIX C. COMMAND REFERENCE MANUAL

66       command 'select \*' to select everything but stay in the current module.

67

68       -none

69       create an empty selection. the current module is unchanged.

70

71       -module <modname>

72       limit the current scope to the specified module.

73       the difference between this and simply selecting the module

74       is that all object names are interpreted relative to this

75       module after this command until the selection is cleared again.

76

77       When this command is called without an argument, the current selection

78       is displayed in a compact form (i.e. only the module name when a whole module

79       is selected).

80

81       The <selection> argument itself is a series of commands for a simple stack

82       machine. Each element on the stack represents a set of selected objects.

83       After this commands have been executed, the union of all remaining sets

84       on the stack is computed and used as selection for the command.

85

86       Pushing (selecting) object when not in -module mode:

87

88       <mod\_pattern>

89       select the specified module(s)

90

91       <mod\_pattern>/<obj\_pattern>

92       select the specified object(s) from the module(s)

93

94       Pushing (selecting) object when in -module mode:

95

96       <obj\_pattern>

97       select the specified object(s) from the current module

98

99       By default, patterns will not match black/white-box modules or their contents. To include such objects, pref

100

101       A <mod\_pattern> can be a module name, wildcard expression (\*, ?, [...])

102       matching module names, or one of the following:

103

104       A:<pattern>, A:<pattern>=<pattern>

105       all modules with an attribute matching the given pattern

106       in addition to = also <, <=, >=, and > are supported

107

108       N:<pattern>

109       all modules with a name matching the given pattern

110       (i.e. 'N:' is optional as it is the default matching rule)

111

112       An <obj\_pattern> can be an object name, wildcard expression, or one of

113       the following:

114

115       w:<pattern>

116       all wires with a name matching the given wildcard pattern

117

118       i:<pattern>, o:<pattern>, x:<pattern>

119       all inputs (i:), outputs (o:) or any ports (x:) with matching names

## APPENDIX C. COMMAND REFERENCE MANUAL

```

120
121 s:<size>, s:<min>:<max>
122     all wires with a matching width
123
124 m:<pattern>
125     all memories with a name matching the given pattern
126
127 c:<pattern>
128     all cells with a name matching the given pattern
129
130 t:<pattern>
131     all cells with a type matching the given pattern
132
133 p:<pattern>
134     all processes with a name matching the given pattern
135
136 a:<pattern>
137     all objects with an attribute name matching the given pattern
138
139 a:<pattern>=<pattern>
140     all objects with a matching attribute name-value-pair.
141     in addition to = also <, <=, >=, and > are supported
142
143 r:<pattern>, r:<pattern>=<pattern>
144     cells with matching parameters. also with <, <=, >= and >.
145
146 n:<pattern>
147     all objects with a name matching the given pattern
148     (i.e. 'n:' is optional as it is the default matching rule)
149
150 @<name>
151     push the selection saved prior with 'select -set <name> ...'
152
153 The following actions can be performed on the top sets on the stack:
154
155 %
156     push a copy of the current selection to the stack
157
158 %%
159     replace the stack with a union of all elements on it
160
161 %n
162     replace top set with its invert
163
164 %u
165     replace the two top sets on the stack with their union
166
167 %i
168     replace the two top sets on the stack with their intersection
169
170 %d
171     pop the top set from the stack and subtract it from the new top
172
173 %D

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

174         like %d but swap the roles of two top sets on the stack
175
176     %c
177         create a copy of the top set from the stack and push it
178
179     %x[<num1>|*][.<num2>][:<rule>[:<rule>..]]
180         expand top set <num1> num times according to the specified rules.
181         (i.e. select all cells connected to selected wires and select all
182         wires connected to selected cells) The rules specify which cell
183         ports to use for this. the syntax for a rule is a '-' for exclusion
184         and a '+' for inclusion, followed by an optional comma separated
185         list of cell types followed by an optional comma separated list of
186         cell ports in square brackets. a rule can also be just a cell or wire
187         name that limits the expansion (is included but does not go beyond).
188         select at most <num2> objects. a warning message is printed when this
189         limit is reached. When '*' is used instead of <num1> then the process
190         is repeated until no further object are selected.
191
192     %ci[<num1>|*][.<num2>][:<rule>[:<rule>..]]
193     %co[<num1>|*][.<num2>][:<rule>[:<rule>..]]
194         similar to %x, but only select input (%ci) or output cones (%co)
195
196     %xe[...] %cie[...] %coe
197         like %x, %ci, and %co but only consider combinatorial cells
198
199     %a
200         expand top set by selecting all wires that are (at least in part)
201         aliases for selected wires.
202
203     %s
204         expand top set by adding all modules that implement cells in selected
205         modules
206
207     %m
208         expand top set by selecting all modules that contain selected objects
209
210     %M
211         select modules that implement selected cells
212
213     %C
214         select cells that implement selected modules
215
216     %R[<num>]
217         select <num> random objects from top selection (default 1)
218
219 Example: the following command selects all wires that are connected to a
220 'GATE' input of a 'SWITCH' cell:
221
222     select */t:SWITCH %x:+[GATE] */t:SWITCH %d

```

### C.175 setattr – set/unset attributes on objects

```

1      setattr [ -mod ] [ -set name value | -unset name ]... [selection]
2
3  Set/unset the given attributes on the selected objects. String values must be
4  passed in double quotes (").
5
6  When called with -mod, this command will set and unset attributes on modules
7  instead of objects within modules.
```

## C.176 setparam – set/unset parameters on objects

```

1      setparam [ -type cell_type ] [ -set name value | -unset name ]... [selection]
2
3  Set/unset the given parameters on the selected cells. String values must be
4  passed in double quotes (").
5
6  The -type option can be used to change the cell type of the selected cells.
```

## C.177 setundef – replace undef values with defined constants

```

1      setundef [options] [selection]
2
3  This command replaces undef (x) constants with defined (0/1) constants.
4
5  -undriven
6      also set undriven nets to constant values
7
8  -expose
9      also expose undriven nets as inputs (use with -undriven)
10
11 -zero
12     replace with bits cleared (0)
13
14 -one
15     replace with bits set (1)
16
17 -undef
18     replace with undef (x) bits, may be used with -undriven
19
20 -anyseq
21     replace with $anyseq drivers (for formal)
22
23 -anyconst
24     replace with $anyconst drivers (for formal)
25
26 -random <seed>
27     replace with random bits using the specified integer as seed
28     value for the random number generator.
29
30 -init
```

```

31     also create/update init values for flip-flops
32
33     -params
34         replace undef in cell parameters

```

## C.178 share – perform sat-based resource sharing

```

1     share [options] [selection]
2
3     This pass merges shareable resources into a single resource. A SAT solver
4     is used to determine if two resources are share-able.
5
6     -force
7         Per default the selection of cells that is considered for sharing is
8         narrowed using a list of cell types. With this option all selected
9         cells are considered for resource sharing.
10
11     IMPORTANT NOTE: If the -all option is used then no cells with internal
12     state must be selected!
13
14     -aggressive
15         Per default some heuristics are used to reduce the number of cells
16         considered for resource sharing to only large resources. This options
17         turns this heuristics off, resulting in much more cells being considered
18         for resource sharing.
19
20     -fast
21         Only consider the simple part of the control logic in SAT solving, resulting
22         in much easier SAT problems at the cost of maybe missing some opportunities
23         for resource sharing.
24
25     -limit N
26         Only perform the first N merges, then stop. This is useful for debugging.

```

## C.179 shell – enter interactive command mode

```

1     shell
2
3     This command enters the interactive command mode. This can be useful
4     in a script to interrupt the script at a certain point and allow for
5     interactive inspection or manual synthesis of the design at this point.
6
7     The command prompt of the interactive shell indicates the current
8     selection (see 'help select'):
9
10     yosys>
11         the entire design is selected
12
13     yosys*>
14         only part of the design is selected

```



```

15
16     yosys [modname]>
17         the entire module 'modname' is selected using 'select -module modname'
18
19     yosys [modname]*>
20         only part of current module 'modname' is selected
21
22 When in interactive shell, some errors (e.g. invalid command arguments)
23 do not terminate yosys but return to the command prompt.
24
25 This command is the default action if nothing else has been specified
26 on the command line.
27
28 Press Ctrl-D or type 'exit' to leave the interactive shell.

```

## C.180 show – generate schematics using graphviz

```

1     show [options] [selection]
2
3 Create a graphviz DOT file for the selected part of the design and compile it
4 to a graphics file (usually SVG or PostScript).
5
6     -viewer <viewer>
7         Run the specified command with the graphics file as parameter.
8         On Windows, this pauses yosys until the viewer exits.
9
10    -format <format>
11        Generate a graphics file in the specified format. Use 'dot' to just
12        generate a .dot file, or other <format> strings such as 'svg' or 'ps'
13        to generate files in other formats (this calls the 'dot' command).
14
15    -lib <verilog_or_rtlil_file>
16        Use the specified library file for determining whether cell ports are
17        inputs or outputs. This option can be used multiple times to specify
18        more than one library.
19
20        note: in most cases it is better to load the library before calling
21        show with 'read_verilog -lib <filename>'. it is also possible to
22        load liberty files with 'read_liberty -lib <filename>'.
23
24    -prefix <prefix>
25        generate <prefix>.* instead of ~/.yosys_show.*
26
27    -color <color> <object>
28        assign the specified color to the specified object. The object can be
29        a single selection wildcard expressions or a saved set of objects in
30        the @<name> syntax (see "help select" for details).
31
32    -label <text> <object>
33        assign the specified label text to the specified object. The object can
34        be a single selection wildcard expressions or a saved set of objects in
35        the @<name> syntax (see "help select" for details).

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
36
37 -colors <seed>
38     Randomly assign colors to the wires. The integer argument is the seed
39     for the random number generator. Change the seed value if the colored
40     graph still is ambiguous. A seed of zero deactivates the coloring.
41
42 -colorattr <attribute_name>
43     Use the specified attribute to assign colors. A unique color is
44     assigned to each unique value of this attribute.
45
46 -width
47     annotate buses with a label indicating the width of the bus.
48
49 -signed
50     mark ports (A, B) that are declared as signed (using the [AB]_SIGNED
51     cell parameter) with an asterisk next to the port name.
52
53 -stretch
54     stretch the graph so all inputs are on the left side and all outputs
55     (including inout ports) are on the right side.
56
57 -pause
58     wait for the user to press enter to before returning
59
60 -enum
61     enumerate objects with internal ($-prefixed) names
62
63 -long
64     do not abbreviate objects with internal ($-prefixed) names
65
66 -notitle
67     do not add the module name as graph title to the dot file
68
69 -nobg
70     don't run viewer in the background, IE wait for the viewer tool to
71     exit before returning
72
73 When no <format> is specified, 'dot' is used. When no <format> and <viewer> is
74 specified, 'xdot' is used to display the schematic (POSIX systems only).
75
76 The generated output files are '~/yosys_show.dot' and '~/yosys_show.<format>',
77 unless another prefix is specified using -prefix <prefix>.
78
79 Yosys on Windows and YosysJS use different defaults: The output is written
80 to 'show.dot' in the current directory and new viewer is launched each time
81 the 'show' command is executed.
```

### C.181 shregmap – map shift registers

```
1 shregmap [options] [selection]
2
3 This pass converts chains of $_DFF_[NP]_ gates to target specific shift register
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
4 primitives. The generated shift register will be of type $__SHREG_DFF_[NP]_ and
5 will use the same interface as the original $_DFF*_ cells. The cell parameter
6 'DEPTH' will contain the depth of the shift register. Use a target-specific
7 'techmap' map file to convert those cells to the actual target cells.
8
9 -minlen N
10     minimum length of shift register (default = 2)
11     (this is the length after -keep_before and -keep_after)
12
13 -maxlen N
14     maximum length of shift register (default = no limit)
15     larger chains will be mapped to multiple shift register instances
16
17 -keep_before N
18     number of DFFs to keep before the shift register (default = 0)
19
20 -keep_after N
21     number of DFFs to keep after the shift register (default = 0)
22
23 -clkpol pos|neg|any
24     limit match to only positive or negative edge clocks. (default = any)
25
26 -enpol pos|neg|none|any_or_none|any
27     limit match to FFs with the specified enable polarity. (default = none)
28
29 -match <cell_type>[:<d_port_name>:<q_port_name>]
30     match the specified cells instead of $_DFF_N_ and $_DFF_P_. If
31     '[:<d_port_name>:<q_port_name>]' is omitted then 'D' and 'Q' is used
32     by default. E.g. the option '-clkpol pos' is just an alias for
33     '-match $_DFF_P_', which is an alias for '-match $_DFF_P_:D:Q'.
34
35 -params
36     instead of encoding the clock and enable polarity in the cell name by
37     deriving from the original cell name, simply name all generated cells
38     $__SHREG_ and use CLKPOL and ENPOL parameters. An ENPOL value of 2 is
39     used to denote cells without enable input. The ENPOL parameter is
40     omitted when '-enpol none' (or no -enpol option) is passed.
41
42 -zinit
43     assume the shift register is automatically zero-initialized, so it
44     becomes legal to merge zero initialized FFs into the shift register.
45
46 -init
47     map initialized registers to the shift reg, add an INIT parameter to
48     generated cells with the initialization value. (first bit to shift out
49     in LSB position)
50
51 -tech greenpak4
52     map to greenpak4 shift registers.
```

### C.182 sim – simulate the circuit

## APPENDIX C. COMMAND REFERENCE MANUAL

```
1  sim [options] [top-level]
2
3  This command simulates the circuit using the given top-level module.
4
5  -vcd <filename>
6      write the simulation results to the given VCD file
7
8  -fst <filename>
9      write the simulation results to the given FST file
10
11 -aiw <filename>
12     write the simulation results to an AIGER witness file
13     (requires a *.aim file via -map)
14
15 -x
16     ignore constant x outputs in simulation file.
17
18 -date
19     include date and full version info in output.
20
21 -clock <portname>
22     name of top-level clock input
23
24 -clockn <portname>
25     name of top-level clock input (inverse polarity)
26
27 -multiclock
28     mark that witness file is multiclock.
29
30 -reset <portname>
31     name of top-level reset input (active high)
32
33 -resetn <portname>
34     name of top-level inverted reset input (active low)
35
36 -rstlen <integer>
37     number of cycles reset should stay active (default: 1)
38
39 -zinit
40     zero-initialize all uninitialized regs and memories
41
42 -timescale <string>
43     include the specified timescale declaration in the vcd
44
45 -n <integer>
46     number of clock cycles to simulate (default: 20)
47
48 -a
49     use all nets in VCD/FST operations, not just those with public names
50
51 -w
52     writeback mode: use final simulation state as new init state
53
54 -r
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
55     read simulation results file (file formats supported: FST, VCD, AIW and WIT)
56         VCD support requires vcd2fst external tool to be present
57
58     -map <filename>
59         read file with port and latch symbols, needed for AIGER witness input
60
61     -scope <name>
62         scope of simulation top model
63
64     -at <time>
65         sets start and stop time
66
67     -start <time>
68         start co-simulation in arbitrary time (default 0)
69
70     -stop <time>
71         stop co-simulation in arbitrary time (default END)
72
73     -sim
74         simulation with stimulus from FST (default)
75
76     -sim-cmp
77         co-simulation expect exact match
78
79     -sim-gold
80         co-simulation, x in simulation can match any value in FST
81
82     -sim-gate
83         co-simulation, x in FST can match any value in simulation
84
85     -q
86         disable per-cycle/sample log message
87
88     -d
89         enable debug output
```

### C.183 simplemap – mapping simple coarse-grain cells

```
1     simplemap [selection]
2
3     This pass maps a small selection of simple coarse-grain cells to yosys gate
4     primitives. The following internal cell types are mapped by this pass:
5
6     $not, $pos, $and, $or, $xor, $xnor
7     $reduce_and, $reduce_or, $reduce_xor, $reduce_xnor, $reduce_bool
8     $logic_not, $logic_and, $logic_or, $mux, $tribuf
9     $sr, $ff, $dff, $dffe, $dffsr, $dffsre, $adff, $adffe, $aldff, $aldffe, $sdff, $sdffe, $sdffce, $dlatch,
```

### C.184 splice – create explicit splicing cells

## APPENDIX C. COMMAND REFERENCE MANUAL

```
1 splice [options] [selection]
2
3 This command adds $slice and $concat cells to the design to make the splicing
4 of multi-bit signals explicit. This for example is useful for coarse grain
5 synthesis, where dedicated hardware is needed to splice signals.
6
7 -sel_by_cell
8     only select the cell ports to rewire by the cell. if the selection
9     contains a cell, than all cell inputs are rewired, if necessary.
10
11 -sel_by_wire
12     only select the cell ports to rewire by the wire. if the selection
13     contains a wire, than all cell ports driven by this wire are wired,
14     if necessary.
15
16 -sel_any_bit
17     it is sufficient if the driver of any bit of a cell port is selected.
18     by default all bits must be selected.
19
20 -wires
21     also add $slice and $concat cells to drive otherwise unused wires.
22
23 -no_outputs
24     do not rewire selected module outputs.
25
26 -port <name>
27     only rewire cell ports with the specified name. can be used multiple
28     times. implies -no_output.
29
30 -no_port <name>
31     do not rewire cell ports with the specified name. can be used multiple
32     times. can not be combined with -port <name>.
33
34 By default selected output wires and all cell ports of selected cells driven
35 by selected wires are rewired.
```

### C.185 splitnets – split up multi-bit nets

```
1 splitnets [options] [selection]
2
3 This command splits multi-bit nets into single-bit nets.
4
5 -format char1[char2[char3]]
6     the first char is inserted between the net name and the bit index, the
7     second char is appended to the netname. e.g. -format () creates net
8     names like 'mysignal(42)'. the 3rd character is the range separation
9     character when creating multi-bit wires. the default is '[:']'.
10
11 -ports
12     also split module ports. per default only internal signals are split.
13
```

```

14     -driver
15         don't blindly split nets in individual bits. instead look at the driver
16         and split nets so that no driver drives only part of a net.

```

## C.186 sta – perform static timing analysis

```

1     sta [options] [selection]
2
3     This command performs static timing analysis on the design. (Only considers
4     paths within a single module, so the design must be flattened.)

```

## C.187 stat – print some statistics

```

1     stat [options] [selection]
2
3     Print some statistics (number of objects) on the selected portion of the
4     design.
5
6     -top <module>
7         print design hierarchy with this module as top. if the design is fully
8         selected and a module has the 'top' attribute set, this module is used
9         default value for this option.
10
11     -liberty <liberty_file>
12         use cell area information from the provided liberty file
13
14     -tech <technology>
15         print area estimate for the specified technology. Currently supported
16         values for <technology>: xilinx, cmos
17
18     -width
19         annotate internal cell types with their word width.
20         e.g. $add_8 for an 8 bit wide $add cell.

```

## C.188 submod – moving part of a module to a new submodule

```

1     submod [options] [selection]
2
3     This pass identifies all cells with the 'submod' attribute and moves them to
4     a newly created module. The value of the attribute is used as name for the
5     cell that replaces the group of cells with the same attribute value.
6
7     This pass can be used to create a design hierarchy in flat design. This can
8     be useful for analyzing or reverse-engineering a design.
9
10    This pass only operates on completely selected modules with no processes
11    or memories.

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
12
13  -copy
14      by default the cells are 'moved' from the source module and the source
15      module will use an instance of the new module after this command is
16      finished. call with -copy to not modify the source module.
17
18  -name <name>
19      don't use the 'submod' attribute but instead use the selection. only
20      objects from one module might be selected. the value of the -name option
21      is used as the value of the 'submod' attribute instead.
22
23  -hidden
24      instead of creating submodule ports with public names, create ports with
25      private names so that a subsequent 'flatten; clean' call will restore the
26      original module with original public names.
```

### C.189 supercover – add hi/lo cover cells for each wire bit

```
1      supercover [options] [selection]
2
3  This command adds two cover cells for each bit of each selected wire, one
4  checking for a hi signal level and one checking for lo level.
```

### C.190 synth – generic synthesis script

```
1      synth [options]
2
3  This command runs the default synthesis script. This command does not operate
4  on partly selected designs.
5
6  -top <module>
7      use the specified module as top module (default='top')
8
9  -auto-top
10     automatically determine the top of the design hierarchy
11
12  -flatten
13     flatten the design before synthesis. this will pass '-auto-top' to
14     'hierarchy' if no top module is specified.
15
16  -encfile <file>
17     passed to 'fsm_recode' via 'fsm'
18
19  -lut <k>
20     perform synthesis for a k-LUT architecture.
21
22  -nofsm
23     do not run FSM optimization
24
25  -noabc
```



## APPENDIX C. COMMAND REFERENCE MANUAL

```
26     do not run abc (as if yosys was compiled without ABC support)
27
28     -noalumacc
29         do not run 'alumacc' pass. i.e. keep arithmetic operators in
30         their direct form ($add, $sub, etc.).
31
32     -nordff
33         passed to 'memory'. prohibits merging of FFs into memory read ports
34
35     -noshare
36         do not run SAT-based resource sharing
37
38     -run <from_label>[:<to_label>]
39         only run the commands between the labels (see below). an empty
40         from label is synonymous to 'begin', and empty to label is
41         synonymous to the end of the command list.
42
43     -abc9
44         use new ABC9 flow (EXPERIMENTAL)
45
46     -flowmap
47         use FlowMap LUT techmapping instead of ABC
48
49     -no-rw-check
50         marks all recognized read ports as "return don't-care value on
51         read/write collision" (same result as setting the no_rw_check
52         attribute on all memories).
53
54
55 The following commands are executed by this synthesis command:
56
57     begin:
58         hierarchy -check [-top <top> | -auto-top]
59
60     coarse:
61         proc
62         flatten      (if -flatten)
63         opt_expr
64         opt_clean
65         check
66         opt -nodffe -nosdff
67         fsm          (unless -nofsm)
68         opt
69         wreduce
70         peepopt
71         opt_clean
72         techmap -map +/cmp2lut.v -map +/cmp2lcu.v      (if -lut)
73         alumacc      (unless -noalumacc)
74         share        (unless -noshare)
75         opt
76         memory -nomap
77         opt_clean
78
79     fine:
```

```

80     opt -fast -full
81     memory_map
82     opt -full
83     techmap
84     techmap -map +/gate2lut.v      (if -noabc and -lut)
85     clean; opt_lut                (if -noabc and -lut)
86     flowmap -maxlut K              (if -flowmap and -lut)
87     opt -fast
88     abc -fast                      (unless -noabc, unless -lut)
89     abc -fast -lut k               (unless -noabc, if -lut)
90     opt -fast                      (unless -noabc)
91
92     check:
93         hierarchy -check
94         stat
95         check

```

## C.191 synth\_achronix – synthesis for Achronix Speedster22i FPGAs.

```

1     synth_achronix [options]
2
3     This command runs synthesis for Achronix Speedster eFPGAs. This work is still experimental.
4
5     -top <module>
6         use the specified module as top module (default='top')
7
8     -vout <file>
9         write the design to the specified Verilog netlist file. writing of an
10        output file is omitted if this parameter is not specified.
11
12     -run <from_label>:<to_label>
13         only run the commands between the labels (see below). an empty
14         from label is synonymous to 'begin', and empty to label is
15         synonymous to the end of the command list.
16
17     -noflatten
18         do not flatten design before synthesis
19
20     -retime
21         run 'abc' with '-dff -D 1' options
22
23
24     The following commands are executed by this synthesis command:
25
26     begin:
27         read_verilog -sv -lib +/achronix/speedster22i/cells_sim.v
28         hierarchy -check -top <top>
29
30     flatten:    (unless -noflatten)
31         proc
32         flatten

```

```

33     tribuf -logic
34     deminout
35
36 coarse:
37     synth -run coarse
38
39 fine:
40     opt -fast -mux_undef -undriven -fine -full
41     memory_map
42     opt -undriven -fine
43     opt -fine
44     techmap -map +/techmap.v
45     opt -full
46     clean -purge
47     setundef -undriven -zero
48     dfflegalize -cell $_DFF_P_ x
49     abc -markgroups -dff -D 1    (only if -retime)
50
51 map_luts:
52     abc -lut 4
53     clean
54
55 map_cells:
56     iopadmap -bits -outpad $__outpad I:0 -inpad $__inpad 0:I
57     techmap -map +/achronix/speedster22i/cells_map.v
58     clean -purge
59
60 check:
61     hierarchy -check
62     stat
63     check -noinit
64     blackbox =A:whitebox
65
66 vout:
67     write_verilog -nodec -attr2comment -defparam -renameprefix syn_ <file-name>

```

## C.192 synth\_anlogic – synthesis for Anlogic FPGAs

```

1     synth_anlogic [options]
2
3     This command runs synthesis for Anlogic FPGAs.
4
5     -top <module>
6         use the specified module as top module
7
8     -edif <file>
9         write the design to the specified EDIF file. writing of an output file
10        is omitted if this parameter is not specified.
11
12    -json <file>
13        write the design to the specified JSON file. writing of an output file
14        is omitted if this parameter is not specified.

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

15
16 -run <from_label>:<to_label>
17     only run the commands between the labels (see below). an empty
18     from label is synonymous to 'begin', and empty to label is
19     synonymous to the end of the command list.
20
21 -noflatten
22     do not flatten design before synthesis
23
24 -retime
25     run 'abc' with '-dff -D 1' options
26
27 -nolutram
28     do not use EG_LOGIC_DRAM16X4 cells in output netlist
29
30 -nobram
31     do not use EG_PHY_BRAM or EG_PHY_BRAM32K cells in output netlist
32
33
34 The following commands are executed by this synthesis command:
35
36 begin:
37     read_verilog -lib +/anlogic/cells_sim.v +/anlogic/eagle_bb.v
38     hierarchy -check -top <top>
39
40 flatten:    (unless -noflatten)
41     proc
42     flatten
43     tribuf -logic
44     deminout
45
46 coarse:
47     synth -run coarse
48
49 map_ram:
50     memory_libmap -lib +/anlogic/lutrams.txt -lib +/anlogic/brams.txt [-no-auto-block] [-no-auto-distrib
(-no-auto-block if -nobram, -no-auto-distributed if -nolutram)
51     techmap -map +/anlogic/lutrams_map.v -map +/anlogic/brams_map.v
52
53 map_ffram:
54     opt -fast -mux_undef -undriven -fine
55     memory_map
56     opt -undriven -fine
57
58 map_gates:
59     techmap -map +/techmap.v -map +/anlogic/arith_map.v
60     opt -fast
61     abc -dff -D 1    (only if -retime)
62
63 map_ffs:
64     dfflegalize -cell $_DFFE_P??P_ r -cell $_SDFFE_P??P_ r -cell $_DLATCH_N??_ r
65     techmap -D NO_LUT -map +/anlogic/cells_map.v
66     opt_expr -mux_undef
67     simplemap

```

```

68
69 map_luts:
70     abc -lut 4:6
71     clean
72
73 map_cells:
74     techmap -map +/anlogic/cells_map.v
75     clean
76
77 map_anlogic:
78     anlogic_fixcarry
79     anlogic_eqn
80
81 check:
82     hierarchy -check
83     stat
84     check -noinit
85     blackbox =A:whitebox
86
87 edif:
88     write_edif <file-name>
89
90 json:
91     write_json <file-name>

```

### C.193 synth\_coolrunner2 – synthesis for Xilinx Coolrunner-II CPLDs

```

1 synth_coolrunner2 [options]
2
3 This command runs synthesis for Coolrunner-II CPLDs. This work is experimental.
4 It is intended to be used with https://github.com/azonenberg/openfpga as the
5 place-and-route.
6
7 -top <module>
8     use the specified module as top module (default='top')
9
10 -json <file>
11     write the design to the specified JSON file. writing of an output file
12     is omitted if this parameter is not specified.
13
14 -run <from_label>:<to_label>
15     only run the commands between the labels (see below). an empty
16     from label is synonymous to 'begin', and empty to label is
17     synonymous to the end of the command list.
18
19 -noflatten
20     do not flatten design before synthesis
21
22 -retime
23     run 'abc' with '-dff -D 1' options
24

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

25
26 The following commands are executed by this synthesis command:
27
28 begin:
29     read_verilog -lib +/coolrunner2/cells_sim.v
30     hierarchy -check -top <top>
31
32 flatten:      (unless -noflatten)
33     proc
34     flatten
35     tribuf -logic
36
37 coarse:
38     synth -run coarse
39
40 fine:
41     extract_counter -dir up -allow_arst no
42     techmap -map +/coolrunner2/cells_counter_map.v
43     clean
44     opt -fast -full
45     techmap -map +/techmap.v -map +/coolrunner2/cells_latch.v
46     opt -fast
47     dfflibmap -prepare -liberty +/coolrunner2/xc2_dff.lib
48
49 map_tff:
50     abc -g AND,XOR
51     clean
52     extract -map +/coolrunner2/tff_extract.v
53
54 map_pla:
55     abc -sop -I 40 -P 56
56     clean
57
58 map_cells:
59     dfflibmap -liberty +/coolrunner2/xc2_dff.lib
60     dffinit -ff FDCP Q INIT
61     dffinit -ff FDCP_N Q INIT
62     dffinit -ff FTCP Q INIT
63     dffinit -ff FTCP_N Q INIT
64     dffinit -ff LDCP Q INIT
65     dffinit -ff LDCP_N Q INIT
66     coolrunner2_sop
67     clean
68     iopadmap -bits -inpad IBUF 0:I -outpad IOBUFE I:IO -inoutpad IOBUFE 0:IO -toutpad IOBUFE E:I:IO -ti
69     attrmvp -attr src -attr LOC t:IOBUFE n:*
70     attrmvp -attr src -attr LOC -driven t:IBUF n:*
71     coolrunner2_fixup
72     splitnets
73     clean
74
75 check:
76     hierarchy -check
77     stat
78     check -noinit

```

```

79     blackbox =A:whitebox
80
81     json:
82         write_json <file-name>

```

## C.194 synth\_easic – synthesis for eASIC platform

```

1     synth_easic [options]
2
3     This command runs synthesis for eASIC platform.
4
5     -top <module>
6         use the specified module as top module
7
8     -vlog <file>
9         write the design to the specified structural Verilog file. writing of
10        an output file is omitted if this parameter is not specified.
11
12     -etools <path>
13         set path to the eTools installation. (default=/opt/eTools)
14
15     -run <from_label>:<to_label>
16         only run the commands between the labels (see below). an empty
17         from label is synonymous to 'begin', and empty to label is
18         synonymous to the end of the command list.
19
20     -noflatten
21         do not flatten design before synthesis
22
23     -retime
24         run 'abc' with '-dff -D 1' options
25
26
27     The following commands are executed by this synthesis command:
28
29     begin:
30         read_liberty -lib <etools_phys_clk_lib>
31         read_liberty -lib <etools_logic_lut_lib>
32         hierarchy -check -top <top>
33
34     flatten:    (unless -noflatten)
35         proc
36         flatten
37
38     coarse:
39         synth -run coarse
40
41     fine:
42         opt -fast -mux_undef -undriven -fine
43         memory_map
44         opt -undriven -fine
45         techmap

```

```

46     opt -fast
47     abc -dff -D 1      (only if -retime)
48     opt_clean      (only if -retime)
49
50 map:
51     dfflibmap -liberty <etools_phys_clk_lib>
52     abc -liberty <etools_logic_lut_lib>
53     opt_clean
54
55 check:
56     hierarchy -check
57     stat
58     check -noinit
59     blackbox =A:whitebox
60
61 vlog:
62     write_verilog -noexpr -attr2comment <file-name>

```

## C.195 synth\_ecp5 – synthesis for ECP5 FPGAs

```

1     synth_ecp5 [options]
2
3 This command runs synthesis for ECP5 FPGAs.
4
5 -top <module>
6     use the specified module as top module
7
8 -blif <file>
9     write the design to the specified BLIF file. writing of an output file
10    is omitted if this parameter is not specified.
11
12 -edif <file>
13    write the design to the specified EDIF file. writing of an output file
14    is omitted if this parameter is not specified.
15
16 -json <file>
17    write the design to the specified JSON file. writing of an output file
18    is omitted if this parameter is not specified.
19
20 -run <from_label>:<to_label>
21    only run the commands between the labels (see below). an empty
22    from label is synonymous to 'begin', and empty to label is
23    synonymous to the end of the command list.
24
25 -noflatten
26    do not flatten design before synthesis
27
28 -dff
29    run 'abc'/'abc9' with -dff option
30
31 -retime
32    run 'abc' with '-dff -D 1' options

```



## APPENDIX C. COMMAND REFERENCE MANUAL

```
33
34 -noccu2
35     do not use CCU2 cells in output netlist
36
37 -nodffe
38     do not use flipflops with CE in output netlist
39
40 -nobram
41     do not use block RAM cells in output netlist
42
43 -nolutram
44     do not use LUT RAM cells in output netlist
45
46 -nowidelut
47     do not use PFU muxes to implement LUTs larger than LUT4s
48
49 -asyncprld
50     use async PRLD mode to implement ALDFF (EXPERIMENTAL)
51
52 -abc2
53     run two passes of 'abc' for slightly improved logic density
54
55 -abc9
56     use new ABC9 flow (EXPERIMENTAL)
57
58 -vpr
59     generate an output netlist (and BLIF file) suitable for VPR
60     (this feature is experimental and incomplete)
61
62 -nodsp
63     do not map multipliers to MULT18X18D
64
65 -no-rw-check
66     marks all recognized read ports as "return don't-care value on
67     read/write collision" (same result as setting the no_rw_check
68     attribute on all memories).
69
70
71 The following commands are executed by this synthesis command:
72
73 begin:
74     read_verilog -lib -specify +/ecp5/cells_sim.v +/ecp5/cells_bb.v
75     hierarchy -check -top <top>
76
77 coarse:
78     proc
79     flatten
80     tribuf -logic
81     deminout
82     opt_expr
83     opt_clean
84     check
85     opt -nodffe -nosdff
86     fsm
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

87     opt
88     wreduce
89     peepopt
90     opt_clean
91     share
92     techmap -map +/cmp2lut.v -D LUT_WIDTH=4
93     opt_expr
94     opt_clean
95     techmap -map +/mul2dsp.v -map +/ecp5/dsp_map.v -D DSP_A_MAXWIDTH=18 -D DSP_B_MAXWIDTH=18
-D DSP_A_MINWIDTH=2 -D DSP_B_MINWIDTH=2 -D DSP_NAME=__MUL18X18 (unless -nodsp)
96     chtype -set $mul t:$__soft_mul (unless -nodsp)
97     alumacc
98     opt
99     memory -nomap [-no-rw-check]
100    opt_clean
101
102    map_ram:
103        memory_libmap -lib +/ecp5/lutrams.txt -lib +/ecp5/brams.txt [-no-auto-block] [-no-auto-distributed]
(-no-auto-block if -nobram, -no-auto-distributed if -nolutram)
104        techmap -map +/ecp5/lutrams_map.v -map +/ecp5/brams_map.v
105
106    map_ffram:
107        opt -fast -mux_undef -undriven -fine
108        memory_map
109        opt -undriven -fine
110
111    map_gates:
112        techmap -map +/techmap.v -map +/ecp5/arith_map.v
113        opt -fast
114        abc -dff -D 1 (only if -retime)
115
116    map_ffs:
117        opt_clean
118        dfflegalize -cell $_DFF_?_ 01 -cell $_DFF_?P?_ r -cell $_SDFF_?P?_ r [-cell $_DFFE_??_ 01 -cell $_D
($_ALDFF_*_ only if -asyncprld, $_DLATCH_*_ only if not -asyncprld, $_*DFFE_*_ only if not -nodffe)
119        zinit -all w:* t:$_DFF_?_ t:$_DFFE_??_ t:$_SDFF* (only if -abc9 and -dff)
120        techmap -D NO_LUT -map +/ecp5/cells_map.v
121        opt_expr -undriven -mux_undef
122        simplemap
123        ecp5_gsr
124        attrmvp -copy -attr syn_useioff
125        opt_clean
126
127    map_luts:
128        abc (only if -abc2)
129        techmap -map +/ecp5/latches_map.v (skip if -asyncprld)
130        abc -dress -lut 4:7
131        clean
132
133    map_cells:
134        techmap -map +/ecp5/cells_map.v (skip if -vpr)
135        opt_lut_ins -tech ecp5
136        clean
137

```

```

138  check:
139      autoname
140      hierarchy -check
141      stat
142      check -noinit
143      blackbox =A:whitebox
144
145  blif:
146      opt_clean -purge                                (vpr mode)
147      write_blif -attr -cname -conn -param <file-name> (vpr mode)
148      write_blif -gates -attr -param <file-name>      (non-vpr mode)
149
150  edif:
151      write_edif <file-name>
152
153  json:
154      write_json <file-name>

```

## C.196 synth\_efinix – synthesis for Efinix FPGAs

```

1  synth_efinix [options]
2
3  This command runs synthesis for Efinix FPGAs.
4
5  -top <module>
6      use the specified module as top module
7
8  -edif <file>
9      write the design to the specified EDIF file. writing of an output file
10     is omitted if this parameter is not specified.
11
12  -json <file>
13     write the design to the specified JSON file. writing of an output file
14     is omitted if this parameter is not specified.
15
16  -run <from_label>:<to_label>
17     only run the commands between the labels (see below). an empty
18     from label is synonymous to 'begin', and empty to label is
19     synonymous to the end of the command list.
20
21  -noflatten
22     do not flatten design before synthesis
23
24  -retime
25     run 'abc' with '-dff -D 1' options
26
27  -nobram
28     do not use EFX_RAM_5K cells in output netlist
29
30
31  The following commands are executed by this synthesis command:
32

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

33 begin:
34     read_verilog -lib +/efinix/cells_sim.v
35     hierarchy -check -top <top>
36
37 flatten:    (unless -noflatten)
38     proc
39     flatten
40     tribuf -logic
41     deminout
42
43 coarse:
44     synth -run coarse
45
46 map_ram:
47     memory_libmap -lib +/efinix/brams.txt
48     techmap -map +/efinix/brams_map.v
49
50 map_ffram:
51     opt -fast -mux_undef -undriven -fine
52     memory_map
53     opt -undriven -fine
54
55 map_gates:
56     techmap -map +/techmap.v -map +/efinix/arith_map.v
57     opt -fast
58     abc -dff -D 1      (only if -retime)
59
60 map_ffs:
61     dfflegalize -cell $_DFFE_????_ 0 -cell $_SDFFE_????_ 0 -cell $_SDFFCFCE_????_ 0 -cell $_DLATCH_?_ x
62     techmap -D NO_LUT -map +/efinix/cells_map.v
63     opt_expr -mux_undef
64     simplemap
65
66 map_luts:
67     abc -lut 4
68     clean
69
70 map_cells:
71     techmap -map +/efinix/cells_map.v
72     clean
73
74 map_gbuf:
75     clkbufmap -buf $_EFX_GBUF 0:I
76     techmap -map +/efinix/gbuf_map.v
77     efinix_fixcarry
78     clean
79
80 check:
81     hierarchy -check
82     stat
83     check -noinit
84     blackbox =A:whitebox
85
86 edif:

```

```

87     write_edif <file-name>
88
89     json:
90         write_json <file-name>

```

## C.197 synth\_gatemate – synthesis for Cologne Chip GateMate FPGAs

```

1     synth_gatemate [options]
2
3     This command runs synthesis for Cologne Chip AG GateMate FPGAs.
4
5     -top <module>
6         use the specified module as top module.
7
8     -vlog <file>
9         write the design to the specified verilog file. Writing of an output
10        file is omitted if this parameter is not specified.
11
12    -json <file>
13        write the design to the specified JSON file. Writing of an output file
14        is omitted if this parameter is not specified.
15
16    -run <from_label>:<to_label>
17        only run the commands between the labels (see below). An empty
18        from label is synonymous to 'begin', and empty to label is
19        synonymous to the end of the command list.
20
21    -noflatten
22        do not flatten design before synthesis.
23
24    -nobram
25        do not use CC_BRAM_20K or CC_BRAM_40K cells in output netlist.
26
27    -noaddf
28        do not use CC_ADDF full adder cells in output netlist.
29
30    -nomult
31        do not use CC_MULT multiplier cells in output netlist.
32
33    -nomx8, -nomx4
34        do not use CC_MX{8,4} multiplexer cells in output netlist.
35
36    -luttree
37        use new LUT tree mapping approach (EXPERIMENTAL).
38
39    -dff
40        run 'abc' with -dff option
41
42    -retime
43        run 'abc' with '-dff -D 1' options
44

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
45     -noiopad
46         disable I/O buffer insertion (useful for hierarchical or
47         out-of-context flows).
48
49     -noclkbuf
50         disable automatic clock buffer insertion.
51
52 The following commands are executed by this synthesis command:
53
54     begin:
55         read_verilog -lib -specify +/gatemate/cells_sim.v +/gatemate/cells_bb.v
56         hierarchy -check -top <top>
57
58     prepare:
59         proc
60         flatten
61         tribuf -logic
62         deminout
63         opt_expr
64         opt_clean
65         check
66         opt -nodffe -nosdff
67         fsm
68         opt
69         wreduce
70         peepopt
71         opt_clean
72         muxpack
73         share
74         techmap -map +/cmp2lut.v -D LUT_WIDTH=4
75         opt_expr
76         opt_clean
77
78     map_mult:      (skip if '-nomult')
79         techmap -map +/gatemate/mul_map.v
80
81     coarse:
82         alumacc
83         opt
84         memory -nomap
85         opt_clean
86
87     map_bram:      (skip if '-nobram')
88         memory_libmap -lib +/gatemate/brams.txt
89         techmap -map +/gatemate/brams_map.v
90
91     map_ffram:
92         opt -fast -mux_undef -undriven -fine
93         memory_map
94         opt -undriven -fine
95
96     map_gates:
97         techmap -map +/techmap.v -map +/gatemate/arith_map.v
98         opt -fast
```

```

99
100 map_io:    (skip if '-noiopad')
101     iopadmap -bits -inpad CC_IBUF Y:I -outpad CC_OBUF A:0 -toutpad CC_TOBUF ~T:A:0 -tinoutpad CC_IOBUF
102     clean
103
104 map_regs:
105     opt_clean
106     dfflegalize -cell $_DFFE_????_ x -cell $_DLATCH_???_ x
107     techmap -map +/gatemate/reg_map.v
108     opt_expr -mux_undef
109     simplemap
110     opt_clean
111
112 map_muxs:
113     muxcover -mux4 -mux8
114     opt -full
115     techmap -map +/gatemate/mux_map.v
116
117 map_luts:
118     abc -genlib +/gatemate/lut_tree_cells.genlib (with -luttree)
119     techmap -map +/gatemate/lut_tree_map.v (with -luttree)
120     gatemate_foldinv (with -luttree)
121     techmap -map +/gatemate/inv_map.v (with -luttree)
122     abc -dress -lut 4 (without -luttree)
123     clean
124
125 map_cells:
126     techmap -map +/gatemate/lut_map.v
127     clean
128
129 map_bufg:    (skip if '-noclkbuf')
130     clkbufmap -buf CC_BUFG 0:I
131     clean
132
133 check:
134     hierarchy -check
135     stat -width
136     check -noinit
137     blackbox =A:whitebox
138
139 vlog:
140     opt_clean -purge
141     write_verilog -noattr <file-name>
142
143 json:
144     write_json <file-name>

```

## C.198 synth\_gowin – synthesis for Gowin FPGAs

```

1 synth_gowin [options]
2
3 This command runs synthesis for Gowin FPGAs. This work is experimental.

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
4
5  -top <module>
6      use the specified module as top module (default='top')
7
8  -vout <file>
9      write the design to the specified Verilog netlist file. writing of an
10     output file is omitted if this parameter is not specified.
11
12  -json <file>
13     write the design to the specified JSON netlist file. writing of an
14     output file is omitted if this parameter is not specified.
15     This disables features not yet supported by nexpnr-gowin.
16
17  -run <from_label>:<to_label>
18     only run the commands between the labels (see below). an empty
19     from label is synonymous to 'begin', and empty to label is
20     synonymous to the end of the command list.
21
22  -nodffe
23     do not use flipflops with CE in output netlist
24
25  -nobram
26     do not use BRAM cells in output netlist
27
28  -nolutram
29     do not use distributed RAM cells in output netlist
30
31  -noflatten
32     do not flatten design before synthesis
33
34  -retime
35     run 'abc' with '-dff -D 1' options
36
37  -nowidelut
38     do not use muxes to implement LUTs larger than LUT4s
39
40  -noiopads
41     do not emit IOB at top level ports
42
43  -noalu
44     do not use ALU cells
45
46  -abc9
47     use new ABC9 flow (EXPERIMENTAL)
48
49  -no-rw-check
50     marks all recognized read ports as "return don't-care value on
51     read/write collision" (same result as setting the no_rw_check
52     attribute on all memories).
```

The following commands are executed by this synthesis command:

```
begin:
```



## APPENDIX C. COMMAND REFERENCE MANUAL

```

58     read_verilog -specify -lib +/gowin/cells_sim.v
59     hierarchy -check -top <top>
60
61     flatten:      (unless -noflatten)
62         proc
63         flatten
64         tribuf -logic
65         deminout
66
67     coarse:
68         synth -run coarse [-no-rw-check]
69
70     map_ram:
71         memory_libmap -lib +/gowin/lutrams.txt -lib +/gowin/brams.txt [-no-auto-block] [-no-auto-distribute
(-no-auto-block if -nobram, -no-auto-distributed if -nolutram)
72         techmap -map +/gowin/lutrams_map.v -map +/gowin/brams_map.v
73
74     map_ffram:
75         opt -fast -mux_undef -undriven -fine
76         memory_map
77         opt -undriven -fine
78
79     map_gates:
80         techmap -map +/techmap.v -map +/gowin/arith_map.v
81         opt -fast
82         abc -dff -D 1      (only if -retime)
83         iopadmap -bits -inpad IBUF 0:I -outpad OBUF I:0 -toutpad TBUF ~OEN:I:0 -tinoutpad IOBUF ~OEN:0:I:IO
(unless -noiopads)
84
85     map_ffs:
86         opt_clean
87         dfflegalize -cell $_DFF_?_ 0 -cell $_DFFE_?P_ 0 -cell $_SDFF_?P?_ r -cell $_SDFFE_?P?P_ r -cell $_D
88         techmap -map +/gowin/cells_map.v
89         opt_expr -mux_undef
90         simplemap
91
92     map_luts:
93         abc -lut 4:8
94         clean
95
96     map_cells:
97         techmap -map +/gowin/cells_map.v
98         opt_lut_ins -tech gowin
99         setundef -undriven -params -zero
100        hilomap -singleton -hicell VCC V -locell GND G
101        splitnets -ports      (only if -vout used)
102        clean
103        autoname
104
105     check:
106         hierarchy -check
107         stat
108         check -noinit
109         blackbox =A:whitebox

```

```

110
111     vout:
112         write_verilog -simple-lhs -decimal -attr2comment -defparam -renameprefix gen <file-name>
113         write_json <file-name>

```

## C.199 synth\_greenpak4 – synthesis for GreenPAK4 FPGAs

```

1     synth_greenpak4 [options]
2
3     This command runs synthesis for GreenPAK4 FPGAs. This work is experimental.
4     It is intended to be used with https://github.com/azonenberg/openfpga as the
5     place-and-route.
6
7     -top <module>
8         use the specified module as top module (default='top')
9
10    -part <part>
11        synthesize for the specified part. Valid values are SLG46140V,
12        SLG46620V, and SLG46621V (default).
13
14    -json <file>
15        write the design to the specified JSON file. writing of an output file
16        is omitted if this parameter is not specified.
17
18    -run <from_label>:<to_label>
19        only run the commands between the labels (see below). an empty
20        from label is synonymous to 'begin', and empty to label is
21        synonymous to the end of the command list.
22
23    -noflatten
24        do not flatten design before synthesis
25
26    -retime
27        run 'abc' with '-dff -D 1' options
28
29
30    The following commands are executed by this synthesis command:
31
32    begin:
33        read_verilog -lib +/greenpak4/cells_sim.v
34        hierarchy -check -top <top>
35
36    flatten:    (unless -noflatten)
37        proc
38        flatten
39        tribuf -logic
40
41    coarse:
42        synth -run coarse
43
44    fine:
45        extract_counter -pout GP_DCMP,GP_DAC -maxwidth 14

```

```

46      clean
47      opt -fast -mux_undef -undriven -fine
48      memory_map
49      opt -undriven -fine
50      techmap -map +/techmap.v -map +/greenpak4/cells_latch.v
51      dfflibmap -prepare -liberty +/greenpak4/gp_dff.lib
52      opt -fast -noclkinv -noff
53      abc -dff -D 1      (only if -retime)
54
55  map_luts:
56      nlutmap -assert -luts 0,6,8,2      (for -part SLG46140V)
57      nlutmap -assert -luts 2,8,16,2      (for -part SLG46620V)
58      nlutmap -assert -luts 2,8,16,2      (for -part SLG46621V)
59      clean
60
61  map_cells:
62      shregmap -tech greenpak4
63      dfflibmap -liberty +/greenpak4/gp_dff.lib
64      dffinit -ff GP_DFF Q INIT
65      dffinit -ff GP_DFFR Q INIT
66      dffinit -ff GP_DFFS Q INIT
67      dffinit -ff GP_DFFSR Q INIT
68      iopadmap -bits -inpad GP_IBUF OUT:IN -outpad GP_OBUF IN:OUT -inoutpad GP_OBUF OUT:IN -toutpad GP_OB
69      attrmvp -attr src -attr LOC t:GP_OBUF t:GP_OBUFT t:GP_IOBUF n:*
70      attrmvp -attr src -attr LOC -driven t:GP_IBUF n:*
71      techmap -map +/greenpak4/cells_map.v
72      greenpak4_dffinv
73      clean
74
75  check:
76      hierarchy -check
77      stat
78      check -noinit
79      blackbox =A:whitebox
80
81  json:
82      write_json <file-name>

```

## C.200 synth\_ice40 – synthesis for iCE40 FPGAs

```

1      synth_ice40 [options]
2
3  This command runs synthesis for iCE40 FPGAs.
4
5      -device < hx | lp | u >
6          relevant only for '-abc9' flow, optimise timing for the specified device.
7          default: hx
8
9      -top <module>
10         use the specified module as top module
11
12      -blif <file>

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
13      write the design to the specified BLIF file. writing of an output file
14      is omitted if this parameter is not specified.
15
16      -edif <file>
17          write the design to the specified EDIF file. writing of an output file
18          is omitted if this parameter is not specified.
19
20      -json <file>
21          write the design to the specified JSON file. writing of an output file
22          is omitted if this parameter is not specified.
23
24      -run <from_label>:<to_label>
25          only run the commands between the labels (see below). an empty
26          from label is synonymous to 'begin', and empty to label is
27          synonymous to the end of the command list.
28
29      -noflatten
30          do not flatten design before synthesis
31
32      -dff
33          run 'abc'/'abc9' with -dff option
34
35      -retime
36          run 'abc' with '-dff -D 1' options
37
38      -nocarry
39          do not use SB_CARRY cells in output netlist
40
41      -nodffe
42          do not use SB_DFFE* cells in output netlist
43
44      -dff_min_ce_use <min_ce_use>
45          do not use SB_DFFE* cells if the resulting CE line would go to less
46          than min_ce_use SB_DFFE* in output netlist
47
48      -nobram
49          do not use SB_RAM40_4K* cells in output netlist
50
51      -spram
52          enable automatic inference of SB_SDRAM256KA
53
54      -dsp
55          use iCE40 UltraPlus DSP cells for large arithmetic
56
57      -noabc
58          use built-in Yosys LUT techmapping instead of abc
59
60      -abc2
61          run two passes of 'abc' for slightly improved logic density
62
63      -vpr
64          generate an output netlist (and BLIF file) suitable for VPR
65          (this feature is experimental and incomplete)
66
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

67     -abc9
68         use new ABC9 flow (EXPERIMENTAL)
69
70     -flowmap
71         use FlowMap LUT techmapping instead of abc (EXPERIMENTAL)
72
73     -no-rw-check
74         marks all recognized read ports as "return don't-care value on
75         read/write collision" (same result as setting the no_rw_check
76         attribute on all memories).
77
78
79 The following commands are executed by this synthesis command:
80
81     begin:
82         read_verilog -D ICE40_HX -lib -specify +/ice40/cells_sim.v
83         hierarchy -check -top <top>
84         proc
85
86     flatten:    (unless -noflatten)
87         flatten
88         tribuf -logic
89         deminout
90
91     coarse:
92         opt_expr
93         opt_clean
94         check
95         opt -nodffe -nosdff
96         fsm
97         opt
98         wreduce
99         peepopt
100        opt_clean
101        share
102        techmap -map +/cmp2lut.v -D LUT_WIDTH=4
103        opt_expr
104        opt_clean
105        memory_dff [-no-rw-check]
106        wreduce t:$mul
107        techmap -map +/mul2dsp.v -map +/ice40/dsp_map.v -D DSP_A_MAXWIDTH=16 -D DSP_B_MAXWIDTH=16 -D DSP_A_
(if -dsp)
108        select a:mul2dsp                                (if -dsp)
109        setattr -unset mul2dsp                          (if -dsp)
110        opt_expr -fine                                    (if -dsp)
111        wreduce                                           (if -dsp)
112        select -clear                                     (if -dsp)
113        ice40_dsp                                         (if -dsp)
114        chtype -set $mul t:$__soft_mul                  (if -dsp)
115        alumacc
116        opt
117        memory -nomap [-no-rw-check]
118        opt_clean
119

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

120     map_ram:
121         memory_libmap -lib +/ice40/brams.txt -lib +/ice40/spram.txt -no-auto-huge [-no-auto-huge] [-no-auto
(-no-auto-huge unless -spram, -no-auto-block if -nobram)
122         techmap -map +/ice40/brams_map.v -map +/ice40/spram_map.v
123         ice40_braminit
124
125     map_ffram:
126         opt -fast -mux_undef -undriven -fine
127         memory_map
128         opt -undriven -fine
129
130     map_gates:
131         ice40_wrapcarry
132         techmap -map +/techmap.v -map +/ice40/arith_map.v
133         opt -fast
134         abc -dff -D 1      (only if -retime)
135         ice40_opt
136
137     map_ffs:
138         dfflegalize -cell $_DFF?_ 0 -cell $_DFFE?P_ 0 -cell $_DFF?P?_ 0 -cell $_DFFE?P?P_ 0 -cell $_SDF
139         techmap -map +/ice40/ff_map.v
140         opt_expr -mux_undef
141         simplemap
142         ice40_opt -full
143
144     map_luts:
145         abc      (only if -abc2)
146         ice40_opt (only if -abc2)
147         techmap -map +/ice40/latches_map.v
148         simplemap      (if -noabc or -flowmap)
149         techmap -map +/gate2lut.v -D LUT_WIDTH=4      (only if -noabc)
150         flowmap -maxlut 4      (only if -flowmap)
151         abc -dress -lut 4      (skip if -noabc)
152         ice40_wrapcarry -unwrap
153         techmap -map +/ice40/ff_map.v
154         clean
155         opt_lut -dlogic SB_CARRY:I0=1:I1=2:CI=3 -dlogic SB_CARRY:CO=3
156
157     map_cells:
158         techmap -map +/ice40/cells_map.v      (skip if -vpr)
159         clean
160
161     check:
162         autaname
163         hierarchy -check
164         stat
165         check -noinit
166         blackbox =A:whitebox
167
168     blif:
169         opt_clean -purge      (vpr mode)
170         write_blif -attr -cname -conn -param <file-name>      (vpr mode)
171         write_blif -gates -attr -param <file-name>      (non-vpr mode)
172

```

```

173     edif:
174         write_edif <file-name>
175
176     json:
177         write_json <file-name>

```

## C.201 synth\_intel – synthesis for Intel (Altera) FPGAs.

```

1     synth_intel [options]
2
3     This command runs synthesis for Intel FPGAs.
4
5     -family <max10 | cyclone10lp | cycloneiv | cycloneive>
6         generate the synthesis netlist for the specified family.
7         MAX10 is the default target if no family argument specified.
8         For Cyclone IV GX devices, use cycloneiv argument; for Cyclone IV E, use cycloneive.
9         For Cyclone V and Cyclone 10 GX, use the synth_intel_alm backend instead.
10
11     -top <module>
12         use the specified module as top module (default='top')
13
14     -vqm <file>
15         write the design to the specified Verilog Quartus Mapping File. Writing of an
16         output file is omitted if this parameter is not specified.
17         Note that this backend has not been tested and is likely incompatible
18         with recent versions of Quartus.
19
20     -vpr <file>
21         write BLIF files for VPR flow experiments. The synthesized BLIF output file is not
22         compatible with the Quartus flow. Writing of an
23         output file is omitted if this parameter is not specified.
24
25     -run <from_label>:<to_label>
26         only run the commands between the labels (see below). an empty
27         from label is synonymous to 'begin', and empty to label is
28         synonymous to the end of the command list.
29
30     -iopads
31         use IO pad cells in output netlist
32
33     -nobram
34         do not use block RAM cells in output netlist
35
36     -noflatten
37         do not flatten design before synthesis
38
39     -retime
40         run 'abc' with '-dff -D 1' options
41
42     The following commands are executed by this synthesis command:
43
44     begin:

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

45
46 family:
47     read_verilog -sv -lib +/intel/max10/cells_sim.v
48     read_verilog -sv -lib +/intel/common/m9k_bb.v
49     read_verilog -sv -lib +/intel/common/altp11_bb.v
50     hierarchy -check -top <top>
51
52 flatten:      (unless -noflatten)
53     proc
54     flatten
55     tribuf -logic
56     deminout
57
58 coarse:
59     synth -run coarse
60
61 map_bram:      (skip if -nobram)
62     memory_bram -rules +/intel/common/brams_m9k.txt      (if applicable for family)
63     techmap -map +/intel/common/brams_map_m9k.v          (if applicable for family)
64
65 map_ffram:
66     opt -fast -mux_undef -undriven -fine -full
67     memory_map
68     opt -undriven -fine
69     techmap -map +/techmap.v
70     opt -full
71     clean -purge
72     setundef -undriven -zero
73     abc -markgroups -dff -D 1      (only if -retime)
74
75 map_ffs:
76     dfflegalize -cell $_DFFE_PN0P_ 01
77     techmap -map +/intel/common/ff_map.v
78
79 map_luts:
80     abc -lut 4
81     clean
82
83 map_cells:
84     iopadmap -bits -outpad $__outpad I:0 -inpad $__inpad 0:I      (if -iopads)
85     techmap -map +/intel/max10/cells_map.v
86     clean -purge
87
88 check:
89     hierarchy -check
90     stat
91     check -noinit
92     blackbox =A:whitebox
93
94 vqm:
95     write_verilog -attr2comment -defparam -nohex -decimal -renameprefix syn_ <file-name>
96
97 vpr:
98     opt_clean -purge

```



```

99     write_blif <file-name>
100
101
102 WARNING: THE 'synth_intel' COMMAND IS EXPERIMENTAL.

```

## C.202 synth\_intel\_alm – synthesis for ALM-based Intel (Altera) FPGAs.

```

1  synth_intel_alm [options]
2
3  This command runs synthesis for ALM-based Intel FPGAs.
4
5  -top <module>
6      use the specified module as top module
7
8  -family <family>
9      target one of:
10     "cyclonev"      - Cyclone V (default)
11     "arriav"        - Arria V (non-GZ)      "cyclone10gx" - Cyclone 10GX
12
13  -vqm <file>
14      write the design to the specified Verilog Quartus Mapping File. Writing of an
15      output file is omitted if this parameter is not specified. Implies -quartus.
16
17  -noflatten
18      do not flatten design before synthesis; useful for per-module area statistics
19
20  -quartus
21      output a netlist using Quartus cells instead of MISTRAL_* cells
22
23  -dff
24      pass DFFs to ABC to perform sequential logic optimisations (EXPERIMENTAL)
25
26  -run <from_label>:<to_label>
27      only run the commands between the labels (see below). an empty
28      from label is synonymous to 'begin', and empty to label is
29      synonymous to the end of the command list.
30
31  -nolutram
32      do not use LUT RAM cells in output netlist
33
34  -nobram
35      do not use block RAM cells in output netlist
36
37  -nodsp
38      do not map multipliers to MISTRAL_MUL cells
39
40  -noiopad
41      do not instantiate IO buffers
42
43  -noclkbuf
44      do not insert global clock buffers

```

## APPENDIX C. COMMAND REFERENCE MANUAL

The following commands are executed by this synthesis command:

```

begin:
  read_verilog -specify -lib -D <family> +/intel_alm/common/alm_sim.v
  read_verilog -specify -lib -D <family> +/intel_alm/common/dff_sim.v
  read_verilog -specify -lib -D <family> +/intel_alm/common/dsp_sim.v
  read_verilog -specify -lib -D <family> +/intel_alm/common/mem_sim.v
  read_verilog -specify -lib -D <family> +/intel_alm/common/misc_sim.v
  read_verilog -specify -lib -D <family> -icells +/intel_alm/common/abc9_model.v
  read_verilog -lib +/intel/common/altp11_bb.v
  read_verilog -lib +/intel_alm/common/megafunction_bb.v
  hierarchy -check -top <top>

coarse:
  proc
  flatten      (skip if -noflatten)
  tribuf -logic
  deminout
  opt_expr
  opt_clean
  check
  opt -nodffe -nosdff
  fsm
  opt
  wreduce
  peepopt
  opt_clean
  share
  techmap -map +/cmp2lut.v -D LUT_WIDTH=6
  opt_expr
  opt_clean
  techmap -map +/mul2dsp.v [...]      (unless -nodsp)
  alumacc
  iopadmap -bits -outpad MISTRAL_OB I:PAD -inpad MISTRAL_IB O:PAD -toutpad MISTRAL_IO OE:O:PAD -tinou
(unless -noiopad)
  techmap -map +/intel_alm/common/arith_alm_map.v -map +/intel_alm/common/dsp_map.v
  opt
  memory -nomap
  opt_clean

map_bram:      (skip if -nobram)
  memory_bram -rules +/intel_alm/common/bram_<bram_type>.txt
  techmap -map +/intel_alm/common/bram_<bram_type>_map.v

map_lutram:    (skip if -nolutram)
  memory_bram -rules +/intel_alm/common/lutram_mlab.txt      (for Cyclone V / Cyclone 10GX)

map_ffram:
  memory_map
  opt -full

map_ffs:
  techmap

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
98     dfflegalize -cell $_DFFE_PN0P_ 0 -cell $_SDFACE_PP0P_ 0
99     techmap -map +/intel_alm/common/dff_map.v
100     opt -full -undriven -mux_undef
101     clean -purge
102     clkbufmap -buf MISTRAL_CLKBUF Q:A      (unless -noclkbuf)
103
104     map_luts:
105         techmap -map +/intel_alm/common/abc9_map.v
106         abc9 [-dff] -maxlut 6 -W 600
107         techmap -map +/intel_alm/common/abc9_unmap.v
108         techmap -map +/intel_alm/common/alm_map.v
109         opt -fast
110         autoname
111         clean
112
113     check:
114         hierarchy -check
115         stat
116         check
117         blackbox =A:whitebox
118
119     quartus:
120         rename -hide w:[* w:]*
121         setundef -zero
122         hilomap -singleton -hicell __MISTRAL_VCC Q -locell __MISTRAL_GND Q
123         techmap -D <family> -map +/intel_alm/common/quartus_rename.v
124
125     vqm:
126         write_verilog -attr2comment -defparam -nohex -decimal <file-name>
```

### C.203 synth\_machxo2 – synthesis for MachXO2 FPGAs. This work is experimental.

```
1     synth_machxo2 [options]
2
3     This command runs synthesis for MachXO2 FPGAs.
4
5     -top <module>
6         use the specified module as top module
7
8     -blif <file>
9         write the design to the specified BLIF file. writing of an output file
10        is omitted if this parameter is not specified.
11
12     -edif <file>
13         write the design to the specified EDIF file. writing of an output file
14        is omitted if this parameter is not specified.
15
16     -json <file>
17         write the design to the specified JSON file. writing of an output file
18        is omitted if this parameter is not specified.
19
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

20  -run <from_label>:<to_label>
21      only run the commands between the labels (see below). an empty
22      from label is synonymous to 'begin', and empty to label is
23      synonymous to the end of the command list.
24
25  -nobram
26      do not use block RAM cells in output netlist
27
28  -nolutram
29      do not use LUT RAM cells in output netlist
30
31  -noflatten
32      do not flatten design before synthesis
33
34  -noiopad
35      do not insert IO buffers
36
37  -vpr
38      generate an output netlist (and BLIF file) suitable for VPR
39      (this feature is experimental and incomplete)
40
41
42  The following commands are executed by this synthesis command:
43
44  begin:
45      read_verilog -lib -icells +/machxo2/cells_sim.v
46      hierarchy -check -top <top>
47
48  flatten:    (unless -noflatten)
49      proc
50      flatten
51      tribuf -logic
52      deminout
53
54  coarse:
55      synth -run coarse
56
57  map_ram:
58      memory_libmap -lib +/machxo2/lutrams.txt -lib +/machxo2/brams.txt [-no-auto-block] [-no-auto-distrib
(-no-auto-block if -nobram, -no-auto-distributed if -nolutram)
59      techmap -map +/machxo2/lutrams_map.v -map +/machxo2/brams_map.v
60
61  fine:
62      memory_map
63      opt -full
64      techmap -map +/techmap.v
65      opt -fast
66
67  map_ios:    (unless -noiopad)
68      iopadmap -bits -outpad $__FACADE_OUTPAD I:0 -inpad $__FACADE_INPAD 0:I -toutpad $__FACADE_TOUTPAD ~
69      attrmvcp -attr src -attr LOC t:$__FACADE_OUTPAD %x:[0] t:$__FACADE_TOUTPAD %x:[0] t:$__FACADE_TIN
70      attrmvcp -attr src -attr LOC -driven t:$__FACADE_INPAD %x:[I]
71
72  map_ffs:

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
73     dfflegalize -cell $_DFF_P_ 0
74
75 map_luts:
76     abc -lut 4 -dress
77     clean
78
79 map_cells:
80     techmap -map +/machxo2/cells_map.v
81     clean
82
83 check:
84     hierarchy -check
85     stat
86     blackbox =A:whitebox
87
88 blif:
89     opt_clean -purge                                (vpr mode)
90     write_blif -attr -cname -conn -param <file-name> (vpr mode)
91     write_blif -gates -attr -param <file-name>      (non-vpr mode)
92
93 edif:
94     write_edif <file-name>
95
96 json:
97     write_json <file-name>
```

### C.204 synth\_nexus – synthesis for Lattice Nexus FPGAs

```
1     synth_nexus [options]
2
3 This command runs synthesis for Lattice Nexus FPGAs.
4
5 -top <module>
6     use the specified module as top module
7
8 -family <device>
9     run synthesis for the specified Nexus device
10    supported values: lifcl, lfd2nx
11
12 -json <file>
13    write the design to the specified JSON file. writing of an output file
14    is omitted if this parameter is not specified.
15
16 -vm <file>
17    write the design to the specified structural Verilog file. writing of
18    an output file is omitted if this parameter is not specified.
19
20 -run <from_label>:<to_label>
21    only run the commands between the labels (see below). an empty
22    from label is synonymous to 'begin', and empty to label is
23    synonymous to the end of the command list.
24
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
25 -noflatten
26     do not flatten design before synthesis
27
28 -dff
29     run 'abc'/'abc9' with -dff option
30
31 -retime
32     run 'abc' with '-dff -D 1' options
33
34 -noccu2
35     do not use CCU2 cells in output netlist
36
37 -nodffe
38     do not use flipflops with CE in output netlist
39
40 -nolram
41     do not use large RAM cells in output netlist
42     note that large RAM must be explicitly requested with a (* lram *)
43     attribute on the memory.
44
45 -nobram
46     do not use block RAM cells in output netlist
47
48 -nolutram
49     do not use LUT RAM cells in output netlist
50
51 -nowidelut
52     do not use PFU muxes to implement LUTs larger than LUT4s
53
54 -noiopad
55     do not insert IO buffers
56
57 -nodsp
58     do not infer DSP multipliers
59
60 -abc9
61     use new ABC9 flow (EXPERIMENTAL)
62
63 The following commands are executed by this synthesis command:
64
65 begin:
66     read_verilog -lib -specify +/nexus/cells_sim.v +/nexus/cells_xtra.v
67     hierarchy -check -top <top>
68
69 coarse:
70     proc
71     flatten
72     tribuf -logic
73     deminout
74     opt_expr
75     opt_clean
76     check
77     opt -nodffe -nosdff
78     fsm
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

79     opt
80     wreduce
81     peepopt
82     opt_clean
83     share
84     techmap -map +/cmp2lut.v -D LUT_WIDTH=4
85     opt_expr
86     opt_clean
87     techmap -map +/mul2dsp.v [...]      (unless -nodsp)
88     techmap -map +/nexus/dsp_map.v      (unless -nodsp)
89     alumacc
90     opt
91     memory -nomap
92     opt_clean
93
94     map_ram:
95         memory_libmap -lib +/nexus/lutrams.txt -lib +/nexus/brams.txt -lib +/nexus/lrams.txt -no-auto-huge
(-no-auto-block if -nobram, -no-auto-distributed if -nolutram)
96         techmap -map +/nexus/lutrams_map.v -map +/nexus/brams_map.v -map +/nexus/lrams_map.v
97
98     map_ffram:
99         opt -fast -mux_undef -undriven -fine
100        memory_map
101        opt -undriven -fine
102
103     map_gates:
104         techmap -map +/techmap.v -map +/nexus/arith_map.v
105         iopadmap -bits -outpad OB I:0 -inpad IB 0:I -toutpad OBZ ~T:I:0 -tinoutpad BB ~T:0:I:B A:top
(skip if '-noiopad')
106         opt -fast
107         abc -dff -D 1      (only if -retime)
108
109     map_ffs:
110         opt_clean
111         dfflegalize -cell $_DFF_P_ 01 -cell $_DFF_PP?_ r -cell $_SDFF_PP?_ r -cell $_DLATCH?_ x [-cell $_D
($_*DFFE_* only if not -nodffe)
112         zinit -all w:* t:$_DFF?_ t:$_DFFE??_ t:$_SDFF*      (only if -abc9 and -dff
113         techmap -D NO_LUT -map +/nexus/cells_map.v
114         opt_expr -undriven -mux_undef
115         simplemap
116         attrmvp -copy -attr syn_useioff
117         opt_clean
118
119     map_luts:
120         techmap -map +/nexus/latches_map.v
121         abc -dress -lut 4:5
122         clean
123
124     map_cells:
125         techmap -map +/nexus/cells_map.v
126         setundef -zero
127         hilomap -singleton -hicell VHI Z -locell VLO Z
128         clean
129

```

```

130     check:
131         autoname
132         hierarchy -check
133         stat
134         check -noinit
135         blackbox =A:whitebox
136
137     json:
138         write_json <file-name>
139
140     vm:
141         write_verilog <file-name>

```

## C.205 synth\_quicklogic – Synthesis for QuickLogic FPGAs

```

1     synth_quicklogic [options]
2     This command runs synthesis for QuickLogic FPGAs
3
4     -top <module>
5         use the specified module as top module
6
7     -family <family>
8         run synthesis for the specified QuickLogic architecture
9         generate the synthesis netlist for the specified family.
10        supported values:
11        - pp3: PolarPro 3
12
13     -blif <file>
14         write the design to the specified BLIF file. writing of an output file
15         is omitted if this parameter is not specified.
16
17     -verilog <file>
18         write the design to the specified verilog file. writing of an output file
19         is omitted if this parameter is not specified.
20
21     -abc
22         use old ABC flow, which has generally worse mapping results but is less
23         likely to have bugs.
24
25     The following commands are executed by this synthesis command:
26
27     begin:
28         read_verilog -lib -specify +/quicklogic/cells_sim.v +/quicklogic/pp3_cells_sim.v
29         read_verilog -lib -specify +/quicklogic/lut_sim.v
30         hierarchy -check -top <top>
31
32     coarse:
33         proc
34         flatten
35         tribuf -logic
36         deminout
37         opt_expr

```



## APPENDIX C. COMMAND REFERENCE MANUAL

```

38     opt_clean
39     check
40     opt -nodffe -nosdff
41     fsm
42     opt
43     wreduce
44     peepopt
45     opt_clean
46     share
47     techmap -map +/cmp2lut.v -D LUT_WIDTH=4
48     opt_expr
49     opt_clean
50     alumacc
51     pmuxtree
52     opt
53     memory -nomap
54     opt_clean
55
56 map_ffram:
57     opt -fast -mux_undef -undriven -fine
58     memory_map -iattr -attr !ram_block -attr !rom_block -attr logic_block -attr syn_ramstyle=auto -attr
59     opt -undriven -fine
60
61 map_gates:
62     techmap
63     opt -fast
64     muxcover -mux8 -mux4
65
66 map_ffs:
67     opt_expr
68     dfflegalize -cell $_DFFSRE_PPPP_ 0 -cell $_DLATCH_?_ x
69     techmap -map +/quicklogic/pp3_cells_map.v -map +/quicklogic/pp3_ffs_map.v
70     opt_expr -mux_undef
71
72 map_luts:
73     techmap -map +/quicklogic/pp3_latches_map.v
74     read_verilog -lib -specify -icells +/quicklogic/abc9_model.v
75     techmap -map +/quicklogic/abc9_map.v
76     abc9 -maxlut 4 -dff
77     techmap -map +/quicklogic/abc9_unmap.v
78     clean
79
80 map_cells:
81     techmap -map +/quicklogic/pp3_lut_map.v
82     clean
83
84 check:
85     autaname
86     hierarchy -check
87     stat
88     check -noinit
89
90 iomap:
91     clkbufmap -inpad ckpad Q:P

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
92         iopadmap -bits -outpad outpad A:P -inpad inpad Q:P -tinoutpad bipad EN:Q:A:P A:top
93
94     finalize:
95         setundef -zero -params -undriven
96         hilomap -hicell logic_1 A -locell logic_0 A -singleton A:top
97         opt_clean -purge
98         check
99         blackbox =A:whitebox
100
101     blif:
102         write_blif -attr -param -auto-top
103
104     verilog:
105         write_verilog -noattr -nohex <file-name>
```

### C.206 synth\_sf2 – synthesis for SmartFusion2 and IGLOO2 FPGAs

```
1     synth_sf2 [options]
2
3     This command runs synthesis for SmartFusion2 and IGLOO2 FPGAs.
4
5     -top <module>
6         use the specified module as top module
7
8     -edif <file>
9         write the design to the specified EDIF file. writing of an output file
10        is omitted if this parameter is not specified.
11
12     -vlog <file>
13        write the design to the specified Verilog file. writing of an output file
14        is omitted if this parameter is not specified.
15
16     -json <file>
17        write the design to the specified JSON file. writing of an output file
18        is omitted if this parameter is not specified.
19
20     -run <from_label>:<to_label>
21        only run the commands between the labels (see below). an empty
22        from label is synonymous to 'begin', and empty to label is
23        synonymous to the end of the command list.
24
25     -noflatten
26        do not flatten design before synthesis
27
28     -noiobs
29        run synthesis in "block mode", i.e. do not insert IO buffers
30
31     -clkbuf
32        insert direct PAD->global_net buffers
33
34     -retime
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

35         run 'abc' with '-dff -D 1' options
36
37
38 The following commands are executed by this synthesis command:
39
40     begin:
41         read_verilog -lib +/sf2/cells_sim.v
42         hierarchy -check -top <top>
43
44     flatten:      (unless -noflatten)
45         proc
46         flatten
47         tribuf -logic
48         deminout
49
50     coarse:
51         synth -run coarse
52
53     fine:
54         opt -fast -mux_undef -undriven -fine
55         memory_map
56         opt -undriven -fine
57         techmap -map +/techmap.v -map +/sf2/arith_map.v
58         opt -fast
59         abc -dff -D 1      (only if -retime)
60
61     map_ffs:
62         dfflegalize -cell $_DFFE_PN?P_ x -cell $_SDFFC_PN?P_ x -cell $_DLATCH_PN?_ x
63         techmap -D NO_LUT -map +/sf2/cells_map.v
64         opt_expr -mux_undef
65         simplemap
66
67     map_luts:
68         abc -lut 4
69         clean
70
71     map_cells:
72         techmap -map +/sf2/cells_map.v
73         clean
74
75     map_iobs:
76         clkbufmap -buf CLKINT Y:A [-inpad CLKBUF Y:PAD]      (unless -noiobs, -inpad only passed if -clkbuf)
77         iopadmap -bits -inpad INBUF Y:PAD -outpad OUTBUF D:PAD -toutpad TRIBUFF E:D:PAD -tinoutpad BIBUF E:
(unless -noiobs
78         clean
79
80     check:
81         hierarchy -check
82         stat
83         check -noinit
84         blackbox =A:whitebox
85
86     edif:
87         write_edif -gndvccy <file-name>

```

```

88
89     vlog:
90         write_verilog <file-name>
91
92     json:
93         write_json <file-name>

```

## C.207 synth\_xilinx – synthesis for Xilinx FPGAs

```

1     synth_xilinx [options]
2
3     This command runs synthesis for Xilinx FPGAs. This command does not operate on
4     partly selected designs. At the moment this command creates netlists that are
5     compatible with 7-Series Xilinx devices.
6
7     -top <module>
8         use the specified module as top module
9
10    -family <family>
11        run synthesis for the specified Xilinx architecture
12        generate the synthesis netlist for the specified family.
13        supported values:
14        - xcup: Ultrascale Plus
15        - xcu: Ultrascale
16        - xc7: Series 7 (default)
17        - xc6s: Spartan 6
18        - xc6v: Virtex 6
19        - xc5v: Virtex 5 (EXPERIMENTAL)
20        - xc4v: Virtex 4 (EXPERIMENTAL)
21        - xc3sda: Spartan 3A DSP (EXPERIMENTAL)
22        - xc3sa: Spartan 3A (EXPERIMENTAL)
23        - xc3se: Spartan 3E (EXPERIMENTAL)
24        - xc3s: Spartan 3 (EXPERIMENTAL)
25        - xc2vp: Virtex 2 Pro (EXPERIMENTAL)
26        - xc2v: Virtex 2 (EXPERIMENTAL)
27        - xcve: Virtex E, Spartan 2E (EXPERIMENTAL)
28        - xcv: Virtex, Spartan 2 (EXPERIMENTAL)
29
30    -edif <file>
31        write the design to the specified edif file. writing of an output file
32        is omitted if this parameter is not specified.
33
34    -blif <file>
35        write the design to the specified BLIF file. writing of an output file
36        is omitted if this parameter is not specified.
37
38    -ise
39        generate an output netlist suitable for ISE
40
41    -nobram
42        do not use block RAM cells in output netlist
43

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
44 -nolutram
45     do not use distributed RAM cells in output netlist
46
47 -nosrl
48     do not use distributed SRL cells in output netlist
49
50 -nocarry
51     do not use XORCY/MUXCY/CARRY4 cells in output netlist
52
53 -nowidelut
54     do not use MUXF[5-9] resources to implement LUTs larger than native for the target
55
56 -nodsp
57     do not use DSP48*s to implement multipliers and associated logic
58
59 -noiopad
60     disable I/O buffer insertion (useful for hierarchical or
61     out-of-context flows)
62
63 -noclkbuf
64     disable automatic clock buffer insertion
65
66 -uram
67     infer URAM288s for large memories (xcup only)
68
69 -widemux <int>
70     enable inference of hard multiplexer resources (MUXF[78]) for muxes at or
71     above this number of inputs (minimum value 2, recommended value >= 5).
72     default: 0 (no inference)
73
74 -run <from_label>:<to_label>
75     only run the commands between the labels (see below). an empty
76     from label is synonymous to 'begin', and empty to label is
77     synonymous to the end of the command list.
78
79 -flatten
80     flatten design before synthesis
81
82 -dff
83     run 'abc'/'abc9' with -dff option
84
85 -retime
86     run 'abc' with '-D 1' option to enable flip-flop retiming.
87     implies -dff.
88
89 -abc9
90     use new ABC9 flow (EXPERIMENTAL)
```

The following commands are executed by this synthesis command:

```
95 begin:
96     read_verilog -lib -specify +/xilinx/cells_sim.v
97     read_verilog -lib +/xilinx/cells_extra.v
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

98     hierarchy -check -auto-top
99
100 prepare:
101     proc
102     flatten      (with '-flatten')
103     tribuf -logic
104     deminout
105     opt_expr
106     opt_clean
107     check
108     opt -nodffe -nosdff
109     fsm
110     opt
111     wreduce [-keepdc]      (option for '-widemux')
112     peepopt
113     opt_clean
114     muxpack          ('-widemux' only)
115     pmux2shiftx      (skip if '-nosrl' and '-widemux=0')
116     clean            (skip if '-nosrl' and '-widemux=0')
117
118 map_dsp:      (skip if '-nodsp')
119     memory_dff
120     techmap -map +/mul2dsp.v -map +/xilinx/{family}_dsp_map.v {options}
121     select a:mul2dsp
122     setattr -unset mul2dsp
123     opt_expr -fine
124     wreduce
125     select -clear
126     xilinx_dsp -family <family>
127     chtype -set $mul t:$_soft_mul
128
129 coarse:
130     techmap -map +/cmp2lut.v -map +/cmp2lcu.v -D LUT_WIDTH=[46]
131     alumacc
132     share
133     opt
134     memory -nomap
135     opt_clean
136
137 map_memory:
138     memory_libmap [...]
139     techmap -map +/xilinx/lutrams_<family>_map.v
140     techmap -map +/xilinx/brams_<family>_map.v
141
142 map_ffram:
143     opt -fast -full
144     memory_map
145
146 fine:
147     simplemap t:$mux      ('-widemux' only)
148     muxcover <internal options>      ('-widemux' only)
149     opt -full
150     xilinx_srl -variable -minlen 3      (skip if '-nosrl')
151     techmap -map +/techmap.v -D LUT_SIZE=[46] [-map +/xilinx/mux_map.v] -map +/xilinx/arith_map.v

```

```

152     opt -fast
153
154     map_cells:
155         iopadmap -bits -outpad OBUF I:0 -inpad IBUF 0:I -toutpad OBUFT ~T:I:0 -tinoutpad IOBUF ~T:0:I:IO A:
(skip if '-noiopad')
156         techmap -map +/techmap.v -map +/xilinx/cells_map.v
157         clean
158
159     map_ffs:
160         dfflegalize -cell $_DFFE_?P?P_ 01 -cell $_SDFFE_?P?P_ 01 -cell $_DLATCH_?P?_ 01      (for xc6v, xc7,
161         zinit -all w:* t:$_SDFFE_*      ('-dff' only)
162         techmap -map +/xilinx/ff_map.v      ('-abc9' only)
163
164     map_luts:
165         opt_expr -mux_undef -noclkinv
166         abc -luts 2:2,3,6:5[,10,20] [-dff] [-D 1]      (option for '-nowidelut', '-dff', '-retime')
167         clean
168         techmap -map +/xilinx/ff_map.v      (only if not '-abc9')
169         xilinx_srl -fixed -minlen 3      (skip if '-nosrl')
170         techmap -map +/xilinx/lut_map.v -map +/xilinx/cells_map.v -D LUT_WIDTH=[46]
171         xilinx_dffopt [-lut4]
172         opt_lut_ins -tech xilinx
173
174     finalize:
175         clkbufmap -buf BUFG 0:I      (skip if '-noclkbuf')
176         extractinv -inv INV 0:I      (only if '-ise')
177         clean
178
179     check:
180         hierarchy -check
181         stat -tech xilinx
182         check -noinit
183         blackbox =A:whitebox
184
185     edif:
186         write_edif -pvector bra
187
188     blif:
189         write_blif

```

## C.208 tcl – execute a TCL script file

```

1     tcl <filename> [args]
2
3     This command executes the tcl commands in the specified file.
4     Use 'yosys cmd' to run the yosys command 'cmd' from tcl.
5
6     The tcl command 'yosys -import' can be used to import all yosys
7     commands directly as tcl commands to the tcl shell. Yosys commands
8     'proc' and 'rename' are wrapped to tcl commands 'procs' and 'renames'
9     in order to avoid a name collision with the built in commands.
10

```

11 If any arguments are specified, these arguments are provided to the script via  
 12 the standard \$argc and \$argv variables.

## C.209 techmap – generic technology mapper

```

1  techmap [-map filename] [selection]
2
3  This pass implements a very simple technology mapper that replaces cells in
4  the design with implementations given in form of a Verilog or RTLIL source
5  file.
6
7  -map filename
8      the library of cell implementations to be used.
9      without this parameter a builtin library is used that
10     transforms the internal RTL cells to the internal gate
11     library.
12
13  -map %<design-name>
14     like -map above, but with an in-memory design instead of a file.
15
16  -extern
17     load the cell implementations as separate modules into the design
18     instead of inlining them.
19
20  -max_iter <number>
21     only run the specified number of iterations on each module.
22     default: unlimited
23
24  -recursive
25     instead of the iterative breadth-first algorithm use a recursive
26     depth-first algorithm. both methods should yield equivalent results,
27     but may differ in performance.
28
29  -autoproc
30     Automatically call "proc" on implementations that contain processes.
31
32  -wb
33     Ignore the 'whitebox' attribute on cell implementations.
34
35  -assert
36     this option will cause techmap to exit with an error if it can't map
37     a selected cell. only cell types that end on an underscore are accepted
38     as final cell types by this mode.
39
40  -D <define>, -I <incdir>
41     this options are passed as-is to the Verilog frontend for loading the
42     map file. Note that the Verilog frontend is also called with the
43     '-nooverwrite' option set.
44
45  When a module in the map file has the 'techmap_celltype' attribute set, it will
46  match cells with a type that match the text value of this attribute. Otherwise
47  the module name will be used to match the cell. Multiple space-separated cell

```



## APPENDIX C. COMMAND REFERENCE MANUAL

types can be listed, and wildcards using [] will be expanded (ie. "\$\_DFF\_[PN]\_" is the same as "\$\_DFF\_P\_ \$\_DFF\_N\_").

When a module in the map file has the 'techmap\_simplemap' attribute set, techmap will use 'simplemap' (see 'help simplemap') to map cells matching the module.

When a module in the map file has the 'techmap\_maccmap' attribute set, techmap will use 'maccmap' (see 'help maccmap') to map cells matching the module.

When a module in the map file has the 'techmap\_wrap' attribute set, techmap will create a wrapper for the cell and then run the command string that the attribute is set to on the wrapper module.

When a port on a module in the map file has the 'techmap\_autopurge' attribute set, and that port is not connected in the instantiation that is mapped, then a cell port connected only to such wires will be omitted in the mapped version of the circuit.

All wires in the modules from the map file matching the pattern \_TECHMAP\_\* or \*\_TECHMAP\_\* are special wires that are used to pass instructions from the mapping module to the techmap command. At the moment the following special wires are supported:

### \_TECHMAP\_FAIL\_

When this wire is set to a non-zero constant value, techmap will not use this module and instead try the next module with a matching 'techmap\_celltype' attribute.

When such a wire exists but does not have a constant value after all \_TECHMAP\_DO\_\* commands have been executed, an error is generated.

### \_TECHMAP\_DO\_\*

This wires are evaluated in alphabetical order. The constant text value of this wire is a yosys command (or sequence of commands) that is run by techmap on the module. A common use case is to run 'proc' on modules that are written using always-statements.

When such a wire has a non-constant value at the time it is to be evaluated, an error is produced. That means it is possible for such a wire to start out as non-constant and evaluate to a constant value during processing of other \_TECHMAP\_DO\_\* commands.

A \_TECHMAP\_DO\_\* command may start with the special token 'CONSTMAP; '. in this case techmap will create a copy for each distinct configuration of constant inputs and shorted inputs at this point and import the constant and connected bits into the map module. All further commands are executed in this copy. This is a very convenient way of creating optimized specializations of techmap modules without using the special parameters described below.

A \_TECHMAP\_DO\_\* command may start with the special token 'RECURSION; '. then techmap will recursively replace the cells in the module with their implementation. This is not affected by the -max\_iter option.

## APPENDIX C. COMMAND REFERENCE MANUAL

It is possible to combine both prefixes to 'RECURSION; CONSTMAP; '.

`_TECHMAP_REMOVEINIT_<port-name>_`  
When this wire is set to a constant value, the init attribute of the wire(s) connected to this port will be consumed. This wire must have the same width as the given port, and for every bit that is set to 1 in the value, the corresponding init attribute bit will be changed to 1'bx. If all bits of an init attribute are left as x, it will be removed.

In addition to this special wires, techmap also supports special parameters in modules in the map file:

`_TECHMAP_CELLTYPE_`  
When a parameter with this name exists, it will be set to the type name of the cell that matches the module.

`_TECHMAP_CELLNAME_`  
When a parameter with this name exists, it will be set to the name of the cell that matches the module.

`_TECHMAP_CONSTMSK_<port-name>_`  
`_TECHMAP_CONSTVAL_<port-name>_`  
When this pair of parameters is available in a module for a port, then former has a 1-bit for each constant input bit and the latter has the value for this bit. The unused bits of the latter are set to undef (x).

`_TECHMAP_WIREINIT_<port-name>_`  
When a parameter with this name exists, it will be set to the initial value of the wire(s) connected to the given port, as specified by the init attribute. If the attribute doesn't exist, x will be filled for the missing bits. To remove the init attribute bits used, use the `_TECHMAP_REMOVEINIT_*_` wires.

`_TECHMAP_BITS_CONNMAP_`  
`_TECHMAP_CONNMAP_<port-name>_`  
For an N-bit port, the `_TECHMAP_CONNMAP_<port-name>_` parameter, if it exists, will be set to an `N*_TECHMAP_BITS_CONNMAP_` bit vector containing N words (of `_TECHMAP_BITS_CONNMAP_` bits each) that assign each single bit driver a unique id. The values 0-3 are reserved for 0, 1, x, and z. This can be used to detect shorted inputs.

When a module in the map file has a parameter where the according cell in the design has a port, the module from the map file is only used if the port in the design is connected to a constant value. The parameter is then set to the constant value.

A cell with the name `_TECHMAP_REPLACE_` in the map file will inherit the name and attributes of the cell that is being replaced.

A cell with a name of the form ``_TECHMAP_REPLACE_.<suffix>`` in the map file will be named thus but with the ``_TECHMAP_REPLACE_`` prefix substituted with the name of the cell being replaced.

Similarly, a wire named in the form ``_TECHMAP_REPLACE_.<suffix>`` will cause a new wire alias to be created and named as above but with the ``_TECHMAP_REPLACE_`` prefix also substituted.

```

156
157 See 'help extract' for a pass that does the opposite thing.
158
159 See 'help flatten' for a pass that does flatten the design (which is
160 essentially techmap but using the design itself as map library).

```

## C.210 tee – redirect command output to file

```

1      tee [-q] [-o logfile|-a logfile] cmd
2
3      Execute the specified command, optionally writing the commands output to the
4      specified logfile(s).
5
6      -q
7          Do not print output to the normal destination (console and/or log file).
8
9      -o logfile
10         Write output to this file, truncate if exists.
11
12      -a logfile
13         Write output to this file, append if exists.
14
15      +INT, -INT
16         Add/subtract INT from the -v setting for this command.

```

## C.211 test\_abcloop – automatically test handling of loops in abc command

```

1      test_abcloop [options]
2
3      Test handling of logic loops in ABC.
4
5      -n {integer}
6          create this number of circuits and test them (default = 100).
7
8      -s {positive_integer}
9          use this value as rng seed value (default = unix time).

```

## C.212 test\_autotb – generate simple test benches

```

1      test_autotb [options] [filename]
2
3      Automatically create primitive Verilog test benches for all modules in the
4      design. The generated testbenches toggle the input pins of the module in
5      a semi-random manner and dumps the resulting output signals.
6
7      This can be used to check the synthesis results for simple circuits by

```

```

8 comparing the testbench output for the input files and the synthesis results.
9
10 The backend automatically detects clock signals. Additionally a signal can
11 be forced to be interpreted as clock signal by setting the attribute
12 'gentb_clock' on the signal.
13
14 The attribute 'gentb_constant' can be used to force a signal to a constant
15 value after initialization. This can e.g. be used to force a reset signal
16 low in order to explore more inner states in a state machine.
17
18 The attribute 'gentb_skip' can be attached to modules to suppress testbench
19 generation.
20
21 -n <int>
22     number of iterations the test bench should run (default = 1000)
23
24 -seed <int>
25     seed used for pseudo-random number generation (default = 0).
26     a value of 0 will cause an arbitrary seed to be chosen, based on
27     the current system time.

```

### C.213 test\_cell – automatically test the implementation of a cell type

```

1 test_cell [options] {cell-types}
2
3 Tests the internal implementation of the given cell type (for example '$add')
4 by comparing SAT solver, EVAL and TECHMAP implementations of the cell types..
5
6 Run with 'all' instead of a cell type to run the test on all supported
7 cell types. Use for example 'all /$add' for all cell types except $add.
8
9 -n {integer}
10     create this number of cell instances and test them (default = 100).
11
12 -s {positive_integer}
13     use this value as rng seed value (default = unix time).
14
15 -f {rtlil_file}
16     don't generate circuits. instead load the specified RTLIL file.
17
18 -w {filename_prefix}
19     don't test anything. just generate the circuits and write them
20     to RTLIL files with the specified prefix
21
22 -map {filename}
23     pass this option to techmap.
24
25 -simlib
26     use "techmap -D SIMLIB_NOCHECKS -map +/simlib.v -max_iter 2 -autoproc"
27
28 -aigmap

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
29         instead of calling "techmap", call "aigmap"
30
31     -muxdiv
32         when creating test benches with dividers, create an additional mux
33         to mask out the division-by-zero case
34
35     -script {script_file}
36         instead of calling "techmap", call "script {script_file}".
37
38     -const
39         set some input bits to random constant values
40
41     -nosat
42         do not check SAT model or run SAT equivalence checking
43
44     -noeval
45         do not check const-eval models
46
47     -edges
48         test cell edges db creator against sat-based implementation
49
50     -v
51         print additional debug information to the console
52
53     -vlog {filename}
54         create a Verilog test bench to test simlib and write_verilog
```

### C.214 test\_pmgen – test pass for pmgen

```
1     test_pmgen -reduce_chain [options] [selection]
2
3     Demo for recursive pmgen patterns. Map chains of AND/OR/XOR to $reduce_*.
4
5
6     test_pmgen -reduce_tree [options] [selection]
7
8     Demo for recursive pmgen patterns. Map trees of AND/OR/XOR to $reduce_*.
9
10
11     test_pmgen -eqpmux [options] [selection]
12
13     Demo for recursive pmgen patterns. Optimize EQ/NE/PMUX circuits.
14
15
16     test_pmgen -generate [options] <pattern_name>
17
18     Create modules that match the specified pattern.
```

### C.215 torder – print cells in topological order

```

1  torder [options] [selection]
2
3  This command prints the selected cells in topological order.
4
5  -stop <cell_type> <cell_port>
6      do not use the specified cell port in topological sorting
7
8  -noautostop
9      by default Q outputs of internal FF cells and memory read port outputs
10     are not used in topological sorting. this option deactivates that.

```

## C.216 trace – redirect command output to file

```

1  trace cmd
2
3  Execute the specified command, logging all changes the command performs on
4  the design in real time.

```

## C.217 tribuf – infer tri-state buffers

```

1  tribuf [options] [selection]
2
3  This pass transforms $mux cells with 'z' inputs to tristate buffers.
4
5  -merge
6      merge multiple tri-state buffers driving the same net
7      into a single buffer.
8
9  -logic
10     convert tri-state buffers that do not drive output ports
11     to non-tristate logic. this option implies -merge.
12
13  -formal
14     convert all tri-state buffers to non-tristate logic and
15     add a formal assertion that no two buffers are driving the
16     same net simultaneously. this option implies -merge.

```

## C.218 uniquify – create unique copies of modules

```

1  uniquify [selection]
2
3  By default, a module that is instantiated by several other modules is only
4  kept once in the design. This preserves the original modularity of the design
5  and reduces the overall size of the design in memory. But it prevents certain
6  optimizations and other operations on the design. This pass creates unique
7  modules for all selected cells. The created modules are marked with the

```

```

8 'unique' attribute.
9
10 This commands only operates on modules that by themselves have the 'unique'
11 attribute set (the 'top' module is unique implicitly).

```

## C.219 `verific` – load Verilog and VHDL designs using Verific

```

1      verific {-vlog95|-vlog2k|-sv2005|-sv2009|-sv2012|-sv} <verilog-file>..
2
3 Load the specified Verilog/SystemVerilog files into Verific.
4
5 All files specified in one call to this command are one compilation unit.
6 Files passed to different calls to this command are treated as belonging to
7 different compilation units.
8
9 Additional -D<macro>[=<value>] options may be added after the option indicating
10 the language version (and before file names) to set additional verilog defines.
11 The macros YOSYS, SYNTHESIS, and VERIFIC are defined implicitly.
12
13
14      verific -formal <verilog-file>..
15
16 Like -sv, but define FORMAL instead of SYNTHESIS.
17
18
19      verific {-vhd187|-vhd193|-vhd12k|-vhd12008|-vhd1} <vhdl-file>..
20
21 Load the specified VHDL files into Verific.
22
23
24      verific {-f|-F} [-vlog95|-vlog2k|-sv2005|-sv2009|
25                    -sv2012|-sv|-formal] <command-file>
26
27 Load and execute the specified command file.
28 Override verilog parsing mode can be set.
29 The macros YOSYS, SYNTHESIS/FORMAL, and VERIFIC are defined implicitly.
30
31 Command file parser supports following commands in file:
32 +define+<MACRO>=<VALUE> - defines macro
33 -u                      - upper case all identifier (makes Verilog parser
34                          case insensitive)
35 -v <filepath>           - register library name (file)
36 -y <filepath>           - register library name (directory)
37 +incdir+<filepath>      - specify include dir
38 +libext+<filepath>      - specify library extension
39 +liborder+<id>          - add library in ordered list
40 +librescan              - unresolved modules will be always searched
41                          starting with the first library specified
42                          by -y/-v options.
43 -f/-file <filepath>     - nested -f option
44 -F <filepath>           - nested -F option (relative path)
45 parse files:

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

46     <filepath>
47     +systemverilogext+<filepath>
48     +verilog1995ext+<filepath>
49     +verilog2001ext+<filepath>
50
51     analysis mode:
52     -ams
53     +v2k
54     -sverilog
55
56
57     verific [-work <libname>] {-sv|-vhdl|...} <hdl-file>
58
59 Load the specified Verilog/SystemVerilog/VHDL file into the specified library.
60 (default library when -work is not present: "work")
61
62
63     verific [-L <libname>] {-sv|-vhdl|...} <hdl-file>
64
65 Look up external definitions in the specified library.
66 (-L may be used more than once)
67
68
69     verific -vlog-incdir <directory>..
70
71 Add Verilog include directories.
72
73
74     verific -vlog-libdir <directory>..
75
76 Add Verilog library directories. Verific will search in this directories to
77 find undefined modules.
78
79
80     verific -vlog-libext <extension>..
81
82 Add Verilog library extensions, used when searching in library directories.
83
84
85     verific -vlog-define <macro>[=<value>]..
86
87 Add Verilog defines.
88
89
90     verific -vlog-undef <macro>..
91
92 Remove Verilog defines previously set with -vlog-define.
93
94
95     verific -set-error <msg_id>..
96     verific -set-warning <msg_id>..
97     verific -set-info <msg_id>..
98     verific -set-ignore <msg_id>..
99

```



## APPENDIX C. COMMAND REFERENCE MANUAL

Set message severity. <msg\_id> is the string in square brackets when a message is printed, such as VERI-1209.

```
verific -import [options] <top-module>..
```

Elaborate the design for the specified top modules, import to Yosys and reset the internal state of Verific.

Import options:

-all

Elaborate all modules, not just the hierarchy below the given top modules. With this option the list of modules to import is optional.

-gates

Create a gate-level netlist.

-flatten

Flatten the design in Verific before importing.

-extnets

Resolve references to external nets by adding module ports as needed.

-autocover

Generate automatic cover statements for all asserts

-fullinit

Keep all register initializations, even those for non-FF registers.

-chparam name value

Elaborate the specified top modules (all modules when -all given) using this parameter value. Modules on which this parameter does not exist will cause Verific to produce a VERI-1928 or VHDL-1676 message. This option can be specified multiple times to override multiple parameters. String values must be passed in double quotes ("").

-v, -vv

Verbose log messages. (-vv is even more verbose than -v.)

-pp <filename>

Pretty print design after elaboration to specified file.

The following additional import options are useful for debugging the Verific bindings (for Yosys and/or Verific developers):

-k

Keep going after an unsupported verific primitive is found. The unsupported primitive is added as blockbox module to the design. This will also add all SVA related cells to the design parallel to the checker logic inferred by it.

-v

Import Verific netlist as-is without translating to Yosys cell types.

```

154
155 -nosva
156     Ignore SVA properties, do not infer checker logic.
157
158 -L <int>
159     Maximum number of ctrl bits for SVA checker FSMs (default=16).
160
161 -n
162     Keep all Verific names on instances and nets. By default only
163     user-declared names are preserved.
164
165 -d <dump_file>
166     Dump the Verific netlist as a verilog file.
167
168
169     verific [-work <libname>] -pp [options] <filename> [<module>]..
170
171 Pretty print design (or just module) to the specified file from the
172 specified library. (default library when -work is not present: "work")
173
174 Pretty print options:
175
176 -verilog
177     Save output for Verilog/SystemVerilog design modules (default).
178
179 -vhdl
180     Save output for VHDL design units.
181
182
183     verific -cfg [<name> [<value>]]
184
185 Get/set Verific runtime flags.
186
187
188 Use YosysHQ Tabby CAD Suite if you need Yosys+Verific.
189 https://www.yosyshq.com/
190
191 Contact office@yosyshq.com for free evaluation
192 binaries of YosysHQ Tabby CAD Suite.

```

## C.220 verilog\_defaults – set default options for read\_verilog

```

1     verilog_defaults -add [options]
2
3 Add the specified options to the list of default options to read_verilog.
4
5
6     verilog_defaults -clear
7
8 Clear the list of Verilog default options.
9
10

```

```

11     verilog_defaults -push
12     verilog_defaults -pop
13
14 Push or pop the list of default options to a stack. Note that -push does
15 not imply -clear.

```

## C.221 verilog\_defines – define and undefine verilog defines

```

1     verilog_defines [options]
2
3 Define and undefine verilog preprocessor macros.
4
5     -Dname[=definition]
6         define the preprocessor symbol 'name' and set its optional value
7         'definition'
8
9     -Uname[=definition]
10        undefine the preprocessor symbol 'name'
11
12     -reset
13        clear list of defined preprocessor symbols
14
15     -list
16        list currently defined preprocessor symbols

```

## C.222 wbflip – flip the whitebox attribute

```

1     wbflip [selection]
2
3 Flip the whitebox attribute on selected cells. I.e. if it's set, unset it, and
4 vice-versa. Blackbox cells are not effected by this command.

```

## C.223 wreduce – reduce the word size of operations if possible

```

1     wreduce [options] [selection]
2
3 This command reduces the word size of operations. For example it will replace
4 the 32 bit adders in the following code with adders of more appropriate widths:
5
6     module test(input [3:0] a, b, c, output [7:0] y);
7         assign y = a + b + c + 1;
8     endmodule
9
10 Options:
11
12     -memx
13         Do not change the width of memory address ports. Use this options in

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
14     flows that use the 'memory_memx' pass.
15
16     -mux_undef
17         remove 'undef' inputs from $mux, $pmux and $_MUX_ cells
18
19     -keepdc
20         Do not optimize explicit don't-care values.
```

### C.224 write\_aiger – write design to AIGER file

```
1     write_aiger [options] [filename]
2
3     Write the current design to an AIGER file. The design must be flattened and
4     must not contain any cell types except $_AND_, $_NOT_, simple FF types,
5     $assert and $assume cells, and $initstate cells.
6
7     $assert and $assume cells are converted to AIGER bad state properties and
8     invariant constraints.
9
10    -ascii
11        write ASCII version of AIGER format
12
13    -zinit
14        convert FFs to zero-initialized FFs, adding additional inputs for
15        uninitialized FFs.
16
17    -miter
18        design outputs are AIGER bad state properties
19
20    -symbols
21        include a symbol table in the generated AIGER file
22
23    -map <filename>
24        write an extra file with port and latch symbols
25
26    -vmap <filename>
27        like -map, but more verbose
28
29    -no-startoffset
30        make indexes zero based, enable using map files with smt solvers.
31
32    -I, -O, -B, -L
33        If the design contains no input/output/assert/flip-flop then create one
34        dummy input/output/bad_state-pin or latch to make the tools reading the
35        AIGER file happy.
```

### C.225 write\_blif – write design to BLIF file

```
1     write_blif [options] [filename]
2
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
3 Write the current design to an BLIF file.
4
5     -top top_module
6         set the specified module as design top module
7
8     -buf <cell-type> <in-port> <out-port>
9         use cells of type <cell-type> with the specified port names for buffers
10
11     -unbuf <cell-type> <in-port> <out-port>
12         replace buffer cells with the specified name and port names with
13         a .names statement that models a buffer
14
15     -true <cell-type> <out-port>
16     -false <cell-type> <out-port>
17     -undef <cell-type> <out-port>
18         use the specified cell types to drive nets that are constant 1, 0, or
19         undefined. when '-' is used as <cell-type>, then <out-port> specifies
20         the wire name to be used for the constant signal and no cell driving
21         that wire is generated. when '+' is used as <cell-type>, then <out-port>
22         specifies the wire name to be used for the constant signal and a .names
23         statement is generated to drive the wire.
24
25     -noalias
26         if a net name is aliasing another net name, then by default a net
27         without fanout is created that is driven by the other net. This option
28         suppresses the generation of this nets without fanout.
29
30 The following options can be useful when the generated file is not going to be
31 read by a BLIF parser but a custom tool. It is recommended to not name the output
32 file *.blif when any of this options is used.
33
34     -icells
35         do not translate Yosys's internal gates to generic BLIF logic
36         functions. Instead create .subckt or .gate lines for all cells.
37
38     -gates
39         print .gate instead of .subckt lines for all cells that are not
40         instantiations of other modules from this design.
41
42     -conn
43         do not generate buffers for connected wires. instead use the
44         non-standard .conn statement.
45
46     -attr
47         use the non-standard .attr statement to write cell attributes
48
49     -param
50         use the non-standard .param statement to write cell parameters
51
52     -cname
53         use the non-standard .cname statement to write cell names
54
55     -iname, -iattr
56         enable -cname and -attr functionality for .names statements
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
57      (the .cname and .attr statements will be included in the BLIF
58      output after the truth table for the .names statement)
59
60      -blackbox
61          write blackbox cells with .blackbox statement.
62
63      -impltf
64          do not write definitions for the $true, $false and $undef wires.
```

### C.226 write\_btor – write design to BTOR file

```
1      write_btor [options] [filename]
2
3      Write a BTOR description of the current design.
4
5      -v
6          Add comments and indentation to BTOR output file
7
8      -s
9          Output only a single bad property for all asserts
10
11     -c
12         Output cover properties using 'bad' statements instead of asserts
13
14     -i <filename>
15         Create additional info file with auxiliary information
16
17     -x
18         Output symbols for internal netnames (starting with '$')
```

### C.227 write\_cxxrtl – convert design to C++ RTL simulation

```
1      write_cxxrtl [options] [filename]
2
3      Write C++ code that simulates the design. The generated code requires a driver
4      that instantiates the design, toggles its clock, and interacts with its ports.
5
6      The following driver may be used as an example for a design with a single clock
7      driving rising edge triggered flip-flops:
8
9      #include "top.cc"
10
11      int main() {
12          cxxrtl_design::p_top top;
13          top.step();
14          while (1) {
15              /* user logic */
16              top.p_clk.set(false);
17              top.step();
18              top.p_clk.set(true);
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

19         top.step();
20     }
21 }
22

```

23 Note that CXXRTL simulations, just like the hardware they are simulating, are  
24 subject to race conditions. If, in the example above, the user logic would run  
25 simultaneously with the rising edge of the clock, the design would malfunction.

26  
27 This backend supports replacing parts of the design with black boxes implemented  
28 in C++. If a module marked as a CXXRTL black box, its implementation is ignored,  
29 and the generated code consists only of an interface and a factory function.  
30 The driver must implement the factory function that creates an implementation of  
31 the black box, taking into account the parameters it is instantiated with.

32  
33 For example, the following Verilog code defines a CXXRTL black box interface for  
34 a synchronous debug sink:

```

35
36     (* cxxrtl_blackbox *)
37     module debug(...);
38         (* cxxrtl_edge = "p" *) input clk;
39         input en;
40         input [7:0] i_data;
41         (* cxxrtl_sync *) output [7:0] o_data;
42     endmodule
43

```

44 For this HDL interface, this backend will generate the following C++ interface:

```

45
46     struct bb_p_debug : public module {
47         value<1> p_clk;
48         bool posedge_p_clk() const { /* ... */ }
49         value<1> p_en;
50         value<8> p_i_data;
51         wire<8> p_o_data;
52
53         bool eval() override;
54         bool commit() override;
55
56         static std::unique_ptr<bb_p_debug>
57             create(std::string name, metadata_map parameters, metadata_map attributes);
58     };
59

```

60 The `create' function must be implemented by the driver. For example, it could  
61 always provide an implementation logging the values to standard error stream:

```

62
63     namespace cxxrtl_design {
64
65     struct stderr_debug : public bb_p_debug {
66         bool eval() override {
67             if (posedge_p_clk() && p_en)
68                 fprintf(stderr, "debug: %02x\n", p_i_data.data[0]);
69             p_o_data.next = p_i_data;
70             return bb_p_debug::eval();
71         }
72     };
73

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

73
74     std::unique_ptr<bb_p_debug>
75     bb_p_debug::create(std::string name, cxxrtl::metadata_map parameters,
76                       cxxrtl::metadata_map attributes) {
77         return std::make_unique<stderr_debug>();
78     }
79
80 }
81

```

For complex applications of black boxes, it is possible to parameterize their port widths. For example, the following Verilog code defines a CXXRTL black box interface for a configurable width debug sink:

```

82
83
84
85
86     (* cxxrtl_blackbox, cxxrtl_template = "WIDTH" *)
87     module debug(...);
88         parameter WIDTH = 8;
89         (* cxxrtl_edge = "p" *) input clk;
90         input en;
91         (* cxxrtl_width = "WIDTH" *) input [WIDTH - 1:0] i_data;
92         (* cxxrtl_width = "WIDTH" *) output [WIDTH - 1:0] o_data;
93     endmodule
94

```

For this parametric HDL interface, this backend will generate the following C++ interface (only the differences are shown):

```

95
96
97
98     template<size_t WIDTH>
99     struct bb_p_debug : public module {
100         // ...
101         value<WIDTH> p_i_data;
102         wire<WIDTH> p_o_data;
103         // ...
104         static std::unique_ptr<bb_p_debug<WIDTH>>
105         create(std::string name, metadata_map parameters, metadata_map attributes);
106     };
107

```

The `create` function must be implemented by the driver, specialized for every possible combination of template parameters. (Specialization is necessary to enable separate compilation of generated code and black box implementations.)

```

111
112     template<size_t SIZE>
113     struct stderr_debug : public bb_p_debug<SIZE> {
114         // ...
115     };
116
117     template<>
118     std::unique_ptr<bb_p_debug<8>>
119     bb_p_debug<8>::create(std::string name, cxxrtl::metadata_map parameters,
120                          cxxrtl::metadata_map attributes) {
121         return std::make_unique<stderr_debug<8>>();
122     }
123

```

The following attributes are recognized by this backend:

```

124
125
126     cxxrtl_blackbox

```



## APPENDIX C. COMMAND REFERENCE MANUAL

only valid on modules. if specified, the module contents are ignored, and the generated code includes only the module interface and a factory function, which will be called to instantiate the module.

**cxxrtl\_edge**  
only valid on inputs of black boxes. must be one of "p", "n", "a". if specified on signal ``clk``, the generated code includes edge detectors ``posedge_p_clk()`` (if "p"), ``negedge_p_clk()`` (if "n"), or both (if "a"), simplifying implementation of clocked black boxes.

**cxxrtl\_template**  
only valid on black boxes. must contain a space separated sequence of identifiers that have a corresponding black box parameters. for each of them, the generated code includes a ``size_t`` template parameter.

**cxxrtl\_width**  
only valid on ports of black boxes. must be a constant expression, which is directly inserted into generated code.

**cxxrtl\_comb, cxxrtl\_sync**  
only valid on outputs of black boxes. if specified, indicates that every bit of the output port is driven, correspondingly, by combinatorial or synchronous logic. this knowledge is used for scheduling optimizations. if neither is specified, the output will be pessimistically treated as driven by both combinatorial and synchronous logic.

The following options are supported by this backend:

- print-wire-types, -print-debug-wire-types**  
enable additional debug logging, for pass developers.
- header**  
generate separate interface (.h) and implementation (.cc) files. if specified, the backend must be called with a filename, and filename of the interface is derived from filename of the implementation. otherwise, interface and implementation are generated together.
- namespace <ns-name>**  
place the generated code into namespace <ns-name>. if not specified, "cxxrtl\_design" is used.
- nohierarchy**  
use design hierarchy as-is. in most designs, a top module should be present as it is exposed through the C API and has unbuffered outputs for improved performance; it will be determined automatically if absent.
- noflatten**  
don't flatten the design. fully flattened designs can evaluate within one delta cycle if they have no combinatorial feedback. note that the debug interface and waveform dumps use full hierarchical names for all wires even in flattened designs.
- noprocs**  
don't convert processes to netlists. in most designs, converting

## APPENDIX C. COMMAND REFERENCE MANUAL

processes significantly improves evaluation performance at the cost of slight increase in compilation time.

**-O <level>**  
set the optimization level. the default is -O6. higher optimization levels dramatically decrease compile and run time, and highest level possible for a design should be used.

**-O0**  
no optimization.

**-O1**  
unbuffer internal wires if possible.

**-O2**  
like -O1, and localize internal wires if possible.

**-O3**  
like -O2, and inline internal wires if possible.

**-O4**  
like -O3, and unbuffer public wires not marked (\*keep\*) if possible.

**-O5**  
like -O4, and localize public wires not marked (\*keep\*) if possible.

**-O6**  
like -O5, and inline public wires not marked (\*keep\*) if possible.

**-g <level>**  
set the debug level. the default is -g4. higher debug levels provide more visibility and generate more code, but do not pessimize evaluation.

**-g0**  
no debug information. the C API is disabled.

**-g1**  
include bare minimum of debug information necessary to access all design state. the C API is enabled.

**-g2**  
like -g1, but include debug information for all public wires that are directly accessible through the C++ interface.

**-g3**  
like -g2, and include debug information for public wires that are tied to a constant or another public wire.

**-g4**  
like -g3, and compute debug information on demand for all public wires that were optimized out.

**C.228 write\_edif – write design to EDIF netlist file**

```

1  write_edif [options] [filename]
2
3  Write the current design to an EDIF netlist file.
4
5  -top top_module
6      set the specified module as design top module
7
8  -nogndvcc
9      do not create "GND" and "VCC" cells. (this will produce an error
10     if the design contains constant nets. use "hilomap" to map to custom
11     constant drivers first)
12
13  -gndvccy
14     create "GND" and "VCC" cells with "Y" outputs. (the default is "G"
15     for "GND" and "P" for "VCC".)
16
17  -attrprop
18     create EDIF properties for cell attributes
19
20  -keep
21     create extra KEEP nets by allowing a cell to drive multiple nets.
22
23  -pvector {par|bra|ang}
24     sets the delimiting character for module port rename clauses to
25     parentheses, square brackets, or angle brackets.
26
27  Unfortunately there are different "flavors" of the EDIF file format. This
28  command generates EDIF files for the Xilinx place&route tools. It might be
29  necessary to make small modifications to this command when a different tool
30  is targeted.

```

**C.229 write\_file – write a text to a file**

```

1  write_file [options] output_file [input_file]
2
3  Write the text from the input file to the output file.
4
5  -a
6      Append to output file (instead of overwriting)
7
8
9  Inside a script the input file can also can a here-document:
10
11  write_file hello.txt <<EOT
12  Hello World!
13  EOT

```

**C.230 write\_firrtl – write design to a FIRRTL file**

```

1  write_firrtl [options] [filename]
2
3  Write a FIRRTL netlist of the current design.
4  The following commands are executed by this command:
5      pmuxtree
6      bmuxmap
7      demuxmap

```

**C.231 write\_ilang – (deprecated) alias of write\_rtlil**

```

1  See `help write_rtlil`.

```

**C.232 write\_intersynth – write design to InterSynth netlist file**

```

1  write_intersynth [options] [filename]
2
3  Write the current design to an 'intersynth' netlist file. InterSynth is
4  a tool for Coarse-Grain Example-Driven Interconnect Synthesis.
5
6  -notypes
7      do not generate celltypes and conntypes commands. i.e. just output
8      the netlists. this is used for postsilicon synthesis.
9
10 -lib <verilog_or_rtlil_file>
11     Use the specified library file for determining whether cell ports are
12     inputs or outputs. This option can be used multiple times to specify
13     more than one library.
14
15 -selected
16     only write selected modules. modules must be selected entirely or
17     not at all.
18
19 http://bygone.clairexen.net/intersynth/

```

**C.233 write\_jny – generate design metadata**

```

1  jny [options] [selection]
2
3  -no-connections
4      Don't include connection information in the netlist output.
5
6  -no-attributes
7      Don't include attributed information in the netlist output.
8
9  -no-properties

```

```

10      Don't include property information in the netlist output.
11
12 Write a JSON metadata for the current design

```

## C.234 write\_json – write design to a JSON file

```

1  write_json [options] [filename]
2
3  Write a JSON netlist of the current design.
4
5  -aig
6      include AIG models for the different gate types
7
8  -compat-int
9      emit 32-bit or smaller fully-defined parameter values directly
10     as JSON numbers (for compatibility with old parsers)
11
12
13 The general syntax of the JSON output created by this command is as follows:
14
15 {
16     "creator": "Yosys <version info>",
17     "modules": {
18         <module_name>: {
19             "attributes": {
20                 <attribute_name>: <attribute_value>,
21                 ...
22             },
23             "parameter_default_values": {
24                 <parameter_name>: <parameter_value>,
25                 ...
26             },
27             "ports": {
28                 <port_name>: <port_details>,
29                 ...
30             },
31             "cells": {
32                 <cell_name>: <cell_details>,
33                 ...
34             },
35             "memories": {
36                 <memory_name>: <memory_details>,
37                 ...
38             },
39             "netnames": {
40                 <net_name>: <net_details>,
41                 ...
42             }
43         }
44     },
45     "models": {
46         ...

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

47     },
48     }

```

50 Where <port\_details> is:

```

51 {
52     "direction": <"input" | "output" | "inout">,
53     "bits": <bit_vector>
54     "offset": <the lowest bit index in use, if non-0>
55     "upto": <1 if the port bit indexing is MSB-first>
56     "signed": <1 if the port is signed>
57 }
58
59

```

60 The "offset" and "upto" fields are skipped if their value would be 0. They don't affect connection semantics

```

61 {
62     "hide_name": <1 | 0>,
63     "type": <cell_type>,
64     "model": <AIG model name, if -aig option used>,
65     "parameters": {
66         <parameter_name>: <parameter_value>,
67         ...
68     },
69     "attributes": {
70         <attribute_name>: <attribute_value>,
71         ...
72     },
73     "port_directions": {
74         <port_name>: <"input" | "output" | "inout">,
75         ...
76     },
77     "connections": {
78         <port_name>: <bit_vector>,
79         ...
80     },
81 }
82
83

```

84 And <memory\_details> is:

```

85 {
86     "hide_name": <1 | 0>,
87     "attributes": {
88         <attribute_name>: <attribute_value>,
89         ...
90     },
91     "width": <memory width>
92     "start_offset": <the lowest valid memory address>
93     "size": <memory size>
94 }
95
96

```

97 And <net\_details> is:

```

98 {
99     "hide_name": <1 | 0>,
100

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

    "bits": <bit_vector>
    "offset": <the lowest bit index in use, if non-0>
    "upto": <1 if the port bit indexing is MSB-first>
    "signed": <1 if the port is signed>
}

```

The "hide\_name" fields are set to 1 when the name of this cell or net is automatically created and is likely not of interest for a regular user.

The "port\_directions" section is only included for cells for which the interface is known.

Module and cell ports and nets can be single bit wide or vectors of multiple bits. Each individual signal bit is assigned a unique integer. The <bit\_vector> values referenced above are vectors of this integers. Signal bits that are connected to a constant driver are denoted as string "0", "1", "x", or "z" instead of a number.

Bit vectors (including integers) are written as string holding the binary representation of the value. String

For example the following Verilog code:

```

module test(input x, y);
    (* keep *) foo #(.P(42), .Q(1337))
        foo_inst (.A({x, y}), .B({y, x}), .C({4'd10, {4{x}}}}));
endmodule

```

Translates to the following JSON output:

```

{
  "creator": "Yosys 0.9+2406 (git sha1 fb1168d8, clang 9.0.1 -fPIC -Os)",
  "modules": {
    "test": {
      "attributes": {
        "cells_not_processed": "00000000000000000000000000000001",
        "src": "test.v:1.1-4.10"
      },
      "ports": {
        "x": {
          "direction": "input",
          "bits": [ 2 ]
        },
        "y": {
          "direction": "input",
          "bits": [ 3 ]
        }
      },
      "cells": {
        "foo_inst": {
          "hide_name": 0,
          "type": "foo",
          "parameters": {
            "P": "0000000000000000000000000000000101010",
            "Q": "000000000000000000000000000000010100111001"
          }
        }
      }
    }
  }
}

```

```

155     },
156     "attributes": {
157         "keep": "00000000000000000000000000000001",
158         "module_not_derived": "00000000000000000000000000000001",
159         "src": "test.v:3.1-3.55"
160     },
161     "connections": {
162         "A": [ 3, 2 ],
163         "B": [ 2, 3 ],
164         "C": [ 2, 2, 2, 2, "0", "1", "0", "1" ]
165     }
166 }
167 },
168 "netnames": {
169     "x": {
170         "hide_name": 0,
171         "bits": [ 2 ],
172         "attributes": {
173             "src": "test.v:1.19-1.20"
174         }
175     },
176     "y": {
177         "hide_name": 0,
178         "bits": [ 3 ],
179         "attributes": {
180             "src": "test.v:1.22-1.23"
181         }
182     }
183 }
184 }
185 }
186 }
187

```

The models are given as And-Inverter-Graphs (AIGs) in the following form:

```

190 "models": {
191     <model_name>: [
192         /* 0 */ [ <node-spec> ],
193         /* 1 */ [ <node-spec> ],
194         /* 2 */ [ <node-spec> ],
195         ...
196     ],
197     ...
198 },
199

```

The following node-types may be used:

```

202 [ "port", <portname>, <bitindex>, <out-list> ]
203     - the value of the specified input port bit
204
205 [ "nport", <portname>, <bitindex>, <out-list> ]
206     - the inverted value of the specified input port bit
207
208 [ "and", <node-index>, <node-index>, <out-list> ]

```



```

209     - the ANDed value of the specified nodes
210
211     [ "nand", <node-index>, <node-index>, <out-list> ]
212     - the inverted ANDed value of the specified nodes
213
214     [ "true", <out-list> ]
215     - the constant value 1
216
217     [ "false", <out-list> ]
218     - the constant value 0
219

```

220 All nodes appear in topological order. I.e. only nodes with smaller indices  
 221 are referenced by "and" and "nand" nodes.

222  
 223 The optional <out-list> at the end of a node specification is a list of  
 224 output portname and bitindex pairs, specifying the outputs driven by this node.

225  
 226 For example, the following is the model for a 3-input 3-output \$reduce\_and cell  
 227 inferred by the following code:

```

228
229     module test(input [2:0] in, output [2:0] out);
230         assign in = &out;
231     endmodule
232
233     "$reduce_and:3U:3": [
234         /* 0 */ [ "port", "A", 0 ],
235         /* 1 */ [ "port", "A", 1 ],
236         /* 2 */ [ "and", 0, 1 ],
237         /* 3 */ [ "port", "A", 2 ],
238         /* 4 */ [ "and", 2, 3, "Y", 0 ],
239         /* 5 */ [ "false", "Y", 1, "Y", 2 ]
240     ]
241

```

242 Future version of Yosys might add support for additional fields in the JSON  
 243 format. A program processing this format must ignore all unknown fields.

## C.235 write\_rtlil – write design to RTLIL file

```

1     write_rtlil [filename]
2
3     Write the current design to an RTLIL file. (RTLIL is a text representation
4     of a design in yosys's internal format.)
5
6     -selected
7         only write selected parts of the design.

```

## C.236 write\_simplec – convert design to simple C code

```

1     write_simplec [options] [filename]
2

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
3 Write simple C code for simulating the design. The C code written can be used to
4 simulate the design in a C environment, but the purpose of this command is to
5 generate code that works well with C-based formal verification.
6
7     -verbose
8         this will print the recursive walk used to export the modules.
9
10    -i8, -i16, -i32, -i64
11        set the maximum integer bit width to use in the generated code.
12
13 THIS COMMAND IS UNDER CONSTRUCTION
```

### C.237 write\_smt2 – write design to SMT-LIBv2 file

```
1     write_smt2 [options] [filename]
2
3 Write a SMT-LIBv2 [1] description of the current design. For a module with name
4 '<mod>' this will declare the sort '<mod>_s' (state of the module) and will
5 define and declare functions operating on that state.
6
7 The following SMT2 functions are generated for a module with name '<mod>'.
8 Some declarations/definitions are printed with a special comment. A prover
9 using the SMT2 files can use those comments to collect all relevant metadata
10 about the design.
11
12     ; yosys-smt2-module <mod>
13     (declare-sort |<mod>_s| 0)
14         The sort representing a state of module <mod>.
15
16     (define-fun |<mod>_h| ((state |<mod>_s|)) Bool (...))
17         This function must be asserted for each state to establish the
18         design hierarchy.
19
20     ; yosys-smt2-input <wirename> <width>
21     ; yosys-smt2-output <wirename> <width>
22     ; yosys-smt2-register <wirename> <width>
23     ; yosys-smt2-wire <wirename> <width>
24     (define-fun |<mod>_n <wirename>| (|<mod>_s|) (_ BitVec <width>))
25     (define-fun |<mod>_n <wirename>| (|<mod>_s|) Bool)
26         For each port, register, and wire with the 'keep' attribute set an
27         accessor function is generated. Single-bit wires are returned as Bool,
28         multi-bit wires as BitVec.
29
30     ; yosys-smt2-cell <submod> <instancename>
31     (declare-fun |<mod>_h <instancename>| (|<mod>_s|) |<submod>_s|)
32         There is a function like that for each hierarchical instance. It
33         returns the sort that represents the state of the sub-module that
34         implements the instance.
35
36     (declare-fun |<mod>_is| (|<mod>_s|) Bool)
37         This function must be asserted 'true' for initial states, and 'false'
38         otherwise.
```

## APPENDIX C. COMMAND REFERENCE MANUAL

```

39
40 (define-fun |<mod>_i| ((state |<mod>_s|)) Bool (...))
41   This function must be asserted 'true' for initial states. For
42   non-initial states it must be left unconstrained.
43
44 (define-fun |<mod>_t| ((state |<mod>_s|) (next_state |<mod>_s|)) Bool (...))
45   This function evaluates to 'true' if the states 'state' and
46   'next_state' form a valid state transition.
47
48 (define-fun |<mod>_a| ((state |<mod>_s|)) Bool (...))
49   This function evaluates to 'true' if all assertions hold in the state.
50
51 (define-fun |<mod>_u| ((state |<mod>_s|)) Bool (...))
52   This function evaluates to 'true' if all assumptions hold in the state.
53
54 ; yosys-smt2-assert <id> <filename:linenum>
55 (define-fun |<mod>_a <id>| ((state |<mod>_s|)) Bool (...))
56   Each $assert cell is converted into one of this functions. The function
57   evaluates to 'true' if the assert statement holds in the state.
58
59 ; yosys-smt2-assume <id> <filename:linenum>
60 (define-fun |<mod>_u <id>| ((state |<mod>_s|)) Bool (...))
61   Each $assume cell is converted into one of this functions. The function
62   evaluates to 'true' if the assume statement holds in the state.
63
64 ; yosys-smt2-cover <id> <filename:linenum>
65 (define-fun |<mod>_c <id>| ((state |<mod>_s|)) Bool (...))
66   Each $cover cell is converted into one of this functions. The function
67   evaluates to 'true' if the cover statement is activated in the state.
68
69 Options:
70
71 -verbose
72   this will print the recursive walk used to export the modules.
73
74 -stbv
75   Use a BitVec sort to represent a state instead of an uninterpreted
76   sort. As a side-effect this will prevent use of arrays to model
77   memories.
78
79 -stdt
80   Use SMT-LIB 2.6 style datatypes to represent a state instead of an
81   uninterpreted sort.
82
83 -nobv
84   disable support for BitVec (FixedSizeBitVectors theory). without this
85   option multi-bit wires are represented using the BitVec sort and
86   support for coarse grain cells (incl. arithmetic) is enabled.
87
88 -nomem
89   disable support for memories (via ArraysEx theory). this option is
90   implied by -nobv. only $mem cells without merged registers in
91   read ports are supported. call "memory" with -nordff to make sure
92   that no registers are merged into $mem read ports. '<mod>_m' functions

```

## APPENDIX C. COMMAND REFERENCE MANUAL

```
93     will be generated for accessing the arrays that are used to represent
94     memories.
95
96     -wires
97         create '<mod>_n' functions for all public wires. by default only ports,
98         registers, and wires with the 'keep' attribute are exported.
99
100    -tpl <template_file>
101        use the given template file. the line containing only the token '%%'
102        is replaced with the regular output of this command.
103
104    -solver-option <option> <value>
105        emit a `; yosys-smt2-solver-option` directive for yosys-smtbmc to write
106        the given option as a `(set-option ...)` command in the SMT-LIBv2.
107
108    [1] For more information on SMT-LIBv2 visit http://smt-lib.org/ or read David
109    R. Cok's tutorial: https://smtlib.github.io/jSMTLIB/SMTLIBTutorial.pdf
110
111    -----
112
113    Example:
114
115    Consider the following module (test.v). We want to prove that the output can
116    never transition from a non-zero value to a zero value.
117
118        module test(input clk, output reg [3:0] y);
119            always @(posedge clk)
120                y <= (y << 1) | ^y;
121        endmodule
122
123    For this proof we create the following template (test.tpl).
124
125        ; we need QF_UFBV for this proof
126        (set-logic QF_UFBV)
127
128        ; insert the auto-generated code here
129        %%
130
131        ; declare two state variables s1 and s2
132        (declare-fun s1 () test_s)
133        (declare-fun s2 () test_s)
134
135        ; state s2 is the successor of state s1
136        (assert (test_t s1 s2))
137
138        ; we are looking for a model with y non-zero in s1
139        (assert (distinct (|test_n y| s1) #b0000))
140
141        ; we are looking for a model with y zero in s2
142        (assert (= (|test_n y| s2) #b0000))
143
144        ; is there such a model?
145        (check-sat)
146
```

```

147 The following yosys script will create a 'test.smt2' file for our proof:
148
149     read_verilog test.v
150     hierarchy -check; proc; opt; check -assert
151     write_smt2 -bv -tpl test.tpl test.smt2
152
153 Running 'cvc4 test.smt2' will print 'unsat' because y can never transition
154 from non-zero to zero in the test design.

```

### C.238 write\_smv – write design to SMV file

```

1     write_smv [options] [filename]
2
3 Write an SMV description of the current design.
4
5     -verbose
6         this will print the recursive walk used to export the modules.
7
8     -tpl <template_file>
9         use the given template file. the line containing only the token '%%'
10        is replaced with the regular output of this command.
11
12 THIS COMMAND IS UNDER CONSTRUCTION

```

### C.239 write\_spice – write design to SPICE netlist file

```

1     write_spice [options] [filename]
2
3 Write the current design to an SPICE netlist file.
4
5     -big_endian
6         generate multi-bit ports in MSB first order
7         (default is LSB first)
8
9     -neg net_name
10        set the net name for constant 0 (default: Vss)
11
12     -pos net_name
13        set the net name for constant 1 (default: Vdd)
14
15     -buf DC|subckt_name
16        set the name for jumper element (default: DC)
17        (used to connect different nets)
18
19     -nc_prefix
20        prefix for not-connected nets (default: _NC)
21
22     -inames
23        include names of internal ($-prefixed) nets in outputs
24        (default is to use net numbers instead)

```

```

25
26     -top top_module
27         set the specified module as design top module

```

## C.240 write\_table – write design as connectivity table

```

1     write_table [options] [filename]
2
3 Write the current design as connectivity table. The output is a tab-separated
4 ASCII table with the following columns:
5
6     module name
7     cell name
8     cell type
9     cell port
10    direction
11    signal
12
13 module inputs and outputs are output using cell type and port '-' and with
14 'pi' (primary input) or 'po' (primary output) or 'pio' as direction.

```

## C.241 write\_verilog – write design to Verilog file

```

1     write_verilog [options] [filename]
2
3 Write the current design to a Verilog file.
4
5     -sv
6         with this option, SystemVerilog constructs like always_comb are used
7
8     -norename
9         without this option all internal object names (the ones with a dollar
10        instead of a backslash prefix) are changed to short names in the
11        format '_<number>_'.
12
13     -renameprefix <prefix>
14         insert this prefix in front of auto-generated instance names
15
16     -noattr
17         with this option no attributes are included in the output
18
19     -attr2comment
20         with this option attributes are included as comments in the output
21
22     -noexpr
23         without this option all internal cells are converted to Verilog
24         expressions.
25
26     -siminit
27         add initial statements with hierarchical refs to initialize FFs when

```

## APPENDIX C. COMMAND REFERENCE MANUAL

28       in -noexpr mode.

29

30       -nodec

31       32-bit constant values are by default dumped as decimal numbers,

32       not bit pattern. This option deactivates this feature and instead

33       will write out all constants in binary.

34

35       -decimal

36       dump 32-bit constants in decimal and without size and radix

37

38       -nohex

39       constant values that are compatible with hex output are usually

40       dumped as hex values. This option deactivates this feature and

41       instead will write out all constants in binary.

42

43       -nostr

44       Parameters and attributes that are specified as strings in the

45       original input will be output as strings by this back-end. This

46       deactivates this feature and instead will write string constants

47       as binary numbers.

48

49       -simple-lhs

50       Connection assignments with simple left hand side without concatenations.

51

52       -extmem

53       instead of initializing memories using assignments to individual

54       elements, use the '\$readmemh' function to read initialization data

55       from a file. This data is written to a file named by appending

56       a sequential index to the Verilog filename and replacing the extension

57       with '.mem', e.g. 'write\_verilog -extmem foo.v' writes 'foo-1.mem',

58       'foo-2.mem' and so on.

59

60       -defparam

61       use 'defparam' statements instead of the Verilog-2001 syntax for

62       cell parameters.

63

64       -blackboxes

65       usually modules with the 'blackbox' attribute are ignored. with

66       this option set only the modules with the 'blackbox' attribute

67       are written to the output file.

68

69       -selected

70       only write selected modules. modules must be selected entirely or

71       not at all.

72

73       -v

74       verbose output (print new names of all renamed wires and cells)

75

76       Note that RTLIL processes can't always be mapped directly to Verilog

77       always blocks. This frontend should only be used to export an RTLIL

78       netlist, i.e. after the "proc" pass has been used to convert all

79       processes to logic networks and registers. A warning is generated when

80       this command is called on a design with RTLIL processes.

**C.242 write\_xaiger – write design to XAIGER file**

```

1  write_xaiger [options] [filename]
2
3  Write the top module (according to the (* top *) attribute or if only one module
4  is currently selected) to an XAIGER file. Any non $_NOT_, $_AND_, (optionally
5  $_DFF_N_, $_DFF_P_), or non (* abc9_box *) cells will be converted into psuedo-
6  inputs and pseudo-outputs. Whitebox contents will be taken from the equivalent
7  module in the '$abc9_holes' design, if it exists.
8
9  -ascii
10     write ASCII version of AIGER format
11
12  -map <filename>
13     write an extra file with port and box symbols
14
15  -dff
16     write $_DFF_[NP]_ cells

```

**C.243 xilinx\_dffopt – Xilinx: optimize FF control signal usage**

```

1  xilinx_dffopt [options] [selection]
2
3  Converts hardware clock enable and set/reset signals on FFs to emulation
4  using LUTs, if doing so would improve area. Operates on post-techmap Xilinx
5  cells (LUT*, FD*).
6
7  -lut4
8     Assume a LUT4-based device (instead of a LUT6-based device).

```

**C.244 xilinx\_dsp – Xilinx: pack resources into DSPs**

```

1  xilinx_dsp [options] [selection]
2
3  Pack input registers (A2, A1, B2, B1, C, D, AD; with optional enable/reset),
4  pipeline registers (M; with optional enable/reset), output registers (P; with
5  optional enable/reset), pre-adder and/or post-adder into Xilinx DSP resources.
6
7  Multiply-accumulate operations using the post-adder with feedback on the 'C'
8  input will be folded into the DSP. In this scenario only, the 'C' input can be
9  used to override the current accumulation result with a new value, which will
10 be added to the multiplier result to form the next accumulation result.
11
12 Use of the dedicated 'PCOUT' -> 'PCIN' cascade path is detected for 'P' -> 'C'
13 connections (optionally, where 'P' is right-shifted by 17-bits and used as an
14 input to the post-adder -- a pattern common for summing partial products to
15 implement wide multipliers). Limited support also exists for similar cascading
16 for A and B using '[AB]COUT' -> '[AB]CIN'. Currently, cascade chains are limited
17 to a maximum length of 20 cells, corresponding to the smallest Xilinx 7 Series

```



## APPENDIX C. COMMAND REFERENCE MANUAL

```
18 device.
19
20 This pass is a no-op if the scratchpad variable 'xilinx_dsp.multonly' is set
21 to 1.
22
23
24 Experimental feature: addition/subtractions less than 12 or 24 bits with the
25 '(* use_dsp="simd" *)' attribute attached to the output wire or attached to
26 the add/subtract operator will cause those operations to be implemented using
27 the 'SIMD' feature of DSPs.
28
29 Experimental feature: the presence of a `$ge' cell attached to the registered
30 P output implementing the operation "(P >= <power-of-2>)" will be transformed
31 into using the DSP48E1's pattern detector feature for overflow detection.
32
33     -family {xcup|xcu|xc7|xc6v|xc5v|xc4v|xc6s|xc3sda}
34         select the family to target
35     default: xc7
```

### C.245 xilinx\_srl – Xilinx shift register extraction

```
1     xilinx_srl [options] [selection]
2
3 This pass converts chains of built-in flops (bit-level: $_DFF_[NP]_, $_DFFE_*
4 and word-level: $dff, $dfffe) as well as Xilinx flops (FDRE, FDRE_1) into a
5 $_XILINX_SHREG cell. Chains must be of the same cell type, clock, clock polarity,
6 enable, and enable polarity (where relevant).
7 Flops with resets cannot be mapped to Xilinx devices and will not be inferred.
8     -minlen N
9         min length of shift register (default = 3)
10
11     -fixed
12         infer fixed-length shift registers.
13
14     -variable
15         infer variable-length shift registers (i.e. fixed-length shifts where
16         each element also fans-out to a $shiftx cell).
```

### C.246 zinit – add inverters so all FF are zero-initialized

```
1     zinit [options] [selection]
2
3 Add inverters as needed to make all FFs zero-initialized.
4
5     -all
6         also add zero initialization to uninitialized FFs
```

## Appendix D

# RTLIL Text Representation

This appendix documents the text representation of RTLIL in extended Backus-Naur form (EBNF).

The grammar is not meant to represent semantic limitations. That is, the grammar is “permissive”, and later stages of processing perform more rigorous checks.

The grammar is also not meant to represent the exact grammar used in the RTLIL frontend, since that grammar is specific to processing by `lex` and `yacc`, is even more permissive, and is somewhat less understandable than simple EBNF notation.

Finally, note that all statements (rules ending in `-stmt`) terminate in an end-of-line. Because of this, a statement cannot be broken into multiple lines.

### D.1 Lexical elements

#### D.1.1 Characters

An RTLIL file is a stream of bytes. Strictly speaking, a “character” in an RTLIL file is a single byte. The lexer treats multi-byte encoded characters as consecutive single-byte characters. While other encodings *may* work, UTF-8 is known to be safe to use. Byte order marks at the beginning of the file will cause an error.

ASCII spaces (32) and tabs (9) separate lexer tokens.

A `nonws` character, used in identifiers, is any character whose encoding consists solely of bytes above ASCII space (32).

An `eol` is one or more consecutive ASCII newlines (10) and carriage returns (13).

#### D.1.2 Identifiers

There are two types of identifiers in RTLIL:

- Publically visible identifiers
- Auto-generated identifiers

$$\langle id \rangle ::= \langle public-id \rangle \mid \langle autogen-id \rangle$$

$$\langle public-id \rangle ::= \backslash \langle nonws \rangle +$$

$$\langle autogen-id \rangle ::= \$ \langle nonws \rangle +$$


---

### D.1.3 Values

A *value* consists of a width in bits and a bit representation, most significant bit first. Bits may be any of:

- 0: A logic zero value
- 1: A logic one value
- x: An unknown logic value (or don't care in case patterns)
- z: A high-impedance value (or don't care in case patterns)
- m: A marked bit (internal use only)
- -: A don't care value

An *integer* is simply a signed integer value in decimal format. **Warning:** Integer constants are limited to 32 bits. That is, they may only be in the range  $[-2147483648, 2147483648]$ . Integers outside this range will result in an error.

---


$$\langle value \rangle ::= \langle decimal-digit \rangle + \mid \langle binary-digit \rangle *$$

$$\langle decimal-digit \rangle ::= 0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7 \mid 8 \mid 9$$

$$\langle binary-digit \rangle ::= 0 \mid 1 \mid x \mid z \mid m \mid -$$

$$\langle integer \rangle ::= -? \langle decimal-digit \rangle +$$


---

### D.1.4 Strings

A string is a series of characters delimited by double-quote characters. Within a string, any character except ASCII NUL (0) may be used. In addition, certain escapes can be used:

- \n: A newline
- \t: A tab
- \ooo: A character specified as a one, two, or three digit octal value

All other characters may be escaped by a backslash, and become the following character. Thus:

- \\: A backslash
- \": A double-quote
- \r: An 'r' character

### D.1.5 Comments

A comment starts with a `#` character and proceeds to the end of the line. All comments are ignored.

## D.2 File

A file consists of an optional autoindex statement followed by zero or more modules.

---

```
<file> ::= <autoidx-stmt>? <module>*
```

---

### D.2.1 Autoindex statements

The autoindex statement sets the global autoindex value used by Yosys when it needs to generate a unique name, e.g. `$flatten$N`. The `N` part is filled with the value of the global autoindex value, which is subsequently incremented. This global has to be dumped into RTLIL, otherwise e.g. dumping and running a pass would have different properties than just running a pass on a warm design.

---

```
<autoidx-stmt> ::= autoidx <integer> <eol>
```

---

### D.2.2 Modules

Declares a module, with zero or more attributes, consisting of zero or more wires, memories, cells, processes, and connections.

---

```
<module> ::= <attr-stmt>* <module-stmt> <module-body> <module-end-stmt>

<module-stmt> ::= module <id> <eol>

<module-body> ::= (<param-stmt>
| <wire>
| <memory>
| <cell>
| <process>
| <conn-stmt>)*

<param-stmt> ::= parameter <id> <constant>? <eol>

<constant> ::= <value> | <integer> | <string>

<module-end-stmt> ::= end <eol>
```

---

### D.2.3 Attribute statements

Declares an attribute with the given identifier and value.

---

```
<attr-stmt> ::= attribute <id> <constant> <eol>
```

---

### D.2.4 Signal specifications

A signal is anything that can be applied to a cell port, i.e. a constant value, all bits or a selection of bits from a wire, or concatenations of those.

**Warning:** When an integer constant is a sigspec, it is always 32 bits wide, 2's complement. For example, a constant of  $-1$  is the same as  $32'11111111111111111111111111111111$ , while a constant of  $1$  is the same as  $32'1$ .

See Sec. 4.2.4 for an overview of signal specifications.

---

```

<sigspec> ::= <constant>
           | <wire-id>
           | <sigspec> [ <integer> (: <integer>)? ]
           | { <sigspec>* }

```

---

### D.2.5 Connections

Declares a connection between the given signals.

---

```

<conn-stmt> ::= connect <sigspec> <sigspec> <eol>

```

---

### D.2.6 Wires

Declares a wire, with zero or more attributes, with the given identifier and options in the enclosing module.

See Sec. 4.2.3 for an overview of wires.

---

```

<wire>      ::= <attr-stmt>* <wire-stmt>

<wire-stmt> ::= wire <wire-option>* <wire-id> <eol>

<wire-id>   ::= <id>

<wire-option> ::= width <integer>
                 | offset <integer>
                 | input <integer>
                 | output <integer>
                 | inout <integer>
                 | upto
                 | signed

```

---

### D.2.7 Memories

Declares a memory, with zero or more attributes, with the given identifier and options in the enclosing module.

See Sec. 4.2.6 for an overview of memory cells, and Sec. 5.1.5 for details about memory cell types.

---

```

<memory>          ::= <attr-stmt>* <memory-stmt>
<memory-stmt>     ::= memory <memory-option>* <id> <eol>
<memory-option>   ::= width <integer>
                   | size <integer>
                   | offset <integer>

```

---

### D.2.8 Cells

Declares a cell, with zero or more attributes, with the given identifier and type in the enclosing module.

Cells perform functions on input signals. See Chap. 5 for a detailed list of cell types.

---

```

<cell>             ::= <attr-stmt>* <cell-stmt> <cell-body-stmt>* <cell-end-stmt>
<cell-stmt>        ::= cell <cell-type> <cell-id> <eol>
<cell-id>          ::= <id>
<cell-type>        ::= <id>
<cell-body-stmt>   ::= parameter (signed | real)? <id> <constant> <eol>
                   | connect <id> <sigspec> <eol>
<cell-end-stmt>    ::= end <eol>

```

---

### D.2.9 Processes

Declares a process, with zero or more attributes, with the given identifier in the enclosing module. The body of a process consists of zero or more assignments, exactly one switch, and zero or more syncs.

See Sec. 4.2.5 for an overview of processes.

---

```

<process>          ::= <attr-stmt>* <proc-stmt> <process-body> <proc-end-stmt>
<proc-stmt>        ::= process <id> <eol>
<process-body>     ::= <assign-stmt>* <switch>? <assign-stmt>* <sync>*
<assign-stmt>      ::= assign <dest-sigspec> <src-sigspec> <eol>
<dest-sigspec>     ::= <sigspec>
<src-sigspec>      ::= <sigspec>
<proc-end-stmt>    ::= end <eol>

```

---

### D.2.10 Switches

Switches test a signal for equality against a list of cases. Each case specifies a comma-separated list of signals to check against. If there are no signals in the list, then the case is the default case. The body of a case consists of zero or more switches and assignments. Both switches and cases may have zero or more attributes.

---

$\langle \text{switch} \rangle$	$::= \langle \text{switch-stmt} \rangle \langle \text{case} \rangle^* \langle \text{switch-end-stmt} \rangle$
$\langle \text{switch-stmt} \rangle$	$::= \langle \text{attr-stmt} \rangle^* \text{switch} \langle \text{sigspec} \rangle \langle \text{eol} \rangle$
$\langle \text{case} \rangle$	$::= \langle \text{attr-stmt} \rangle^* \langle \text{case-stmt} \rangle \langle \text{case-body} \rangle$
$\langle \text{case-stmt} \rangle$	$::= \text{case} \langle \text{compare} \rangle? \langle \text{eol} \rangle$
$\langle \text{compare} \rangle$	$::= \langle \text{sigspec} \rangle (, \langle \text{sigspec} \rangle)^*$
$\langle \text{case-body} \rangle$	$::= (\langle \text{switch} \rangle \mid \langle \text{assign-stmt} \rangle)^*$
$\langle \text{switch-end-stmt} \rangle$	$::= \text{end} \langle \text{eol} \rangle$

---

### D.2.11 Syncs

Syncs update signals with other signals when an event happens. Such an event may be:

- An edge or level on a signal
- Global clock ticks
- Initialization
- Always

---

$\langle \text{sync} \rangle$	$::= \langle \text{sync-stmt} \rangle \langle \text{update-stmt} \rangle^*$
$\langle \text{sync-stmt} \rangle$	$::= \text{sync} \langle \text{sync-type} \rangle \langle \text{sigspec} \rangle \langle \text{eol} \rangle$ $\mid \text{sync global} \langle \text{eol} \rangle$ $\mid \text{sync init} \langle \text{eol} \rangle$ $\mid \text{sync always} \langle \text{eol} \rangle$
$\langle \text{sync-type} \rangle$	$::= \text{low} \mid \text{high} \mid \text{posedge} \mid \text{negedge} \mid \text{edge}$
$\langle \text{update-stmt} \rangle$	$::= \text{update} \langle \text{dest-sigspec} \rangle \langle \text{src-sigspec} \rangle \langle \text{eol} \rangle$

---

## Appendix E

# Application Notes

This appendix contains copies of the Yosys application notes.

- Yosys AppNote 010: Converting Verilog to BLIF .....Page [265](#)
- Yosys AppNote 011: Interactive Design Investigation ..... Page [268](#)
- Yosys AppNote 012: Converting Verilog to BTOR ..... Page [277](#)



## Yosys Application Note 010: Converting Verilog to BLIF

Claire Xenia Wolf  
November 2013

**Abstract**—Verilog-2005 is a powerful Hardware Description Language (HDL) that can be used to easily create complex designs from small HDL code. It is the preferred method of design entry for many designers<sup>1</sup>.

The Berkeley Logic Interchange Format (BLIF) [6] is a simple file format for exchanging sequential logic between programs. It is easy to generate and easy to parse and is therefore the preferred method of design entry for many authors of logic synthesis tools.

Yosys [1] is a feature-rich Open-Source Verilog synthesis tool that can be used to bridge the gap between the two file formats. It implements most of Verilog-2005 and thus can be used to import modern behavioral Verilog designs into BLIF-based design flows without dependencies on proprietary synthesis tools.

The scope of Yosys goes of course far beyond Verilog logic synthesis. But it is a useful and important feature and this Application Note will focus on this aspect of Yosys.

### I. INSTALLATION

Yosys written in C++ (using features from C++11) and is tested on modern Linux. It should compile fine on most UNIX systems with a C++11 compiler. The README file contains useful information on building Yosys and its prerequisites.

Yosys is a large and feature-rich program with a couple of dependencies. It is, however, possible to deactivate some of the dependencies in the Makefile, resulting in features in Yosys becoming unavailable. When problems with building Yosys are encountered, a user who is only interested in the features of Yosys that are discussed in this Application Note may deactivate TCL, Qt and MinisAT support in the Makefile and may opt against building `yosys-abc`.

This Application Note is based on GIT Rev. e216e0e from 2013-11-23 of Yosys [1]. The Verilog sources used for the examples are taken from `yosys-bigsim` [2], a collection of real-world designs used for regression testing Yosys.

### II. GETTING STARTED

We start our tour with the Navré processor from `yosys-bigsim`. The Navré processor [3] is an Open Source AVR clone. It is a single module (`softusb_navre`) in a single design file (`softusb_navre.v`). It also is using only features that map nicely to the BLIF format, for example it only uses synchronous resets.

Converting `softusb_navre.v` to `softusb_navre.blif` could not be easier:

```
1 | yosys -o softusb_navre.blif -S softusb_navre.v
```

Listing 1. Calling Yosys without script file

Behind the scenes Yosys is controlled by synthesis scripts that execute commands that operate on Yosys' internal state. For example, the `-o softusb_navre.blif` option just adds the command `write_blif softusb_navre.blif` to the end of the script. Likewise a file on the command line `- softusb_navre.v` in this case `-` adds the command `read_verilog softusb_navre.v` to the beginning of the synthesis script. In both cases the file type is detected from the file extension.

<sup>1</sup>The other half prefers VHDL, a very different but – of course – equally powerful language.

Finally the option `-s` instantiates a built-in default synthesis script. Instead of using `-s` one could also specify the synthesis commands for the script on the command line using the `-p` option, either using individual options for each command or by passing one big command string with a semicolon-separated list of commands. But in most cases it is more convenient to use an actual script file.

### III. USING A SYNTHESIS SCRIPT

With a script file we have better control over Yosys. The following script file replicates what the command from the last section did:

```
1 | read_verilog softusb_navre.v
2 | hierarchy
3 | proc; opt; memory; opt; techmap; opt
4 | write_blif softusb_navre.blif
```

Listing 2. `softusb_navre.js`

The first and last line obviously read the Verilog file and write the BLIF file.

The 2nd line checks the design hierarchy and instantiates parametrized versions of the modules in the design, if necessary. In the case of this simple design this is a no-op. However, as a general rule a synthesis script should always contain this command as first command after reading the input files.

The 3rd line does most of the actual work:

- The command `opt` is the Yosys' built-in optimizer. It can perform some simple optimizations such as const-folding and removing unconnected parts of the design. It is common practice to call `opt` after each major step in the synthesis procedure. In cases where too much optimization is not appreciated (for example when analyzing a design), it is recommended to call `clean` instead of `opt`.
- The command `proc` converts *processes* (Yosys' internal representation of Verilog `always-` and `initial-`blocks) to circuits of multiplexers and storage elements (various types of flip-flops).
- The command `memory` converts Yosys' internal representations of arrays and array accesses to multi-port block memories, and then maps this block memories to address decoders and flip-flops, unless the option `-nomap` is used, in which case the multi-port block memories stay in the design and can then be mapped to architecture-specific memory primitives using other commands.
- The command `techmap` turns a high-level circuit with coarse grain cells such as wide adders and multipliers to a fine-grain circuit of simple logic primitives and single-bit storage elements. The command does that by substituting the complex cells by circuits of simpler cells. It is possible to provide a custom set of rules for this process in the form of a Verilog source file, as we will see in the next section.

Now Yosys can be run with the filename of the synthesis script as argument:

```
1 | yosys softusb_navre.js
```

Listing 3. Calling Yosys with script file

Now that we are using a synthesis script we can easily modify how Yosys synthesizes the design. The first thing we should customize is the call to the `hierarchy` command:

Whenever it is known that there are no implicit blackboxes in the design, i.e. modules that are referenced but are not defined, the `hierarchy` command should be called with the `-check` option. This will then cause synthesis to fail when implicit blackboxes are found in the design.

The 2nd thing we can improve regarding the `hierarchy` command is that we can tell it the name of the top level module of the design hierarchy. It will then automatically remove all modules that are not referenced from this top level module.

For many designs it is also desired to optimize the encodings for the finite state machines (FSMs) in the design. The `fsm` command finds FSMs, extracts them, performs some basic optimizations and then generate a circuit from the extracted and optimized description. It would also be possible to tell the `fsm` command to leave the FSMs in their extracted form, so they can be further processed using custom commands. But in this case we don't want that.

So now we have the final synthesis script for generating a BLIF file for the Navré CPU:

```
1 read_verilog softusb_navre.v
2 hierarchy -check -top softusb_navre
3 proc; opt; memory; opt; fsm; opt; techmap; opt
4 write_blif softusb_navre.blif
```

Listing 4. `softusb_navre.ys` (improved)

#### IV. ADVANCED EXAMPLE: THE AMBER23 ARMV2A CPU

Our 2nd example is the Amber23 [4] ARMv2a CPU. Once again we base our example on the Verilog code that is included in `yosys-bigsim` [2].

The problem with this core is that it contains no dedicated reset logic. Instead the coding techniques shown in Listing 6 are used to define reset values for the global asynchronous reset in an FPGA implementation. This design can not be expressed in BLIF as it is. Instead we need to use a synthesis script that transforms this form to synchronous resets that can be expressed in BLIF.

```
1 read_verilog a23_alu.v
2 read_verilog a23_barrel_shift_fpga.v
3 read_verilog a23_barrel_shift.v
4 read_verilog a23_cache.v
5 read_verilog a23_coprocessor.v
6 read_verilog a23_core.v
7 read_verilog a23_decode.v
8 read_verilog a23_execute.v
9 read_verilog a23_fetch.v
10 read_verilog a23_multiply.v
11 read_verilog a23_ram_register_bank.v
12 read_verilog a23_register_bank.v
13 read_verilog a23_wishbone.v
14 read_verilog generic_sram_byte_en.v
15 read_verilog generic_sram_line_en.v
16 hierarchy -check -top a23_core
17 add -global_input globrst 1
18 proc -global_arst globrst
19 techmap -map adff2dff.v
20 opt; memory; opt; fsm; opt; techmap
21 write_blif amber23.blif
```

Listing 5. `amber23.ys`

```
1 reg [7:0] a = 13, b;
2 initial b = 37;
```

Listing 6. Implicit coding of global asynchronous resets

(Note that there is no problem if this coding techniques are used to model ROM, where the register is initialized using this syntax but is never updated otherwise.)

Listing 5 shows the synthesis script for the Amber23 core. In line 17 the `add` command is used to add a 1-bit wide global input signal with the name `globrst`. That means that an input with that name is added to each module in the design hierarchy and then all module instantiations are altered so that this new signal is connected throughout the whole design hierarchy.

In line 18 the `proc` command is called. But in this script the signal name `globrst` is passed to the command as a global reset signal for resetting the registers to their assigned initial values.

Finally in line 19 the `techmap` command is used to replace all instances of flip-flops with asynchronous resets with flip-flops with synchronous resets. The map file used for this is shown in Listing 7. Note how the `techmap_celltype` attribute is used in line 1 to tell the `techmap` command which cells to replace in the design, how the `_TECHMAP_FAIL_` wire in lines 15 and 16 (which evaluates to a constant value) determines if the parameter set is compatible with this replacement circuit, and how the `_TECHMAP_DO_` wire in line 13 provides a mini synthesis-script to be used to process this cell.

#### V. VERIFICATION OF THE AMBER23 CPU

The BLIF file for the Amber23 core, generated using Listings 5 and 7 and the version of the Amber23 RTL source that is bundled with `yosys-bigsim`, was verified using the test-bench from `yosys-bigsim`. It successfully executed the program shown in Listing 8 in the test-bench.

```
1 (* techmap_celltype = "$adff" *)
2 module adff2dff (CLK, ARST, D, Q);
3
4 parameter WIDTH = 1;
5 parameter CLK_POLARITY = 1;
6 parameter ARST_POLARITY = 1;
7 parameter ARST_VALUE = 0;
8
9 input CLK, ARST;
10 input [WIDTH-1:0] D;
11 output reg [WIDTH-1:0] Q;
12
13 wire [1023:0] _TECHMAP_DO_ = "proc";
14
15 wire _TECHMAP_FAIL_ =
16     !CLK_POLARITY || !ARST_POLARITY;
17
18 always @(posedge CLK)
19     if (ARST)
20         Q <= ARST_VALUE;
21     else
22         Q <= D;
23
24 endmodule
```

Listing 7. `adff2dff.v`

```

1  #include <stdint.h>
2  #include <stdbool.h>
3
4  #define BITMAP_SIZE 64
5  #define OUTPUT 0x10000000
6
7  static uint32_t bitmap[BITMAP_SIZE/32];
8
9  static void bitmap_set(uint32_t idx) { bitmap[idx/32] |= 1 << (idx % 32); }
10 static bool bitmap_get(uint32_t idx) { return (bitmap[idx/32] & (1 << (idx % 32))) != 0; }
11 static void output(uint32_t val) { *((volatile uint32_t*)OUTPUT) = val; }
12
13 int main() {
14     uint32_t i, j, k;
15     output(2);
16     for (i = 0; i < BITMAP_SIZE; i++) {
17         if (bitmap_get(i)) continue;
18         output(3+2*i);
19         for (j = 2*(3+2*i); j += 3+2*i) {
20             if (j%2 == 0) continue;
21             k = (j-3)/2;
22             if (k >= BITMAP_SIZE) break;
23             bitmap_set(k);
24         }
25     }
26     output(0);
27     return 0;
28 }

```

Listing 8. Test program for the Amber23 CPU (Sieve of Eratosthenes). Compiled using GCC 4.6.3 for ARM with `-Os -marm -march=armv2a -mno-thumb-interwork -ffreestanding`, linked with `--fix-v4bx` set and booted with a custom setup routine written in ARM assembler.

For simulation the BLIF file was converted back to Verilog using ABC [5]. So this test includes the successful transformation of the BLIF file into ABC’s internal format as well.

The only thing left to write about the simulation itself is that it probably was one of the most energy inefficient and time consuming ways of successfully calculating the first 31 primes the author has ever conducted.

## VI. LIMITATIONS

At the time of this writing Yosys does not support multi-dimensional memories, does not support writing to individual bits of array elements, does not support initialization of arrays with `$readmemb` and `$readmemh`, and has only limited support for tristate logic, to name just a few limitations.

That being said, Yosys can synthesize an overwhelming majority of real-world Verilog RTL code. The remaining cases can usually be modified to be compatible with Yosys quite easily.

The various designs in `yosys-bigsim` are a good place to look for examples of what is within the capabilities of Yosys.

## VII. CONCLUSION

Yosys is a feature-rich Verilog-2005 synthesis tool. It has many uses, but one is to provide an easy gateway from high-level Verilog code to low-level logic circuits.

The command line option `-s` can be used to quickly synthesize Verilog code to BLIF files without a hassle.

With custom synthesis scripts it becomes possible to easily perform high-level optimizations, such as re-encoding FSMs. In some extreme cases, such as the Amber23 ARMv2 CPU, the more advanced Yosys

features can be used to change a design to fit a certain need without actually touching the RTL code.

## REFERENCES

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## Yosys Application Note 011: Interactive Design Investigation

Claire Xenia Wolf  
Original Version December 2013

**Abstract**—Yosys [1] can be a great environment for building custom synthesis flows. It can also be an excellent tool for teaching and learning Verilog based RTL synthesis. In both applications it is of great importance to be able to analyze the designs it produces easily.

This Yosys application note covers the generation of circuit diagrams with the Yosys `show` command, the selection of interesting parts of the circuit using the `select` command, and briefly discusses advanced investigation commands for evaluating circuits and solving SAT problems.

### I. INSTALLATION AND PREREQUISITES

This Application Note is based on the Yosys [1] GIT Rev. 2b90ba1 from 2013-12-08. The README file covers how to install Yosys. The `show` command requires a working installation of GraphViz [2] and [3] for generating the actual circuit diagrams.

### II. OVERVIEW

This application note is structured as follows:

Sec. III introduces the `show` command and explains the symbols used in the circuit diagrams generated by it.

Sec. IV introduces additional commands used to navigate in the design, select portions of the design, and print additional information on the elements in the design that are not contained in the circuit diagrams.

Sec. V introduces commands to evaluate the design and solve SAT problems within the design.

Sec. VI concludes the document and summarizes the key points.

### III. INTRODUCTION TO THE `show` COMMAND

The `show` command generates a circuit diagram for the design in its current state. Various options can be used to change the appearance of the circuit diagram, set the name and format for the output file, and so forth. When called without any special options, it saves the circuit diagram in a temporary file and launches `xdot` to display the diagram. Subsequent calls to `show` re-use the `xdot` instance (if still running).

```

1  $ cat example.y
2  read_verilog example.v
3  show -pause
4  proc
5  show -pause
6  opt
7  show -pause
8
9  $ cat example.v
10 module example(input clk, a, b, c,
11                output reg [1:0] y);
12     always @(posedge clk)
13         if (c)
14             y <= c ? a + b : 2'd0;
15 endmodule

```

Figure 1. Yosys script with `show` commands and example design

#### A. A simple circuit

Fig. 1 shows a simple synthesis script and a Verilog file that demonstrate the usage of `show` in a simple setting. Note that `show` is called with the `-pause` option, that halts execution of the Yosys script until the user presses the Enter key. The `show -pause` command also allows the user to enter an interactive shell to further investigate the circuit before continuing synthesis.

So this script, when executed, will show the design after each of the three synthesis commands. The generated circuit diagrams are shown in Fig. 2.

The first diagram (from top to bottom) shows the design directly after being read by the Verilog front-end. Input and output ports are displayed as octagonal shapes. Cells are displayed as rectangles with inputs on the left and outputs on the right side. The cell labels are two lines long: The first line contains a unique identifier for the cell and the second line contains the cell type. Internal cell types are prefixed with a dollar sign. The Yosys manual contains a chapter on the internal cell library used in Yosys.

Constants are shown as ellipses with the constant value as label. The syntax `<bit_width>'<bits>` is used for constants that are not 32-bit wide and/or contain bits that are not 0 or 1 (i.e. x or z). Ordinary 32-bit constants are written using decimal numbers.

Single-bit signals are shown as thin arrows pointing from the driver to the load. Signals that are multiple bits wide are shown as thick arrows.

Finally *processes* are shown in boxes with round corners. Processes are Yosys' internal representation of the decision-trees and synchronization events modelled in a Verilog `always`-block. The label reads PROC followed by a unique identifier in the first line and contains the source code location of the original `always`-block in the 2nd line. Note how the multiplexer from the `?:`-expression is represented as a `$mux` cell but the multiplexer from the `if`-statement is yet still hidden within the process.

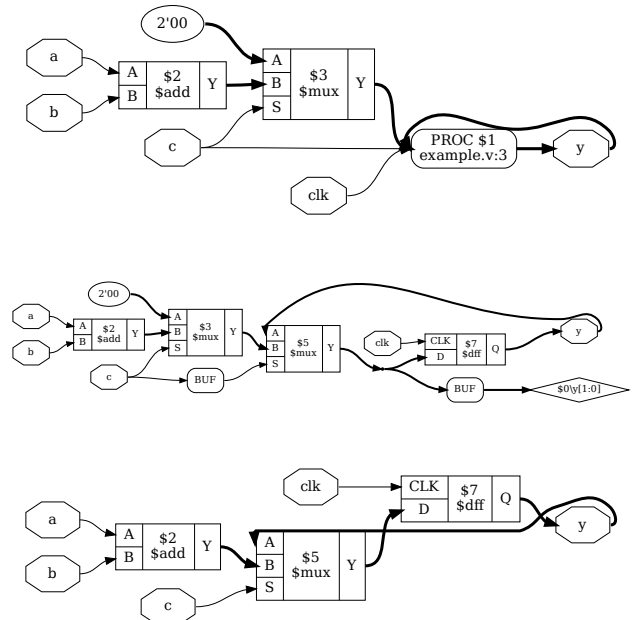


Figure 2. Output of the three `show` commands from Fig. 1

The `proc` command transforms the process from the first diagram into a multiplexer and a d-type flip-flop, which brings us to the 2nd diagram.

The Rhombus shape to the right is a dangling wire. (Wire nodes are only shown if they are dangling or have “public” names, for example names assigned from the Verilog input.) Also note that the design now contains two instances of a `BUF`-node. This are artefacts left behind by the `proc`-command. It is quite usual to see such artefacts after calling commands that perform changes in the design, as most commands only care about doing the transformation in the least complicated way, not about cleaning up after them. The next call to `clean` (or `opt`, which includes `clean` as one of its operations) will clean up this artefacts. This operation is so common in Yosys scripts that it can simply be abbreviated with the `;` token, which doubles as separator for commands. Unless one wants to specifically analyze this artefacts left behind some operations, it is therefore recommended to always call `clean` before calling `show`.

In this script we directly call `opt` as next step, which finally leads us to the 3rd diagram in Fig. 2. Here we see that the `opt` command not only has removed the artifacts left behind by `proc`, but also determined correctly that it can remove the first `$mux` cell without changing the behavior of the circuit.

### B. Break-out boxes for signal vectors

As has been indicated by the last example, Yosys is can manage signal vectors (aka. multi-bit wires or buses) as native objects. This provides great advantages when analyzing circuits that operate on wide integers. But it also introduces some additional complexity when the individual bits of of a signal vector are accessed. The example show in Fig. 3 and 4 demonstrates how such circuits are visualized by the `show` command.

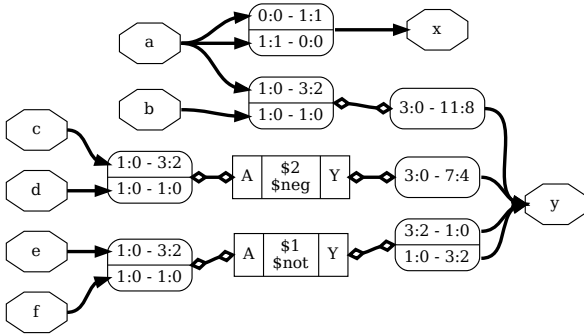


Figure 3. Output of `yosys -p 'proc; opt; show' splice.v`

```

1 module splice_demo(a, b, c, d, e, f, x, y);
2
3 input [1:0] a, b, c, d, e, f;
4 output [1:0] x = {a[0], a[1]};
5
6 output [11:0] y;
7 assign {y[11:4], y[1:0], y[3:2]} =
8     {a, b, ~{c, d}, ~{e, f}};
9
10 endmodule

```

Figure 4. `splice.v`

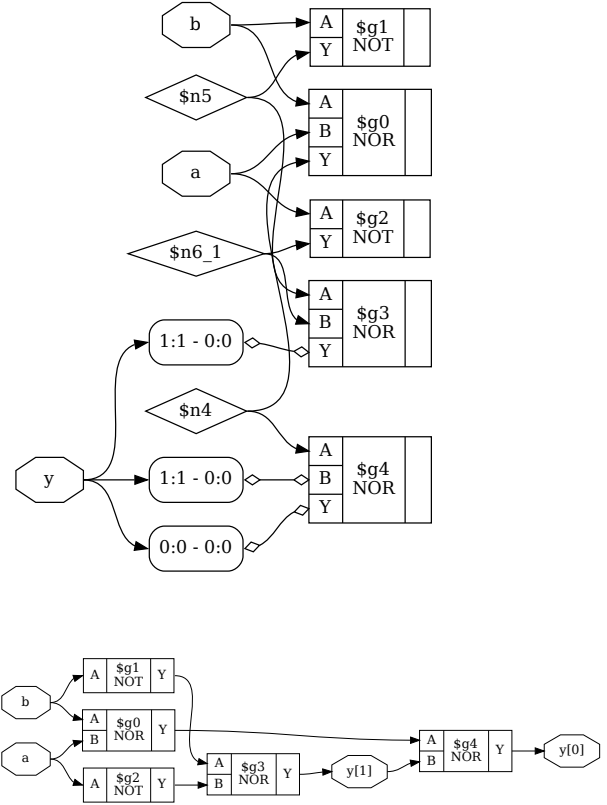


Figure 5. Effects of `splitnets` command and of providing a cell library. (The circuit is a half-adder built from simple CMOS gates.)

The key elements in understanding this circuit diagram are of course the boxes with round corners and rows labeled `<MSB_LEFT>:<LSB_LEFT>` - `<MSB_RIGHT>:<LSB_RIGHT>`. Each of this boxes has one signal per row on one side and a common signal for all rows on the other side. The `<MSB>:<LSB>` tuples specify which bits of the signals are broken out and connected. So the top row of the box connecting the signals `a` and `x` indicates that the bit 0 (i.e. the range 0:0) from signal `a` is connected to bit 1 (i.e. the range 1:1) of signal `x`.

Lines connecting such boxes together and lines connecting such boxes to cell ports have a slightly different look to emphasise that they are not actual signal wires but a necessity of the graphical representation. This distinction seems like a technicality, until one wants to debug a problem related to the way Yosys internally represents signal vectors, for example when writing custom Yosys commands.

### C. Gate level netlists

Finally Fig. 5 shows two common pitfalls when working with designs mapped to a cell library. The top figure has two problems: First Yosys did not have access to the cell library when this diagram was generated, resulting in all cell ports defaulting to being inputs. This is why all ports are drawn on the left side the cells are awkwardly arranged in a large column. Secondly the two-bit vector `y` requires breakout-boxes for its individual bits, resulting in an unnecessary complex diagram.

For the 2nd diagram Yosys has been given a description of the cell library as Verilog file containing blackbox modules. There are two ways to load cell descriptions into Yosys: First the Verilog file for the cell library can be passed directly to the `show` command using the `-lib <filename>` option. Secondly it is possible to load cell libraries into the design with the `read_verilog -lib <filename>` command. The 2nd method has the great advantage that the library only needs to be loaded once and can then be used in all subsequent calls to the `show` command.

In addition to that, the 2nd diagram was generated after `splitnet -ports` was run on the design. This command splits all signal vectors into individual signal bits, which is often desirable when looking at gate-level circuits. The `-ports` option is required to also split module ports. Per default the command only operates on interior signals.

#### D. Miscellaneous notes

Per default the `show` command outputs a temporary dot file and launches `xdot` to display it. The options `-format`, `-viewer` and `-prefix` can be used to change format, viewer and filename prefix. Note that the `pdf` and `ps` format are the only formats that support plotting multiple modules in one run.

In densely connected circuits it is sometimes hard to keep track of the individual signal wires. For this cases it can be useful to call `show` with the `-colors <integer>` argument, which randomly assigns colors to the nets. The integer ( $> 0$ ) is used as seed value for the random color assignments. Sometimes it is necessary it try some values to find an assignment of colors that looks good.

The command `help show` prints a complete listing of all options supported by the `show` command.

### IV. NAVIGATING THE DESIGN

Plotting circuit diagrams for entire modules in the design brings us only helps in simple cases. For complex modules the generated circuit diagrams are just stupidly big and are no help at all. In such cases one first has to select the relevant portions of the circuit.

In addition to *what* to display one also needs to carefully decide *when* to display it, with respect to the synthesis flow. In general it is a good idea to troubleshoot a circuit in the earliest state in which a problem can be reproduced. So if, for example, the internal state before calling the `techmap` command already fails to verify, it is better to troubleshoot the coarse-grain version of the circuit before `techmap` than the gate-level circuit after `techmap`.

Note: It is generally recommended to verify the internal state of a design by writing it to a Verilog file using `write_verilog -noexpr` and using the simulation models from `simlib.v` and `simcells.v` from the Yosys data directory (as printed by `yosys-config --datdir`).

#### A. Interactive Navigation

Once the right state within the synthesis flow for debugging the circuit has been identified, it is recommended to simply add the `shell` command to the matching place in the synthesis script. This command will stop the synthesis at the specified moment and go to shell mode, where the user can interactively enter commands.

For most cases, the shell will start with the whole design selected (i.e. when the synthesis script does not already narrow the selection). The command `ls` can now be used to create a list of all modules. The command `cd` can be used to switch to one of the modules (type `cd ..` to switch back). Now the `ls` command lists the objects within that module. Fig. 6 demonstrates this using the design from Fig. 1.

There is a thing to note in Fig. 6: We can see that the cell names from Fig. 2 are just abbreviations of the actual cell names, namely

```

1  yosys> ls
2
3  1 modules:
4      example
5
6  yosys> cd example
7
8  yosys [example]> ls
9
10 7 wires:
11  $0\y[1:0]
12  $add$example.v:5$2_Y
13  a
14  b
15  c
16  clk
17  y
18
19 3 cells:
20  $add$example.v:5$2
21  $procdff$7
22  $procmux$5

```

Figure 6. Demonstration of `ls` and `cd` using `example.v` from Fig. 1

the part after the last dollar-sign. Most auto-generated names (the ones starting with a dollar sign) are rather long and contains some additional information on the origin of the named object. But in most cases those names can simply be abbreviated using the last part.

Usually all interactive work is done with one module selected using the `cd` command. But it is also possible to work from the design-context (`cd ..`). In this case all object names must be prefixed with `<module_name>/. For example a*/b* would refer to all objects whose names start with b from all modules whose names start with a.`

The `dump` command can be used to print all information about an object. For example `dump $2` will print Fig. 7. This can for example be useful to determine the names of nets connected to cells, as the net-names are usually suppressed in the circuit diagram if they are auto-generated.

For the remainder of this document we will assume that the commands are run from module-context and not design-context.

#### B. Working with selections

When a module is selected using the `cd` command, all commands (with a few exceptions, such as the `read_*` and `write_*` commands) operate only on the selected module. This can also be useful for

```

1  attribute \src "example.v:5"
2  cell $add $add$example.v:5$2
3      parameter \A_SIGNED 0
4      parameter \A_WIDTH 1
5      parameter \B_SIGNED 0
6      parameter \B_WIDTH 1
7      parameter \Y_WIDTH 2
8      connect \A \a
9      connect \B \b
10     connect \Y $add$example.v:5$2_Y
11 end

```

Figure 7. Output of `dump $2` using the design from Fig. 1 and Fig. 2

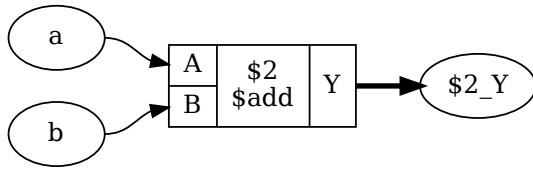


Figure 8. Output of show after select \$2 or select t:\$add (see also Fig. 2)

synthesis scripts where different synthesis strategies should be applied to different modules in the design.

But for most interactive work we want to further narrow the set of selected objects. This can be done using the `select` command.

For example, if the command `select $2` is executed, a subsequent `show` command will yield the diagram shown in Fig. 8. Note that the nets are now displayed in ellipses. This indicates that they are not selected, but only shown because the diagram contains a cell that is connected to the net. This of course makes no difference for the circuit that is shown, but it can be a useful information when manipulating selections.

Objects can not only be selected by their name but also by other properties. For example `select t:$add` will select all cells of type `$add`. In this case this is also yields the diagram shown in Fig. 8.

The output of `help select` contains a complete syntax reference for matching different properties.

Many commands can operate on explicit selections. For example the command `dump t:$add` will print information on all `$add` cells in the active module. Whenever a command has `[selection]` as last argument in its usage help, this means that it will use the engine behind the `select` command to evaluate additional arguments and use the resulting selection instead of the selection created by the last `select` command.

Normally the `select` command overwrites a previous selection. The commands `select -add` and `select -del` can be used to add or remove objects from the current selection.

The command `select -clear` can be used to reset the selection to the default, which is a complete selection of everything in the current module.

### C. Operations on selections

The `select` command is actually much more powerful than it might seem on the first glimpse. When it is called with multiple arguments, each argument is evaluated and pushed separately on a stack. After all arguments have been processed it simply creates the union of all

```

1 module foobaraddsub(a, b, c, d, fa, fs, ba, bs);
2   input [7:0] a, b, c, d;
3   output [7:0] fa, fs, ba, bs;
4   assign fa = a + (* foo *) b;
5   assign fs = a - (* foo *) b;
6   assign ba = c + (* bar *) d;
7   assign bs = c - (* bar *) d;
8 endmodule

```

Figure 9. Test module for operations on selections

```

1 module sumprod(a, b, c, sum, prod);
2
3   input [7:0] a, b, c;
4   output [7:0] sum, prod;
5
6   (* sumstuff *)
7   assign sum = a + b + c;
8   (* *)
9
10  assign prod = a * b * c;
11
12 endmodule

```

Figure 10. Another test module for operations on selections

elements on the stack. So the following command will select all `$add` cells and all objects with the `foo` attribute set:

```
select t:$add a:foo
```

(Try this with the design shown in Fig. 9. Use the `select -list` command to list the current selection.)

In many cases simply adding more and more stuff to the selection is an ineffective way of selecting the interesting part of the design. Special arguments can be used to combine the elements on the stack. For example the `%i` argument pops the last two elements from the stack, intersects them, and pushes the result back on the stack. So the following command will select all `$add` cells that have the `foo` attribute set:

```
select t:$add a:foo %i
```

The listing in Fig. 10 uses the Yosys non-standard `{* ... *}` syntax to set the attribute `sumstuff` on all cells generated by the first `assign` statement. (This works on arbitrary large blocks of Verilog code and can be used to mark portions of code for analysis.)

Selecting `a:sumstuff` in this module will yield the circuit diagram shown in Fig. 11. As only the cells themselves are selected, but not the temporary wire `$1_Y`, the two adders are shown as two disjunct parts. This can be very useful for global signals like clock and reset signals: just unselect them using a command such as `select -del clk rst` and each cell using them will get its own net label.

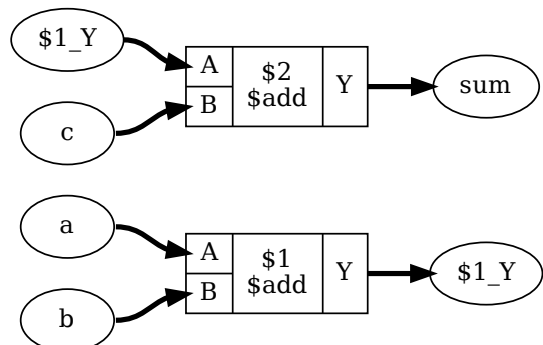
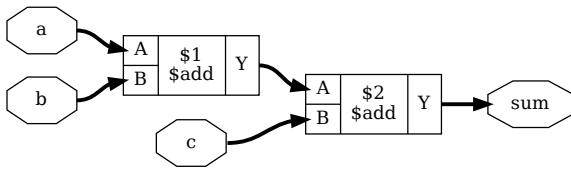


Figure 11. Output of show a:sumstuff on Fig. 10

Figure 12. Output of `show a:sumstuff %x` on Fig. 10

In this case however we would like to see the cells connected properly. This can be achieved using the `%x` action, that broadens the selection, i.e. for each selected wire it selects all cells connected to the wire and vice versa. So `show a:sumstuff %x` yields the diagram shown in Fig. 12.

#### D. Selecting logic cones

Fig. 12 shows what is called the *input cone* of `sum`, i.e. all cells and signals that are used to generate the signal `sum`. The `%ci` action can be used to select the input cones of all object in the top selection in the stack maintained by the `select` command.

As the `%x` action, this commands broadens the selection by one “step”. But this time the operation only works against the direction of data flow. That means, wires only select cells via output ports and cells only select wires via input ports.

Fig. 13 show the sequence of diagrams generated by the following commands:

```
show prod
show prod %ci
show prod %ci %ci
show prod %ci %ci %ci
```

When selecting many levels of logic, repeating `%ci` over and over again can be a bit dull. So there is a shortcut for that: the number of iterations can be appended to the action. So for example the action `%ci3` is identical to performing the `%ci` action three times.

The action `%ci*` performs the `%ci` action over and over again until it has no effect anymore.

In most cases there are certain cell types and/or ports that should not be considered for the `%ci` action, or we only want to follow certain cell types and/or ports. This can be achieved using additional patterns that can be appended to the `%ci` action.

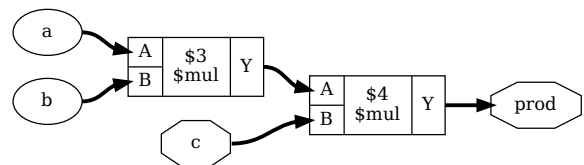
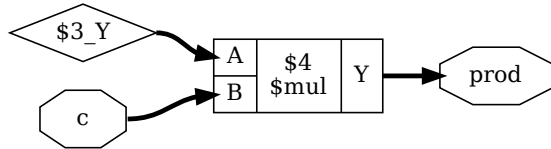
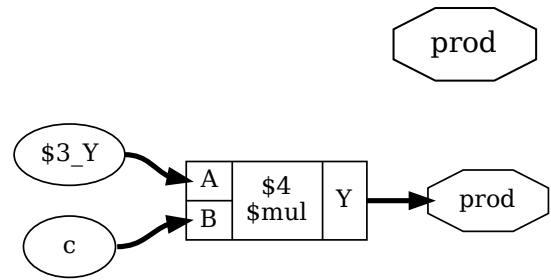
Lets consider the design from Fig. 14. It serves no purpose other than being a non-trivial circuit for demonstrating some of the advanced Yosys features. We synthesize the circuit using `proc; opt; memory; opt` and change to the `memdemo` module with `cd memdemo`. If we type `show now` we see the diagram shown in Fig. 15.

But maybe we are only interested in the tree of multiplexers that select the output value. In order to get there, we would start by just showing the output signal and its immediate predecessors:

```
show y %ci2
```

From this we would learn that `y` is driven by a `$dff` cell, that `y` is connected to the output port `Q`, that the `clk` signal goes into the `CLK` input port of the cell, and that the data comes from a auto-generated wire into the input `D` of the flip-flop cell.

As we are not interested in the clock signal we add an additional pattern to the `%ci` action, that tells it to only follow ports `Q` and `D` of `$dff` cells:

Figure 13. Objects selected by `select prod %ci...`

```
show y %ci2:+$dff[Q,D]
```

To add a pattern we add a colon followed by the pattern to the `%ci` action. The pattern it self starts with `-` or `+`, indicating if it is an include or exclude pattern, followed by an optional comma separated list of cell types, followed by an optional comma separated list of port names in square brackets.

Since we know that the only cell considered in this case is a `$dff` cell, we could as well only specify the port names:

```
show y %ci2:+[Q,D]
```

```
1 module memdemo(clk, d, y);
2
3 input clk;
4 input [3:0] d;
5 output reg [3:0] y;
6
7 integer i;
8 reg [1:0] s1, s2;
9 reg [3:0] mem [0:3];
10
11 always @(posedge clk) begin
12     for (i = 0; i < 4; i = i+1)
13         mem[i] <= mem[(i+1) % 4] + mem[(i+2) % 4];
14     { s2, s1 } = d ? { s1, s2 } ^ d : 4'b0;
15     mem[s1] <= d;
16     y <= mem[s2];
17 end
18
19 endmodule
```

Figure 14. Demo circuit for demonstrating some advanced Yosys features



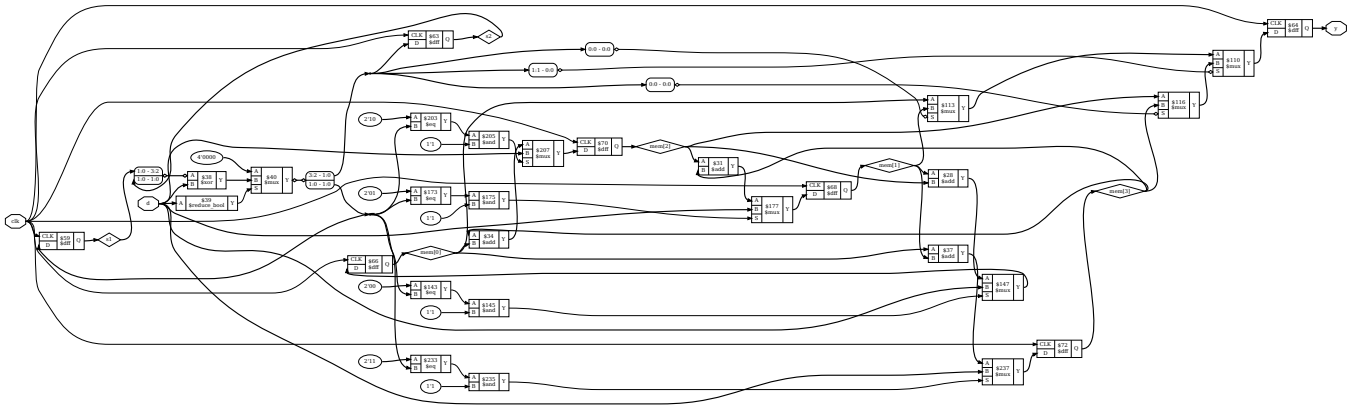


Figure 15. Complete circuit diagram for the design shown in Fig. 14

Or we could decide to tell the %ci action to not follow the CLK input:

```
show y %ci2:-[CLK]
```

Next we would investigate the next logic level by adding another %ci2 to the command:

```
show y %ci2:-[CLK] %ci2
```

From this we would learn that the next cell is a \$mux cell and we would add additional pattern to narrow the selection on the path we are interested. In the end we would end up with a command such as

```
show y %ci2:+$dff[Q,D] %ci*:-$mux[S]:-$dff
```

in which the first %ci jumps over the initial d-type flip-flop and the 2nd action selects the entire input cone without going over multiplexer select inputs and flip-flop cells. The diagram produces by this command is shown in Fig. 16.

Similar to %ci exists an action %co to select output cones that accepts the same syntax for pattern and repetition. The %x action mentioned previously also accepts this advanced syntax.

This actions for traversing the circuit graph, combined with the actions for boolean operations such as intersection (%i) and difference (%d) are powerful tools for extracting the relevant portions of the circuit under investigation.

See help select for a complete list of actions available in selections.

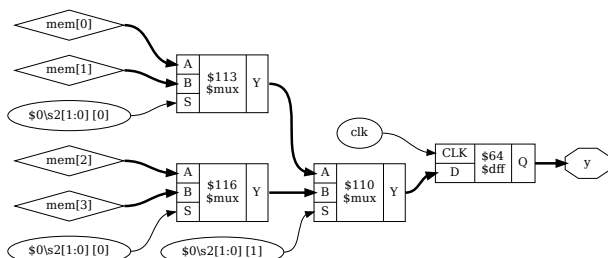


Figure 16. Output of show y %ci2:+\$dff[Q,D] %ci\*:-\$mux[S]:-\$dff

### E. Storing and recalling selections

The current selection can be stored in memory with the command `select -set <name>`. It can later be recalled using `select @<name>`. In fact, the @<name> expression pushes the stored selection on the stack maintained by the select command. So for example

```
select @foo @bar %i
```

will select the intersection between the stored selections `foo` and `bar`.

In larger investigation efforts it is highly recommended to maintain a script that sets up relevant selections, so they can easily be recalled, for example when Yosys needs to be re-run after a design or source code change.

The history command can be used to list all recent interactive commands. This feature can be useful for creating such a script from the commands used in an interactive session.

## V. ADVANCED INVESTIGATION TECHNIQUES

When working with very large modules, it is often not enough to just select the interesting part of the module. Instead it can be useful to extract the interesting part of the circuit into a separate module. This can for example be useful if one wants to run a series of synthesis commands on the critical part of the module and wants to carefully read all the debug output created by the commands in order to spot a problem. This kind of troubleshooting is much easier if the circuit under investigation is encapsulated in a separate module.

Fig. 17 shows how the `submod` command can be used to split the circuit from Fig. 14 and 15 into its components. The `-name` option is used to specify the name of the new module and also the name of the new cell in the current module.

### A. Evaluation of combinatorial circuits

The `eval` command can be used to evaluate combinatorial circuits. For example (see Fig. 17 for the circuit diagram of `selstage`):

```
yosys [selstage]> eval -set s2,s1 4'b1001 -set d 4'hc -show n2 -show n1
```

```
9. Executing EVAL pass (evaluate the circuit given an input).
Full command line: eval -set s2,s1 4'b1001 -set d 4'hc -show n2 -show n1
Eval result: \n2 = 2'10.
Eval result: \n1 = 2'10.
```

So the `-set` option is used to set input values and the `-show` option is used to specify the nets to evaluate. If no `-show` option is specified, all selected output ports are used per default.

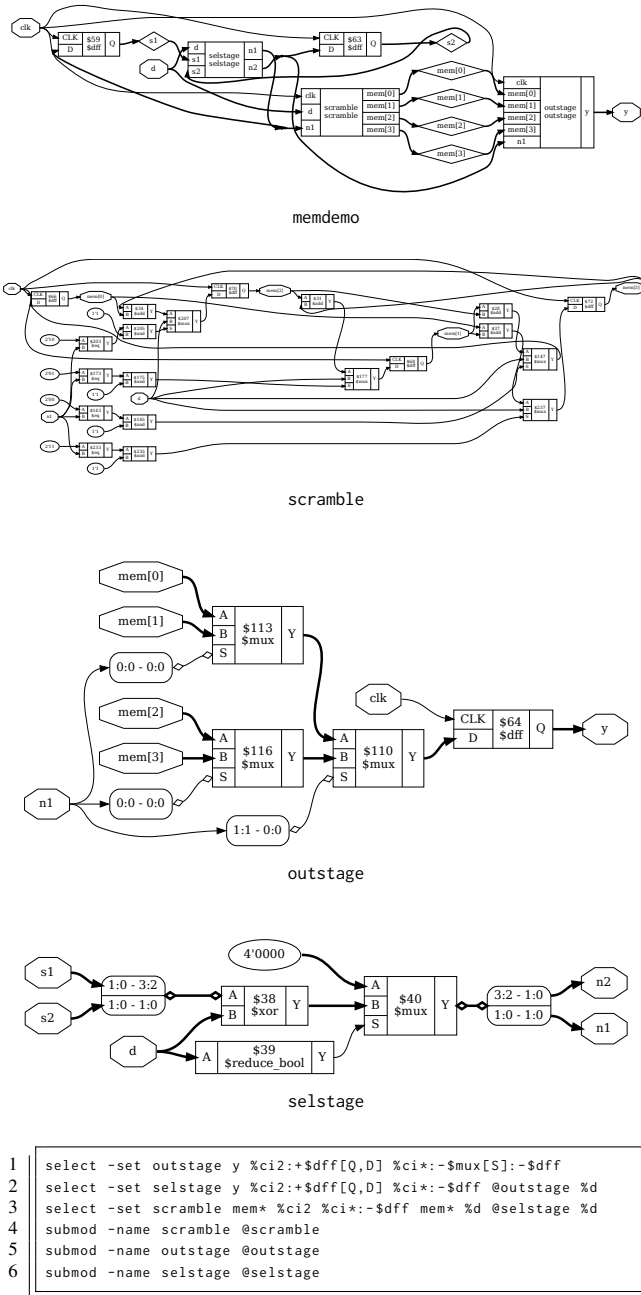


Figure 17. The circuit from Fig. 14 and 15 broken up using submod

If a necessary input value is not given, an error is produced. The option `-set-undef` can be used to instead set all unspecified input nets to undef (x).

The `-table` option can be used to create a truth table. For example:

```
yosys [selstage]> eval -set-undef -set d[3:1] 0 -table s1,d[0]
```

10. Executing EVAL pass (evaluate the circuit given an input).  
Full command line: eval -set-undef -set d[3:1] 0 -table s1,d[0]

```

\s1 \d [0] | \n1 \n2
-----|-----
2'00 1'0 | 2'00 2'00
2'00 1'1 | 2'xx 2'00
2'01 1'0 | 2'00 2'00
2'01 1'1 | 2'xx 2'01
2'10 1'0 | 2'00 2'00
2'10 1'1 | 2'xx 2'10
2'11 1'0 | 2'00 2'00

```

```
2'11 1'1 | 2'xx 2'11
```

Assumed undef (x) value for the following signals: \s2

Note that the eval command (as well as the sat command discussed in the next sections) does only operate on flattened modules. It can not analyze signals that are passed through design hierarchy levels. So the flatten command must be used on modules that instantiate other modules before this commands can be applied.

### B. Solving combinatorial SAT problems

Often the opposite of the eval command is needed, i.e. the circuits output is given and we want to find the matching input signals. For small circuits with only a few input bits this can be accomplished by trying all possible input combinations, as it is done by the eval -table command. For larger circuits however, Yosys provides the sat command that uses a SAT [4] solver [5] to solve this kind of problems.

The sat command works very similar to the eval command. The main difference is that it is now also possible to set output values and find the corresponding input values. For Example:

```
yosys [selstage]> sat -show s1,s2,d -set s1 s2 -set n2,n1 4'b1001
```

11. Executing SAT pass (solving SAT problems in the circuit).  
Full command line: sat -show s1,s2,d -set s1 s2 -set n2,n1 4'b1001

Setting up SAT problem:

```

Import set-constraint: \s1 = \s2
Import set-constraint: { \n2 \n1 } = 4'b1001
Final constraint equation: { \n2 \n1 \s1 } = { 4'b1001 \s2 }
Imported 3 cells to SAT database.
Import show expression: { \s1 \s2 \d }

```

Solving problem with 81 variables and 207 clauses..  
SAT solving finished - model found:

Signal Name	Dec	Hex	Bin
\d	9	9	1001
\s1	0	0	00
\s2	0	0	00

Note that the sat command supports signal names in both arguments to the `-set` option. In the above example we used `-set s1 s2` to constraint s1 and s2 to be equal. When more complex constraints are needed, a wrapper circuit must be constructed that checks the constraints and signals if the constraint was met using an extra output port, which then can be forced to a value using the `-set` option. (Such a circuit that contains the circuit under test plus additional constraint checking circuitry is called a *miter* circuit.)

Fig. 18 shows a miter circuit that is supposed to be used as a prime number test. If ok is 1 for all input values a and b for a given p, then p is prime, or at least that is the idea.

The Yosys shell session shown in Fig. 19 demonstrates that SAT solvers can even find the unexpected solutions to a problem: Using integer overflow there actually is a way of “factorizing” 31. The clean solution would of course be to perform the test in 32 bits, for example by replacing `p != a*b` in the miter with `p != {16'd0,a}*b`, or by using a temporary variable for the 32 bit product `a*b`. But as 31 fits well into 8 bits (and as the purpose of this document is to show off Yosys

```

1 module primetest(p, a, b, ok);
2   input [15:0] p, a, b;
3   output ok = p != a*b || a == 1 || b == 1;
4 endmodule

```

Figure 18. A simple miter circuit for testing if a number is prime. But it has a problem (see main text and Fig. 19).

```

1 yosys [primetest]> sat -prove ok 1 -set p 31
2
3 8. Executing SAT pass (solving SAT problems in the circuit).
4 Full command line: sat -prove ok 1 -set p 31
5
6 Setting up SAT problem:
7 Import set-constraint: \p = 16'0000000000011111
8 Final constraint equation: \p = 16'0000000000011111
9 Imported 6 cells to SAT database.
10 Import proof-constraint: \ok = 1'1
11 Final proof equation: \ok = 1'1
12
13 Solving problem with 2790 variables and 8241 clauses..
14 SAT proof finished - model found: FAIL!
15
16
17 (----- \      / ---) / ---) (-) |      | | |
18 (-----) )--- --- --- -| |--- -| |--- --- -| | | |
19 | ---/ ---) - \ / - (- ---) (- ---|--- | | | | --- | / - | -|
20 | | | | | | | | | | | | | | | | | | | | | | | | | | | | | |
21 | -| | -| \---/ \---/ | -| | -| \---| -| \---) \---| -|
22
23
24 Signal Name          Dec          Hex          Bin
25 -----
26 \a                    15029        3ab5        0011101010110101
27 \b                     4099        1003        0001000000000011
28 \ok                     0           0           0
29 \p                     31           1f        0000000000011111
30
31 yosys [primetest]> sat -prove ok 1 -set p 31 -set a[15:8],b[15:8] 0
32
33 9. Executing SAT pass (solving SAT problems in the circuit).
34 Full command line: sat -prove ok 1 -set p 31 -set a[15:8],b[15:8] 0
35
36 Setting up SAT problem:
37 Import set-constraint: \p = 16'0000000000011111
38 Import set-constraint: { \a [15:8] \b [15:8] } = 16'0000000000000000
39 Final constraint equation: { \a [15:8] \b [15:8] \p } = { 16'0000000000000000 16'0000000000011111 }
40 Imported 6 cells to SAT database.
41 Import proof-constraint: \ok = 1'1
42 Final proof equation: \ok = 1'1
43
44 Solving problem with 2790 variables and 8257 clauses..
45 SAT proof finished - no model found: SUCCESS!
46
47
48 /$$$$$ /$$$$$$ /$$$$$
49 /$__$ $$ | $$_---/ | $$_ $$
50 | $$ \ $$ | $$ | $$ \ $$
51 | $$ | $$ | $$$$ | $$ | $$
52 | $$ | $$ | $$_- / | $$ | $$
53 | $$/$$ $$ | $$ | $$ | $$
54 | $$$$$$ / $$ | $$$$$$ / $$ | $$$$$$ / $$
55 \___ $$_- / |___- / |___- / |___- /
    \___ /

```

Figure 19. Experiments with the miter circuit from Fig. 18. The first attempt of proving that 31 is prime failed because the SAT solver found a creative way of factorizing 31 using integer overflow.

features) we can also simply force the upper 8 bits of *a* and *b* to zero for the `sat` call, as is done in the second command in Fig. 19 (line 31).

The `-prove` option used in this example works similar to `-set`, but tries to find a case in which the two arguments are not equal. If such a case is not found, the property is proven to hold for all inputs that satisfy the other constraints.

It might be worth noting, that SAT solvers are not particularly

efficient at factorizing large numbers. But if a small factorization problem occurs as part of a larger circuit problem, the Yosys SAT solver is perfectly capable of solving it.

### C. Solving sequential SAT problems

The SAT solver functionality in Yosys can not only be used to solve combinatorial problems, but can also solve sequential problems. Let's

```

1  yosys [memdemo]> sat -seq 6 -show y -show d -set-init-undef \
2      -max_undef -set-at 4 y 1 -set-at 5 y 2 -set-at 6 y 3
3
4  6. Executing SAT pass (solving SAT problems in the circuit).
5  Full command line: sat -seq 6 -show y -show d -set-init-undef
6      -max_undef -set-at 4 y 1 -set-at 5 y 2 -set-at 6 y 3
7
8  Setting up time step 1:
9  Final constraint equation: { } = { }
10 Imported 29 cells to SAT database.
11
12 Setting up time step 2:
13 Final constraint equation: { } = { }
14 Imported 29 cells to SAT database.
15
16 Setting up time step 3:
17 Final constraint equation: { } = { }
18 Imported 29 cells to SAT database.
19
20 Setting up time step 4:
21 Import set-constraint for timestep: \y = 4'0001
22 Final constraint equation: \y = 4'0001
23 Imported 29 cells to SAT database.
24
25 Setting up time step 5:
26 Import set-constraint for timestep: \y = 4'0010
27 Final constraint equation: \y = 4'0010
28 Imported 29 cells to SAT database.
29
30 Setting up time step 6:
31 Import set-constraint for timestep: \y = 4'0011
32 Final constraint equation: \y = 4'0011
33 Imported 29 cells to SAT database.
34
35 Setting up initial state:
36 Final constraint equation: { \y \s2 \s1 \mem[3] \mem[2] \mem[1]
37     \mem[0] } = 24'xxxxxxxxxxxxxxxxxxxxxxxx
38
39 Import show expression: \y
40 Import show expression: \d
41
42 Solving problem with 10322 variables and 27881 clauses..
43 SAT model found, maximizing number of undefs.
44 SAT solving finished - model found:
45
46   Time Signal Name      Dec      Hex      Bin
47   -----
48   init \mem[0]          --      --      xxxx
49   init \mem[1]          --      --      xxxx
50   init \mem[2]          --      --      xxxx
51   init \mem[3]          --      --      xxxx
52   init \s1              --      --      xx
53   init \s2              --      --      xx
54   init \y                --      --      xxxx
55   -----
56   1 \d                   0       0       0000
57   1 \y                   --      --      xxxx
58   -----
59   2 \d                   1       1       0001
60   2 \y                   --      --      xxxx
61   -----
62   3 \d                   2       2       0010
63   3 \y                   0       0       0000
64   -----
65   4 \d                   3       3       0011
66   4 \y                   1       1       0001
67   -----
68   5 \d                   --      --      001x
69   5 \y                   2       2       0010
70   -----
71   6 \d                   --      --      xxxx
72   6 \y                   3       3       0011

```

Figure 20. Solving a sequential SAT problem in the memdemo module from Fig. 14.

consider the entire memdemo module from Fig. 14 and suppose we want to know which sequence of input values for *d* will cause the output *y* to produce the sequence 1, 2, 3 from any initial state. Fig. 20 show the solution to this question, as produced by the following command:

```
sat -seq 6 -show y -show d -set-init-undef \
```

```
-max_undef -set-at 4 y 1 -set-at 5 y 2 -set-at 6 y 3
```

The `-seq 6` option instructs the `sat` command to solve a sequential problem in 6 time steps. (Experiments with lower number of steps have show that at least 3 cycles are necessary to bring the circuit in a state from which the sequence 1, 2, 3 can be produced.)

The `-set-init-undef` option tells the `sat` command to initialize all registers to the undef (x) state. The way the x state is treated in Verilog will ensure that the solution will work for any initial state.

The `-max_undef` option instructs the `sat` command to find a solution with a maximum number of undefs. This way we can see clearly which inputs bits are relevant to the solution.

Finally the three `-set-at` options add constraints for the *y* signal to play the 1, 2, 3 sequence, starting with time step 4.

It is not surprising that the solution sets *d* = 0 in the first step, as this is the only way of setting the *s1* and *s2* registers to a known value. The input values for the other steps are a bit harder to work out manually, but the SAT solver finds the correct solution in an instant.

There is much more to write about the `sat` command. For example, there is a set of options that can be used to performs sequential proofs using temporal induction [6]. The command `help sat` can be used to print a list of all options with short descriptions of their functions.

## VI. CONCLUSION

Yosys provides a wide range of functions to analyze and investigate designs. For many cases it is sufficient to simply display circuit diagrams, maybe use some additional commands to narrow the scope of the circuit diagrams to the interesting parts of the circuit. But some cases require more than that. For this applications Yosys provides commands that can be used to further inspect the behavior of the circuit, either by evaluating which output values are generated from certain input values (eval) or by evaluation which input values and initial conditions can result in a certain behavior at the outputs (sat). The SAT command can even be used to prove (or disprove) theorems regarding the circuit, in more advanced cases with the additional help of a miter circuit.

This features can be powerful tools for the circuit designer using Yosys as a utility for building circuits and the software developer using Yosys as a framework for new algorithms alike.

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## Yosys Application Note 012: Converting Verilog to BTOR

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**Abstract**—Verilog-2005 is a powerful Hardware Description Language (HDL) that can be used to easily create complex designs from small HDL code. BTOR [3] is a bit-precise word-level format for model checking. It is a simple format and easy to parse. It allows to model the model checking problem over the theory of bit-vectors with one-dimensional arrays, thus enabling to model Verilog designs with registers and memories. Yosys [1] is an Open-Source Verilog synthesis tool that can be used to convert Verilog designs with simple assertions to BTOR format.

### I. INSTALLATION

Yosys written in C++ (using features from C++11) and is tested on modern Linux. It should compile fine on most UNIX systems with a C++11 compiler. The README file contains useful information on building Yosys and its prerequisites.

Yosys is a large and feature-rich program with some dependencies. For this work, we may deactivate other extra features such as TCL and ABC support in the Makefile.

This Application Note is based on GIT Rev. 082550f from 2015-04-04 of Yosys [1].

### II. QUICK START

We assume that the Verilog design is synthesizable and we also assume that the design does not have multi-dimensional memories. As BTOR implicitly initializes registers to zero value and memories stay uninitialized, we assume that the Verilog design does not contain initial blocks. For more details about the BTOR format, please refer to [3].

We provide a shell script `verilog2btor.sh` which can be used to convert a Verilog design to BTOR. The script can be found in the `backends/btor` directory. The following example shows its usage:

```
verilog2btor.sh fsm.v fsm.btor test
```

Listing 1. Using `verilog2btor` script

The script `verilog2btor.sh` takes three parameters. In the above example, the first parameter `fsm.v` is the input design, the second parameter `fsm.btor` is the file name of BTOR output, and the third parameter `test` is the name of top module in the design.

To specify the properties (that need to be checked), we have two options:

- We can use the Verilog `assert` statement in the procedural block or module body of the Verilog design, as shown in Listing 2. This is the preferred option.
- We can use a single-bit output wire, whose name starts with `safety`. The value of this output wire needs to be driven low when the property is met, i.e. the solver will try to find a model that makes the safety pin go high. This is demonstrated in Listing 3.

```
module test(input clk, input rst, output y);

    reg [2:0] state;

    always @(posedge clk) begin
        if (rst || state == 3) begin
            state <= 0;
        end else begin
            assert(state < 3);
            state <= state + 1;
        end
    end

    assign y = state[2];

    assert property (y != 1'b1);

endmodule
```

Listing 2. Specifying property in Verilog design with `assert`

```
module test(input clk, input rst,
            output y, output safety1);

    reg [2:0] state;

    always @(posedge clk) begin
        if (rst || state == 3)
            state <= 0;
        else
            state <= state + 1;
    end

    assign y = state[2];

    assign safety1 = !(y != 1'b1);

endmodule
```

Listing 3. Specifying property in Verilog design with output wire

We can run Boolector [2] 1.4.1<sup>1</sup> on the generated BTOR file:

```
$ boolector fsm.btor
unsat
```

Listing 4. Running boolector on BTOR file

We can also use nuXmv [4], but on BTOR designs it does not support memories yet. With the next release of nuXmv, we will be also able to verify designs with memories.

### III. DETAILED FLOW

Yosys is able to synthesize Verilog designs up to the gate level. We are interested in keeping registers and memories when synthesizing the design. For this purpose, we describe a customized Yosys synthesis flow, that is also provided by the `verilog2btor.sh` script. Listing 5 shows the Yosys commands that are executed by `verilog2btor.sh`.

<sup>1</sup>Newer version of Boolector do not support sequential models. Boolector 1.4.1 can be built with picosat-951. Newer versions of picosat have an incompatible API.

```

1 read_verilog -sv $1;
2 hierarchy -top $3; hierarchy -libdir $DIR;
3 hierarchy -check;
4 proc; opt;
5 opt_expr -mux_undef; opt;
6 rename -hide;;;
7 splice; opt;
8 memory_dff -wr_only; memory_collect;;
9 flatten;;
10 memory_unpack;
11 splitnets -driver;
12 setundef -zero -undriven;
13 opt;;;
14 write_btor $2;

```

Listing 5. Synthesis Flow for BTOR with memories

Here is short description of what is happening in the script line by line:

- 1) Reading the input file.
- 2) Setting the top module in the hierarchy and trying to read automatically the files which are given as include in the file read in first line.
- 3) Checking the design hierarchy.
- 4) Converting processes to multiplexers (muxs) and flip-flops.
- 5) Removing undef signals from muxs.
- 6) Hiding all signal names that are not used as module ports.
- 7) Explicit type conversion, by introducing slice and concat cells in the circuit.
- 8) Converting write memories to synchronous memories, and collecting the memories to multi-port memories.
- 9) Flattening the design to get only one module.
- 10) Separating read and write memories.
- 11) Splitting the signals that are partially assigned
- 12) Setting undef to zero value.
- 13) Final optimization pass.
- 14) Writing BTOR file.

For detailed description of the commands mentioned above, please refer to the Yosys documentation, or run `yosys -h command_name`.

The script presented earlier can be easily modified to have a BTOR file that does not contain memories. This is done by removing the line number 8 and 10, and introduces a new command `memory` at line number 8. Listing 6 shows the modified Yosys script file:

```

read_verilog -sv $1;
hierarchy -top $3; hierarchy -libdir $DIR;
hierarchy -check;
proc; opt;
opt_expr -mux_undef; opt;
rename -hide;;;
splice; opt;
memory;;
flatten;;
splitnets -driver;
setundef -zero -undriven;
opt;;;
write_btor $2;

```

Listing 6. Synthesis Flow for BTOR without memories

## IV. EXAMPLE

Here is an example Verilog design that we want to convert to BTOR:

```

module array(input clk);

    reg [7:0] counter;
    reg [7:0] mem [7:0];

    always @(posedge clk) begin
        counter <= counter + 8'd1;
        mem[counter] <= counter;
    end

    assert property (!(counter > 8'd0) ||
        mem[counter - 8'd1] == counter - 8'd1);

endmodule

```

Listing 7. Example - Verilog Design

The generated BTOR file that contain memories, using the script shown in Listing 5:

```

1 var 1 clk
2 array 8 3
3 var 8 $auto$rename.cc:150:execute$20
4 const 8 00000001
5 sub 8 3 4
6 slice 3 5 2 0
7 read 8 2 6
8 slice 3 3 2 0
9 add 8 3 4
10 const 8 00000000
11 ugt 1 3 10
12 not 1 11
13 const 8 11111111
14 slice 1 13 0 0
15 one 1
16 eq 1 1 15
17 and 1 16 14
18 write 8 3 2 8 3
19 acond 8 3 17 18 2
20 anex 8 3 2 19
21 eq 1 7 5
22 or 1 12 21
23 const 1 1
24 one 1
25 eq 1 23 24
26 cond 1 25 22 24
27 root 1 -26
28 cond 8 1 9 3
29 next 8 3 28

```

Listing 8. Example - Converted BTOR with memory

And the BTOR file obtained by the script shown in Listing 6, which expands the memory into individual elements:

```

1 var 1 clk
2 var 8 mem[0]
3 var 8 $auto$rename.cc:150:execute$20
4 slice 3 3 2 0
5 slice 1 4 0 0
6 not 1 5
7 slice 1 4 1 1
8 not 1 7
9 slice 1 4 2 2
10 not 1 9
11 and 1 8 10
12 and 1 6 11
13 cond 8 12 3 2
14 cond 8 1 13 2
15 next 8 2 14
16 const 8 00000001
17 add 8 3 16
18 const 8 00000000
19 ugt 1 3 18
20 not 1 19
21 var 8 mem[2]
22 and 1 7 10
23 and 1 6 22
24 cond 8 23 3 21
25 cond 8 1 24 21
26 next 8 21 25
27 sub 8 3 16
:
:
54 cond 1 53 50 52
55 root 1 -54
:
:
77 cond 8 76 3 44
78 cond 8 1 77 44
79 next 8 44 78

```

Listing 9. Example - Converted BTOR without memory

## V. LIMITATIONS

BTOR does not support initialization of memories and registers, i.e. they are implicitly initialized to value zero, so the initial block for memories need to be removed when converting to BTOR. It should also be kept in consideration that BTOR does not support the `x` or `z` values of Verilog.

Another thing to bear in mind is that Yosys will convert multi-dimensional memories to one-dimensional memories and address decoders. Therefore out-of-bounds memory accesses can yield unexpected results.

## VI. CONCLUSION

Using the described flow, we can use Yosys to generate word-level verification benchmarks with or without memories from Verilog designs.

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